

➔ Towards PaaS and Clouds Our experience with BonjourGrid and PastryGrid – AOC Team –

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Bi-lateral China-France workshop



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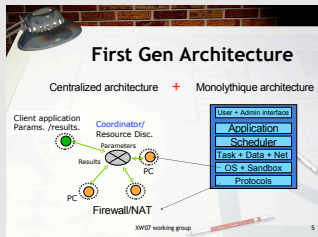
➔ Objectives

1. Motivate research projects in Grids & Clouds ;
2. Starting from recent advances in Desktop Grid Middleware:
 - ➔ BonjourGrid (orchestration of multiple instances of DG middleware) and PastryGrid (fully distributed execution of applications)
 - ➔ Joint works with UTIC lab., Tunisia (Heithem Abbes and Mohamed Jemni)
3. Before keeping innovative ideas to reuse in Cloud Architectures / Systems:
 - ➔ decentralized architectures and services;
 - ➔ large scale systems (FT);
 - ➔ interoperability of services (the client is not a prisoner, or if it is, he can choose his prison(s)!)



Desktop Grid Architectures

Desktop Grid

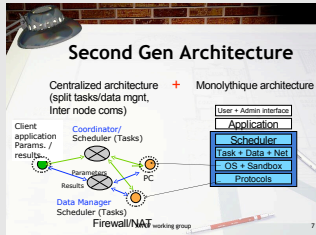


Key Points

- ⊕ Federation of thousand of nodes;
- ⊕ Internet as the communication layer: no trust!
- ⊕ Volatility; local IP; Firewall

⊕ Desktop Grid Architectures

Desktop Grid



Future Generation (in 2006)

- ⊕ **Distributed Architecture**
- ⊕ Architecture with modularity: every component is “configurable”: scheduler, storage, transport protocole
- ⊕ Direct communications between peers;
- ⊕ Security;
- ⊕ Applications coming from any sciences (e-Science applications)



➔ In search of distributed architecture

First line: publish/subscribe system to notify and coordinate services and multiple DG without a central broker ⇒ BonjourGrid;

Second line: approach based on structured overlay network to discover (on the fly) the next node executing the next task ⇒ PastryGrid;

(main contributions of Heithem Abbes in his PhD)



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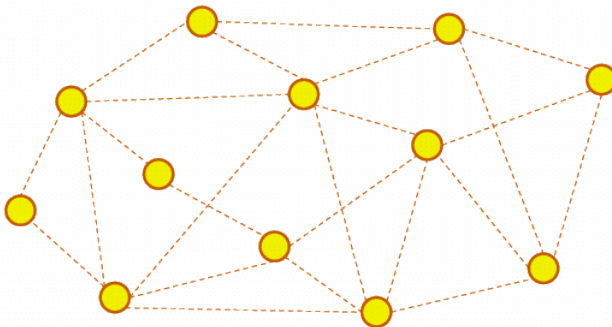


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- ⊕ Each coordinator searches, in a concurrent way, participants (idle machines)

