

Towards theorem proving graph grammars

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Motivation

- * **Theorem proving** is a **powerful technique** for the analysis of computational systems
- * It is possible to prove properties of reachable states for **infinite state systems**
- * This approach is **complementary** to other existing analysis techniques for graph grammars, like model checking and analysis based on approximations

Personal motivation: I kept trying to convince colleagues that using formal methods is nice, they allow to prove many relevant properties, hopefully using some tool, but I myself never used theorem provers before...

But...

- * Theorem provers are typically hard to use because
 - the system under analysis must be faithfully encoded in a logical framework (in a rather low level way)
 - many existing tools are hard to install and use
 - proving properties needs a lot of user assistance
- * Graphs (with types, attributes, ...) are complex structures, and the properties we wish to prove easily become cryptic logical formulae

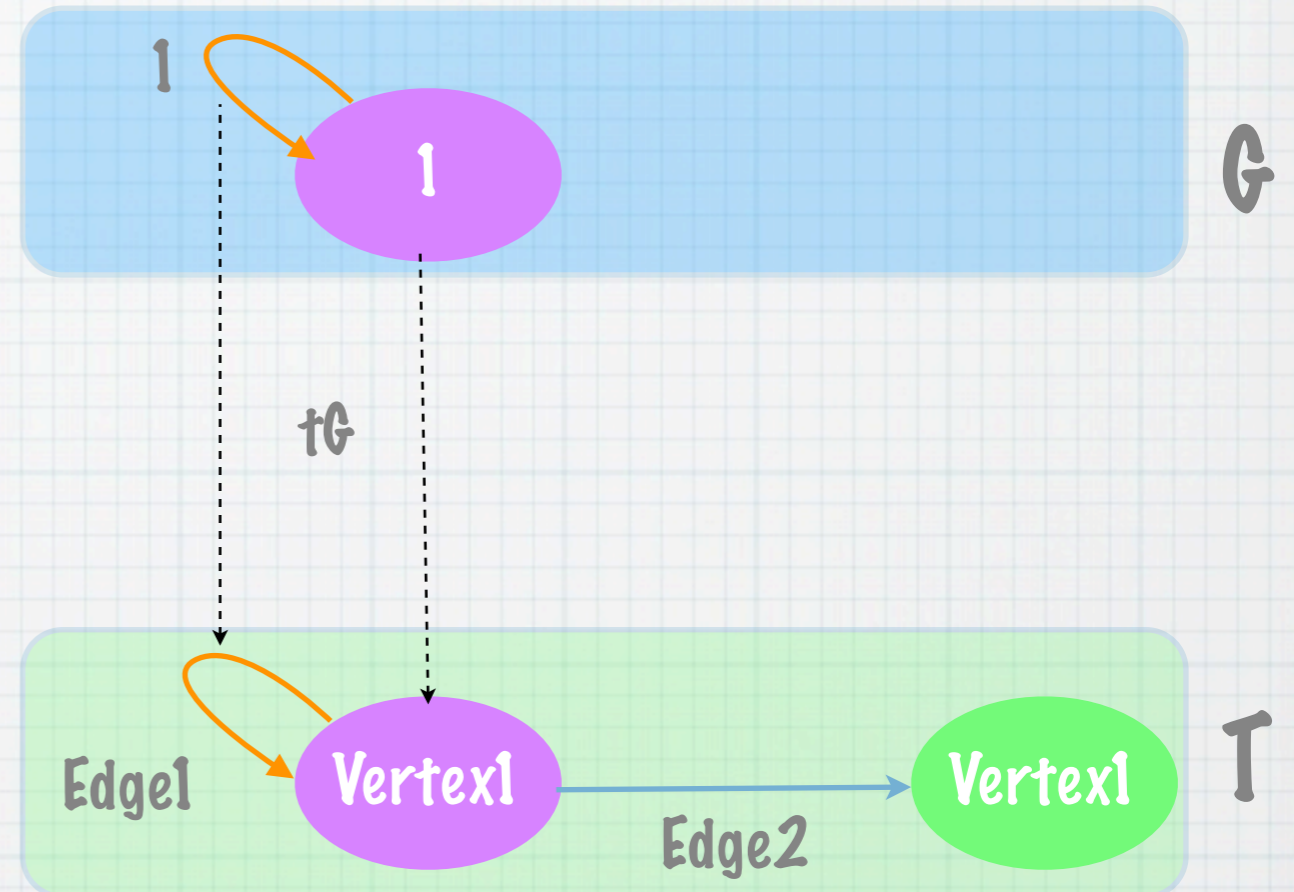
The idea...