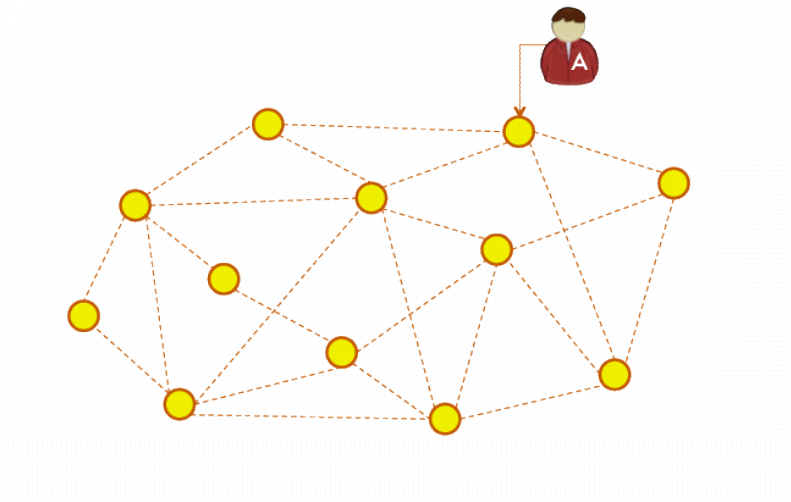


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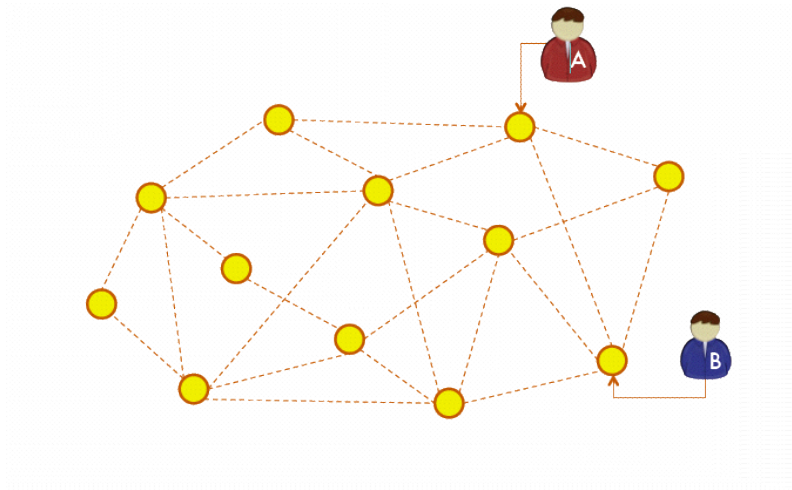


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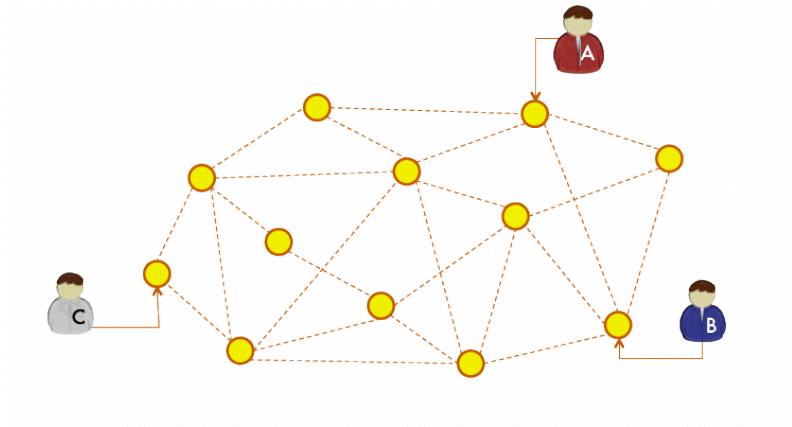


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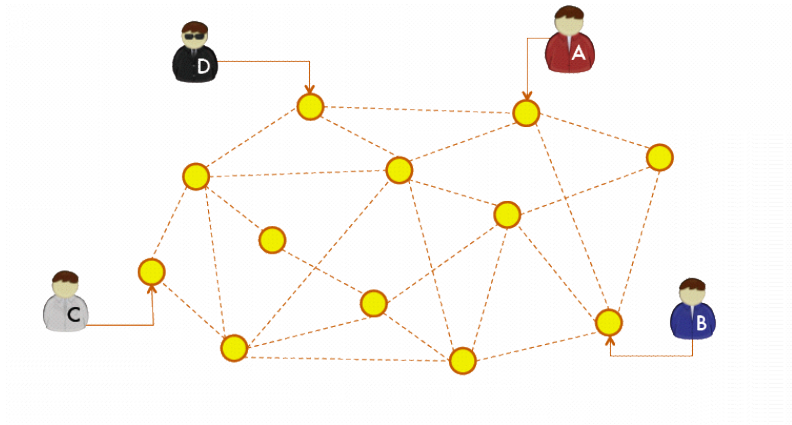


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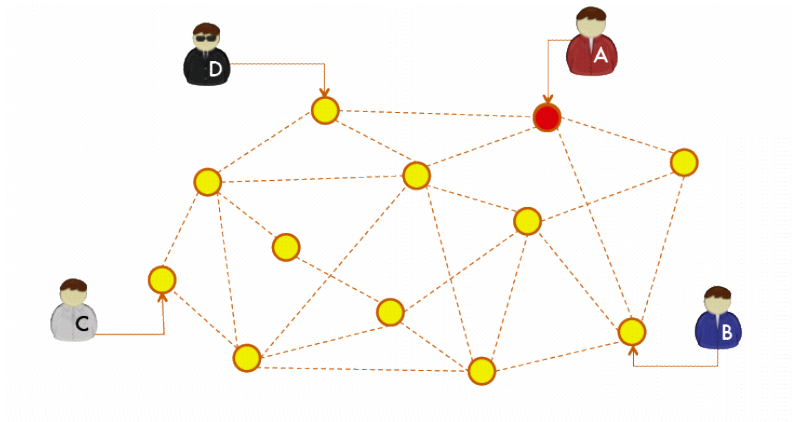


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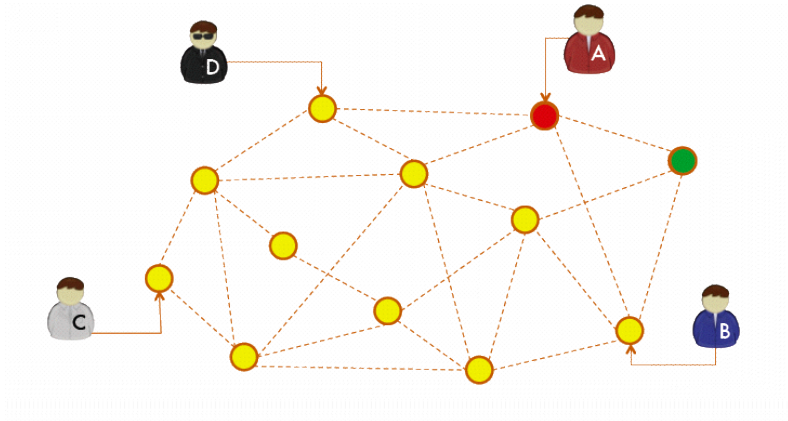


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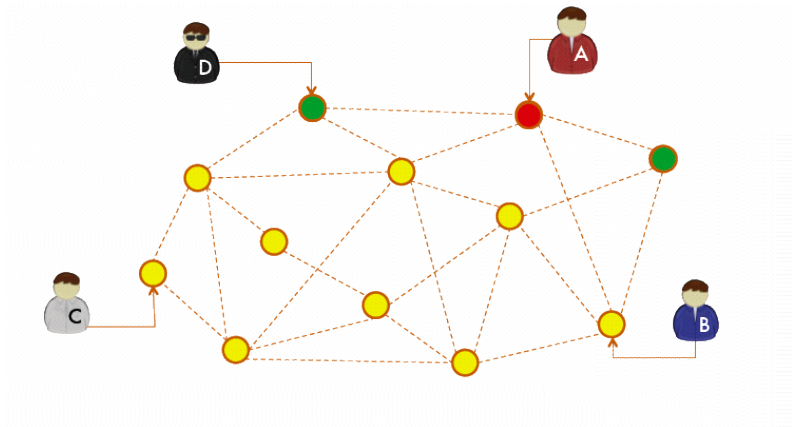


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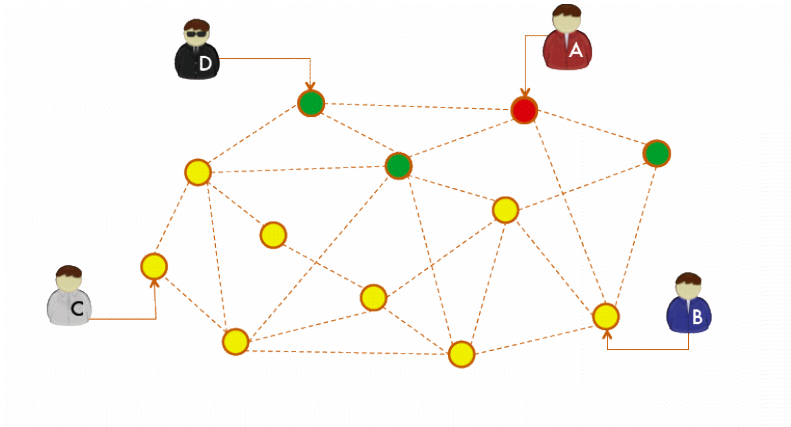