- * Encode GGs in such a way they become closer to the input languages of theorem provers
- * Use existing theorem provers to prove properties of graph grammars

Typed Graphs

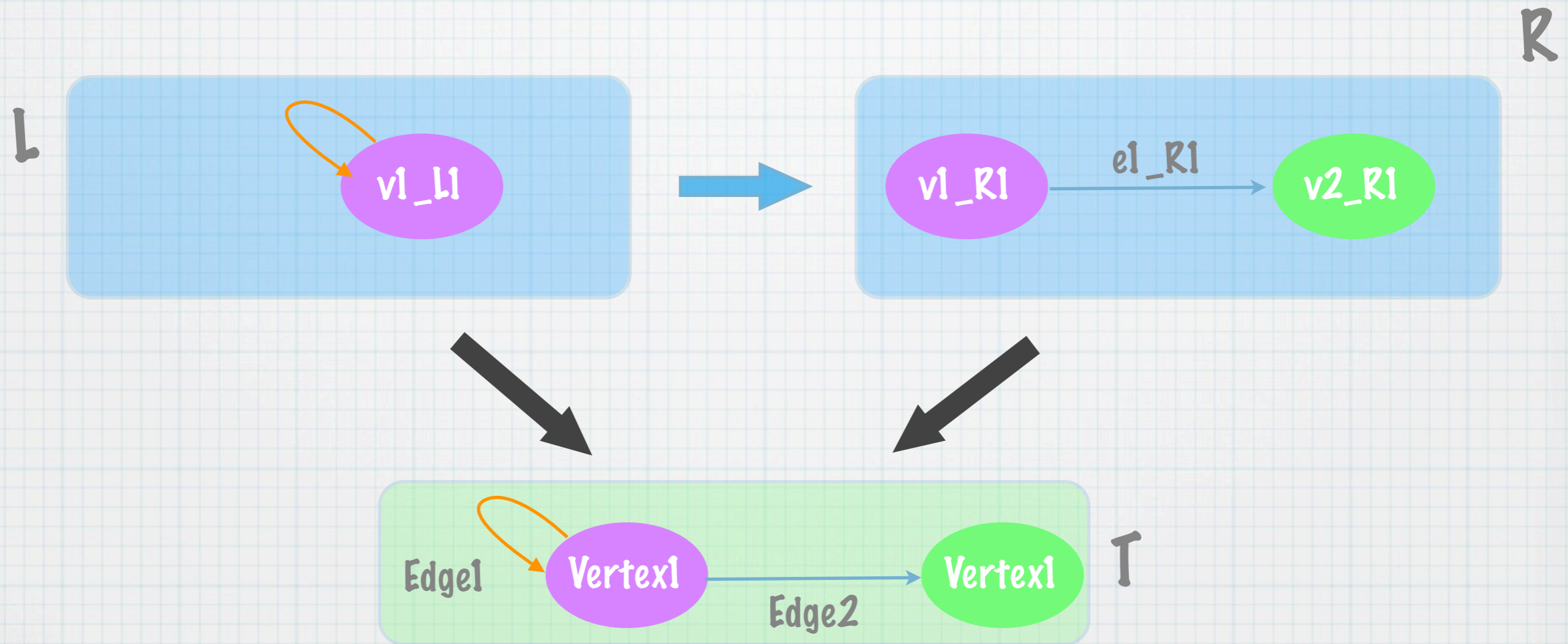
Typed Graph: tG

Type Graph: T



Category of typed graphs and partial graph morphisms : $\text{GraP}(T)$

(SPO) Rule



 : partial injective graph homomorphism, no vertex deletion

SPO Graph Grammar

- * Type Graph T
- * Initial (start) graph G typed over T
- * Set of rules typed over T

Semantics based on rule applications (POs in $\text{GraP}(T)$)

Relational Graph Grammars

- * Definition of **graph grammars** based on **relations and restrictions** on these relations
- * **Faithful encoding of SPO**, considering injective rules that do not delete vertices and matches that are injective on edges
- * Based on Courcelle's relational structures: domain+relations, transduction

Relational Graph

V_{GG} : set of vertex ids

E_{GG} : set of edge ids

$vertT \subseteq V_{GG}$

$edgeT \subseteq E_{GG}$

$sourceT \in edgeT \rightarrow vertT$

$targetT \in edgeT \rightarrow vertT$

$tG_V \in vertG \rightarrow vertT$

$tG_E \in edgeG \rightarrow edgeT$

source/target compat.

$vertG \subseteq V_{GG}$

$edgeG \subseteq E_{GG}$

$sourceG \in edgeG \rightarrow vertG$

$targetG \in edgeG \rightarrow vertG$

Graph T

tG

Graph G

Relational Rule

$$\alpha : L \longrightarrow R$$

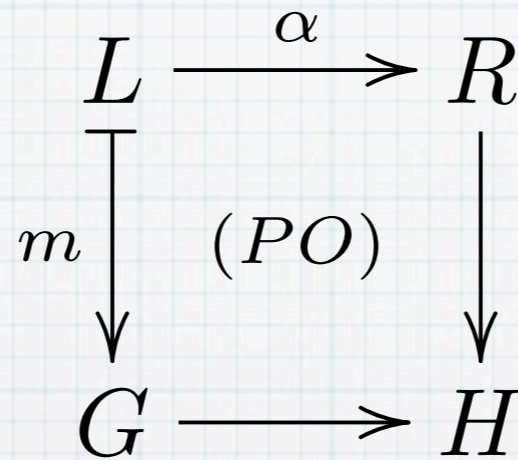
$$\alpha_V : \text{vert}L \rightsquigarrow \text{vert}R$$

$$\alpha_E : \text{edge}L \rightsquigarrow \text{edge}R$$

type compat.

Match is analogous, but total.

Rule Application



$$\text{vert}H = \text{vert}G \uplus (\text{vert}R - \alpha_V(\text{vert}L))$$

$$\text{edge}H = (\text{edge}G - m_E(\text{edge}L)) \uplus E_R$$

+ source, target, typing functions...

Now, implementing....

- * First attempt: Use Isabelle
- * Second attempt: Event-B (Rodin platform)

Event-B

- * Formalism based on B

- * Event-B models are composed by

- ▶ **CONTEXT** : definition of type graph and rules facts that may be used in the model

- ▶ **MACHINE** : definition of start graph and rule application (transitions)

Type Graph in Event-B

CONTEXT ctx_gg

SETS

vertT (Type Graph T) Verteces

edgeT (Type Graph T) Edges

CONSTANTS

sourceT

targetT

Vertex1

Edge1

Vertex2

Edge2

AXIOMS

axm1 : $partition(vertT, \{Vertex1\}, \{Vertex2\})$

axm2 : $partition(edgeT, \{Edge1\}, \{Edge2\})$

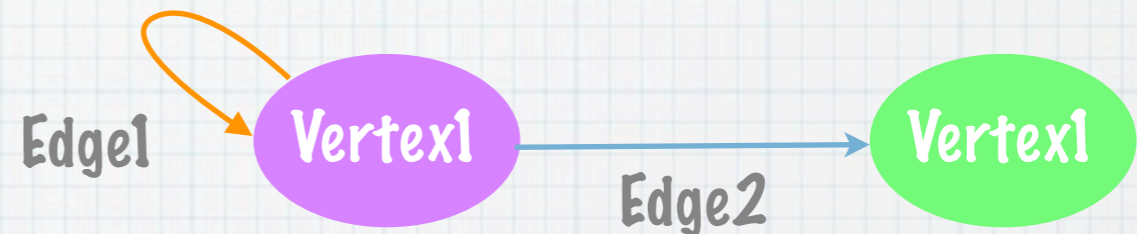
axm_src : $sourceT \in edgeT \rightarrow vertT$

axm_srcD : $partition(sourceT, \{Edge1 \mapsto Vertex1\}, \{Edge2 \mapsto Vertex1\})$