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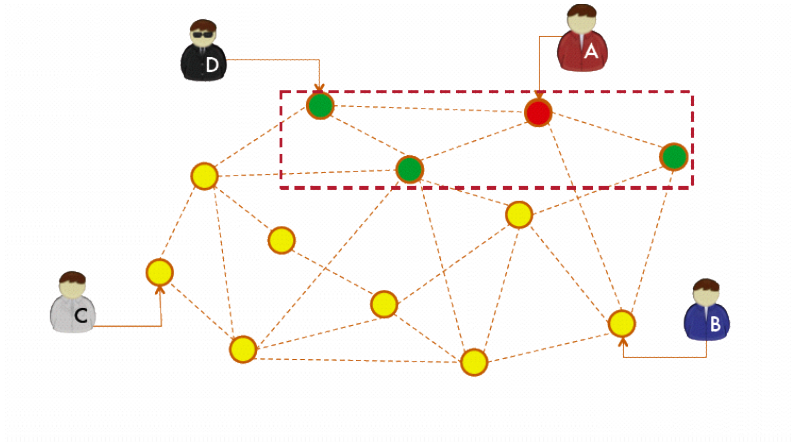


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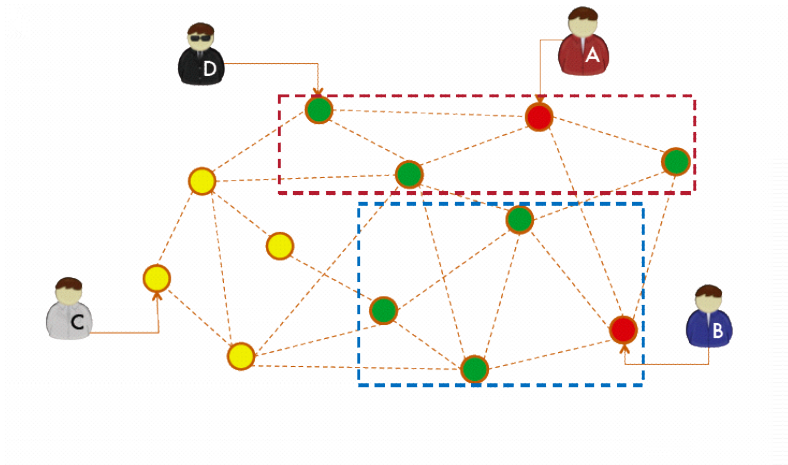


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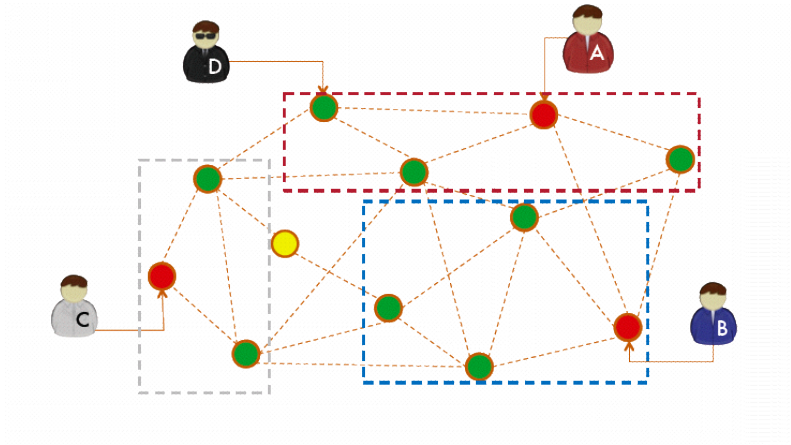


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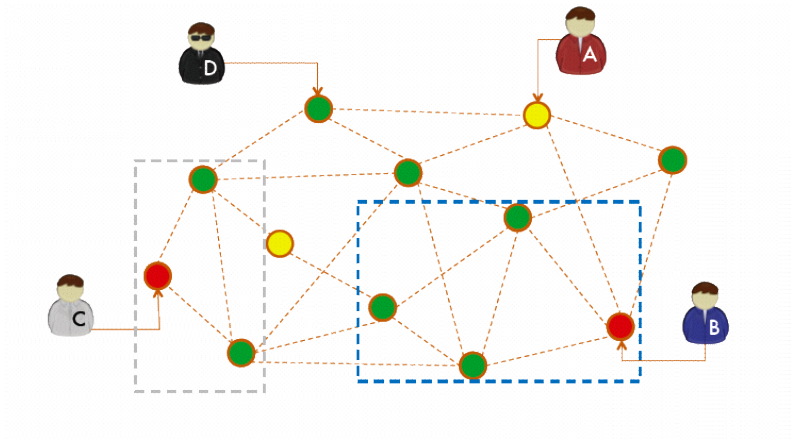