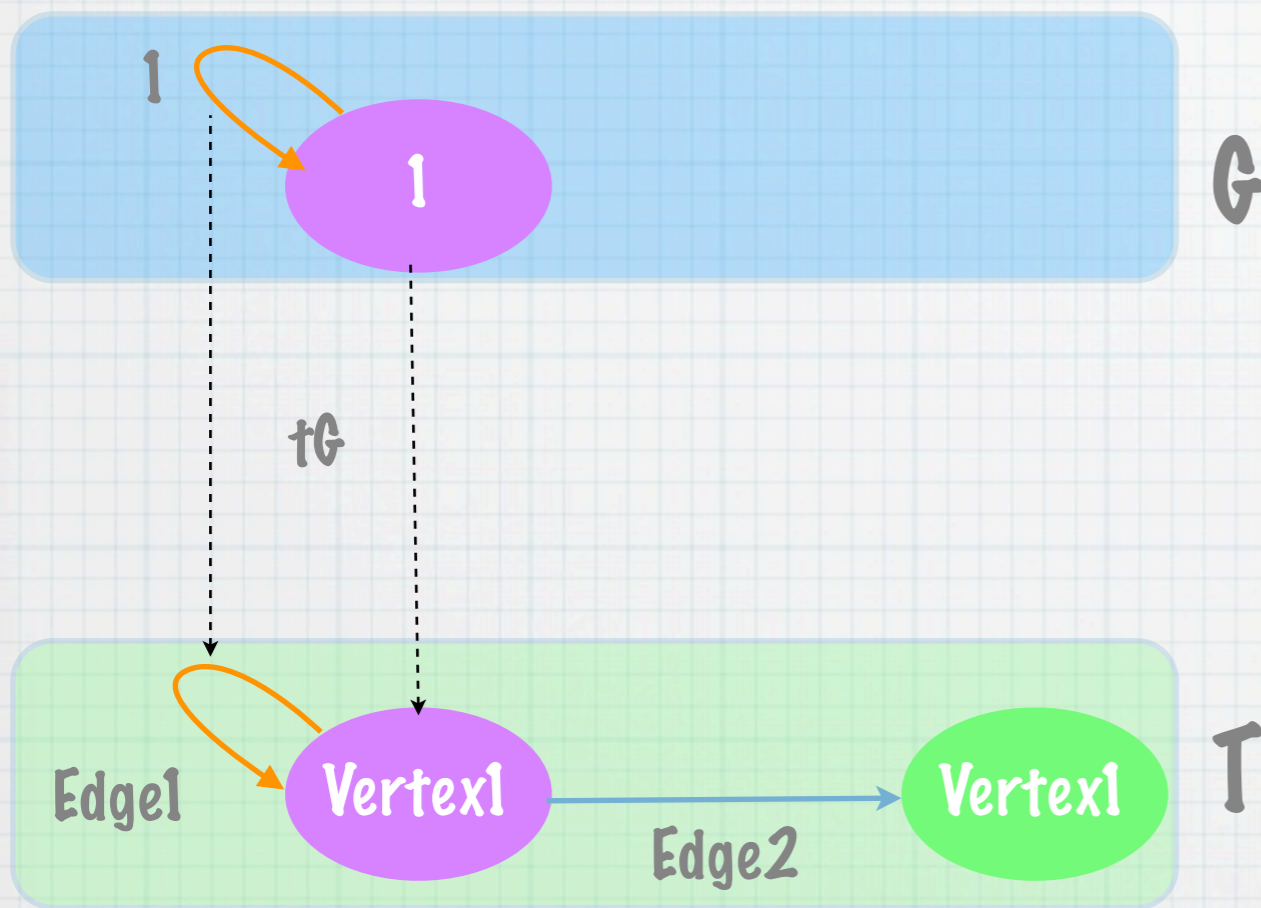
axm_trg : $targetT \in edgeT \rightarrow vertT$

axm10 : $partition(targetT, \{Edge1 \mapsto Vertex1\}, \{Edge2 \mapsto Vertex2\})$

END



State Graph



MACHINE mch_gg

SEES ctx_gg

VARIABLES

vertG

edgeG

sourceG

targetG

tG_V Typing vertices, tG_V

tG_E Typing edges, tG_E

INVARIANTS

type_vertG : $vertG \in \mathbb{P}(\mathbb{N})$

type_edgeG : $edgeG \in \mathbb{P}(\mathbb{N})$

type_sourceG : $sourceG \in edgeG \rightarrow vertG$

type_targetG : $targetG \in edgeG \rightarrow vertG$

type_tG_V : $tG_V \in vertG \rightarrow vertT$

type_tG_E : $tG_E \in edgeG \rightarrow edgeT$

EVENTS

Initialisation

begin

act1 : $vertG := \{1\}$

actE : $edgeG := \{1\}$

acts : $sourceG := \{1 \mapsto 1\}$

actt : $targetG := \{1 \mapsto 1\}$

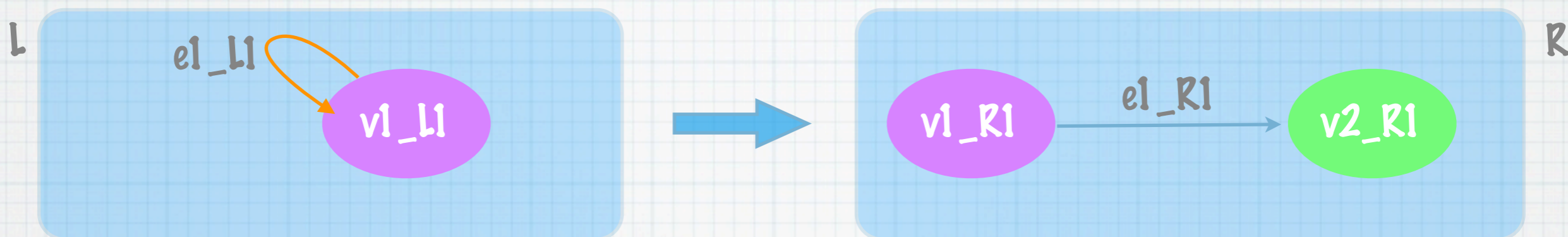
act3 : $tG_V := \{1 \mapsto Vertex1\}$

act4 : $tG_E := \{1 \mapsto Edge1\}$

end

END

Rule



CONTEXT ctx_gg

SETS

vertL1
edgeL1
vertR1
edgeR1

CONSTANTS

v1_L1
e1_L1
v1_R1
v2_R1
e1_R1
tL1_V (LHS1) Typing vertices, tL1_V
tL1_E (LHS1) Typing edges, tL1_E
tR1_V (RHS1) Typing vertices, tR1_V
tR1_E (RHS1) Typing edges, tR1_E
alpha1V (Rule 1) Rule alpha1: mapping vertices
alpha1E (Rule 1) Rule alpha1: mapping edges

AXIOMS

axm15 : $partition(vertL1, \{v1_L1\})$
 axm16 : $partition(edgeL1, \{e1_L1\})$
 axm_srcL1 : $sourceL1 \in edgeL1 \rightarrow vertL1$
 axm_srcDL1 : $partition(sourceL1, \{e1_L1 \mapsto v1_L1\})$
 axm_tarL1 : $targetL1 \in edgeL1 \rightarrow vertL1$
 axmtarDL1 : $partition(targetL1, \{e1_L1 \mapsto v1_L1\})$
 axmtL1V : $tL1_V \in vertL1 \rightarrow vertT$
 act9 : $partition(tL1_V, \{v1_L1 \mapsto Vertex1\})$
 axmtL1E : $tL1_E \in edgeL1 \rightarrow edgeT$
 act10 : $partition(tL1_E, \{e1_L1 \mapsto Edge1\})$
 axm17 : $partition(vertR1, \{v1_R1\}, \{v2_R1\})$
 axm18 : $partition(edgeR1, \{e1_R1\})$
 axm_srcR1 : $sourceR1 \in edgeR1 \rightarrow vertR1$
 axm_srcDR1 : $partition(sourceR1, \{e1_R1 \mapsto v1_R1\})$
 axm_tarR1 : $targetR1 \in edgeR1 \rightarrow vertR1$
 axmrafDR1 : $partition(targetR1, \{e1_R1 \mapsto v2_R1\})$
 axmtR1 : $tR1_V \in vertR1 \rightarrow vertT$
 act13 : $partition(tR1_V, \{v1_R1 \mapsto Vertex1\}, \{v2_R1 \mapsto Vertex2\})$
 axmtR1E : $tR1_E \in edgeR1 \rightarrow edgeT$
 act14 : $partition(tR1_E, \{e1_R1 \mapsto Edge2\})$
 axmA1V : $alpha1V \in vertL1 \mapsto vertR1$
 act15 : $partition(alpha1V, \{v1_L1 \mapsto v1_R1\})$
 axmA1E : $alpha1E \in edgeL1 \mapsto edgeR1$
 act16 : $alpha1E = \emptyset$

Rule (Behavior)

EVENTS

Event $\alpha_1 \hat{=}$

any

mV

mE

$newV$

$newE$

where

$grd1 : mV \in vertL1 \rightarrow vertG$

$grd2 : mE \in edgeL1 \mapsto edgeG$

$grd3 : newV \in \mathbb{N} \setminus vertG$

$grd4 : newE \in \mathbb{N} \setminus edgeG$

$grd5 : \forall v \cdot v \in vertL1 \Rightarrow tL1_V(v) = tG_V(mV(v))$

$grd6 : \forall e \cdot e \in edgeL1 \Rightarrow tL1_E(e) = tG_E(mE(e))$

$grd7 : \forall e \cdot e \in edgeL1 \Rightarrow mV(sourceL1(e)) = sourceG(mE(e)) \wedge$
 $mV(targetL1(e)) = targetG(mE(e))$

then

$act3 : vertG := vertG \cup \{newV\}$

$actE : edgeG := (edgeG \setminus \{mE(e1_L1)\}) \cup \{newE\}$

$acts : sourceG := (\{mE(e1_L1)\} \triangleleft sourceG) \cup \{newE \mapsto (sourceG(mE(e1_L1)))\}$

$actt : targetG := (\{mE(e1_L1)\} \triangleleft targetG) \cup \{newE \mapsto newV\}$

$acttV : tG_V := tG_V \cup \{newV \mapsto Vertex2\}$

$act6 : tG_E := (\{mE(e1_L1)\} \triangleleft tG_E) \cup \{newE \mapsto Edge2\}$

end

m is a match

vertex type compatibility
edge type compatibility
source/target compatibility

Properties

- * Properties are stated as invariants:

INVARIANTS

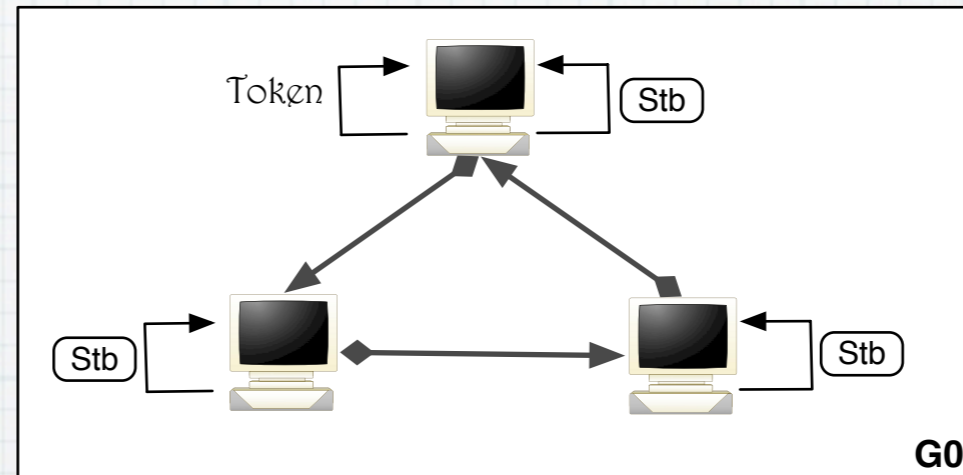
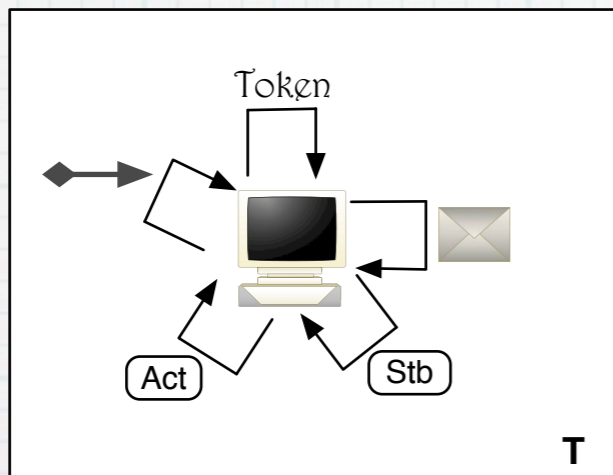
... : ...

prop1 : *finite(edgeG)*

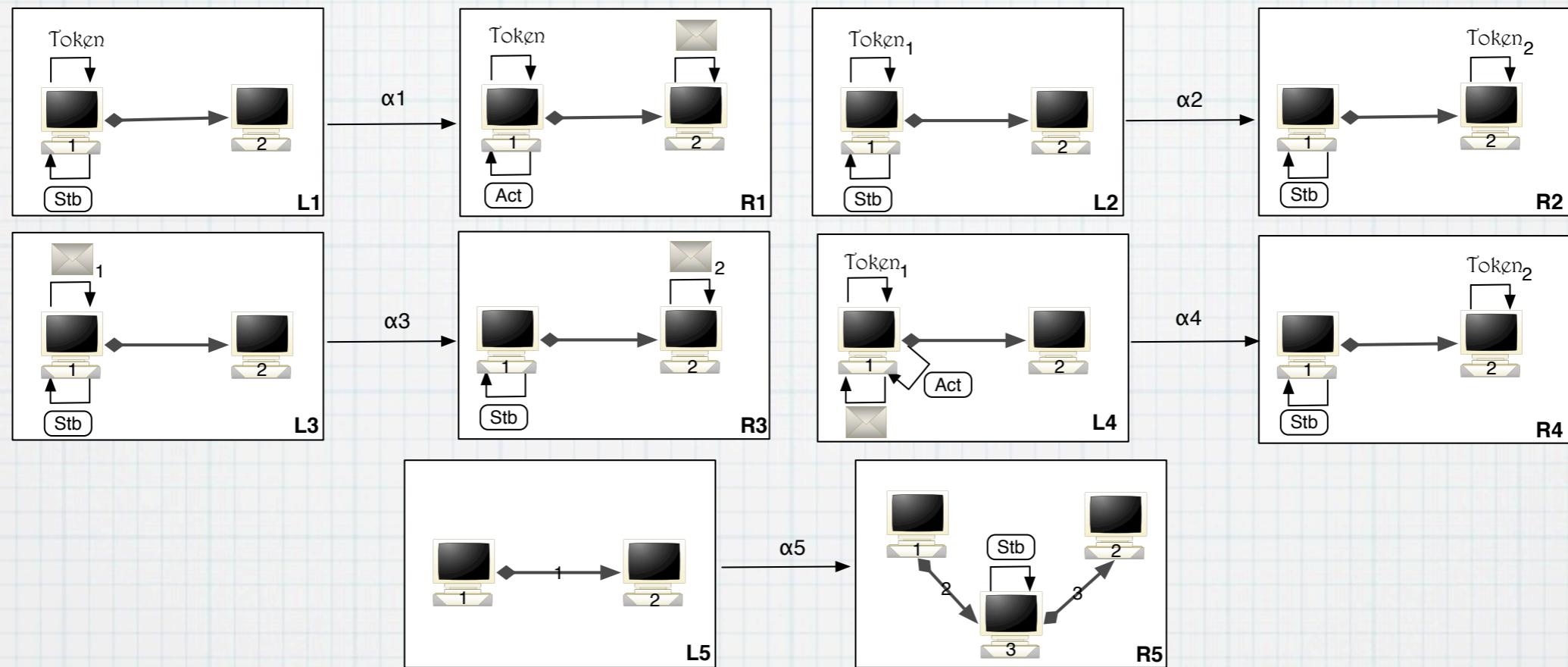
prop2 : *card(edgeG) ≤ 2*

- * In Event-B, proof obligations are generated to guarantee that all **invariants are well-defined**, **are valid in any initial state** and **are preserved by all events**

Another Example: Token Ring



Another Example: Token Ring



Rodin Platform File Edit Navigate Search Project Theory Run ProB Refactor Event-B Window Help

Proving - TR_RingPropertyAllRulesTeseSimone/mch_trAll.bum - Rodin Platform - /Users/scosta/Documents/Doutorado/Rodin/GG_ICGMT

Proof Tree Event-B Explorer

- Events
 - Proof Obligations
 - prop1/WD
 - INITIALISATION/inv_vertG/INV
 - INITIALISATION/inv_edgeG/INV
 - INITIALISATION/inv_srcGtype/INV
 - INITIALISATION/inv_tgtGtype/INV
 - INITIALISATION/inv_tG_V/INV
 - INITIALISATION/inv_tG_E/INV
 - INITIALISATION/prop1fin/INV
 - INITIALISATION/prop1/INV**
 - rule1/grd_vertices/WD
 - rule1/grd_edges/WD
 - rule1/grd_srctgt/WD
 - rule1/inv_edgeG/INV
 - rule1/inv_srcGtype/INV
 - rule1/inv_tgtGtype/INV
 - rule1/inv_tG_E/INV
 - rule1/prop1fin/INV
 - rule1/prop1/INV**
 - rule1/act_E/WD
 - rule1/act_src/WD
 - rule1/act_tgt/WD
 - rule1/act_tE/WD
 - rule2/grd_vertices/WD
 - rule2/grd_edges/WD
 - rule2/grd_srctgt/WD
 - rule2/inv_edgeG/INV
 - rule2/inv_srcGtype/INV
 - rule2/inv_tgtGtype/INV
 - rule2/inv_tG_E/INV
 - rule2/prop1fin/INV

1 items selected

ctx_trAll mch_trAll mch_trAll

INVARIANTS

```

inv_vertG : vertG ∈ P(N) // // (Graph) Vertices are natural numbers.
inv_edgeG : edgeG ∈ P(N) // // (Graph) Edges are natural numbers.
inv_srcGtype : sourceG ∈ edgeG → vertG // // (Graph) function sourceG
inv_tgtGtype : targetG ∈ edgeG → vertG // // (Graph) function targetG
inv_tG_V : tG_V ∈ vertG → vertT // // (Graph) function tG_V
inv_tG_E : tG_E ∈ edgeG → edgeT // // (Graph) function tG_E
prop1fin : finite(dom(tG_E ▷ {Tok})) // // Property 0: The set of edges of
prop1 : card(dom(tG_E ▷ {Tok})) = 1 // // Property 1: Any reachable graph t

```

EVENTS

INITIALISATION ≡

STATUS
ordinary

Pretty Print Edit Synthesis Dependencies

Goal

```

ct card(dom(((mE(Stb11)) ◀ tG_E)u(newEact ↦ Act,newEmsg ↦ Msg)) ▷ {Tok}))=1

```

Proof Control Statistics Rodin Problems

86 POs: 57 (A) + 29 (SA)

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Final Remarks

- ☑ GGs can be faithfully encoded in Event-B
- ☑ Rodin offers a set of theorem provers that can be used to verify graph grammars
- ☐ However, we need
 - * libraries for commonly used operations;
 - * theories for graph structures;
 - * property patterns (tactics patterns);
 - * integration of data types (attributed GTS), NACs, ...
 - * more “intelligent” theorem prover: Isabelle???

Thank you!!!

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