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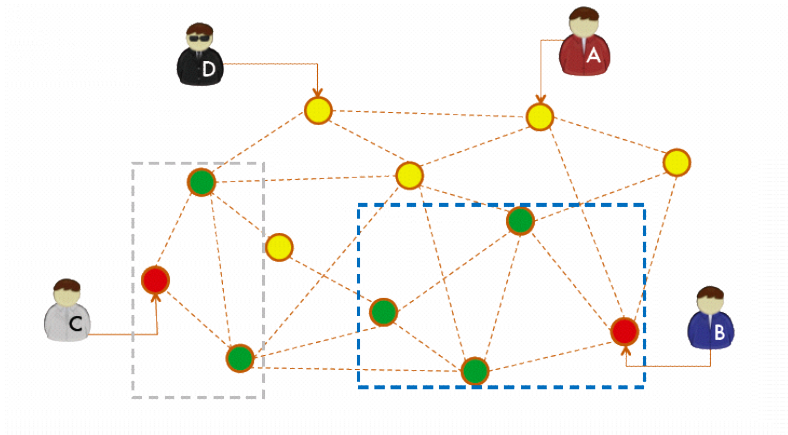


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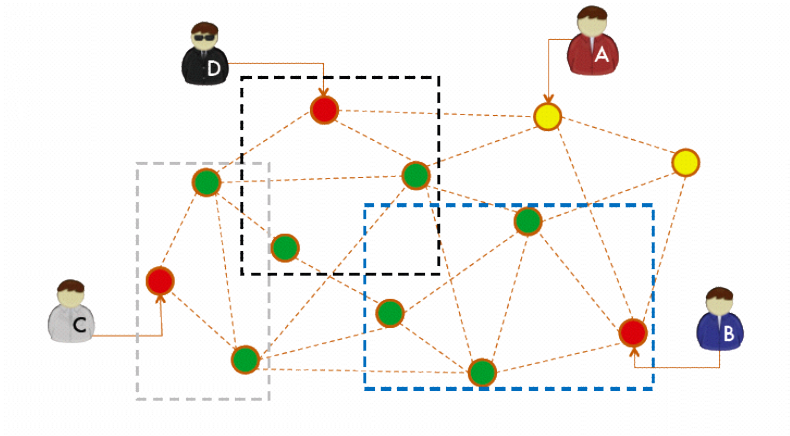


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- ⊕ The current protocol has been developed/specified with 'ad-hoc' methods → we need to consolidate the trust (ongoing project to verify it, based on Colored Petri Nets)



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- ⊕ Our solution: tolerate the failure of coordinators
 - ⊕ For any application we create and manage dynamically copies of the coordinator;
 - ⊕ We manage k copies; based on passive replication.
 - ⊕ When a service disappears: we added a special status flag to distinguish between 'end of the application' / 'failure' \Rightarrow slaves can redirect the communication to a copy.



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 - ⊕ in terms of robustness in supporting FT;

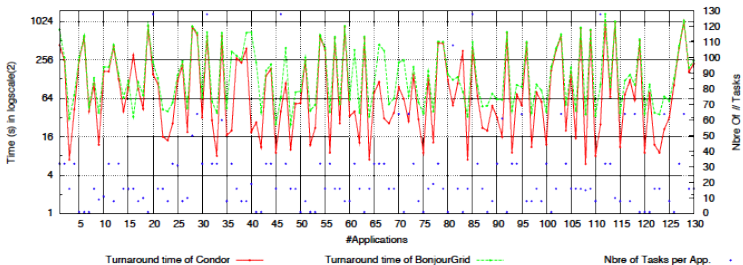


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 - ⊕ in terms of robustness in supporting FT;
- ⊕ Example Condor: 130 applications (2 to 128 // tasks), 200 nodes, application task: 1s to 500s. Result: with BonjourGrid, 35% of applications generate a delay of about 30s.



⊕ Experiments: one example





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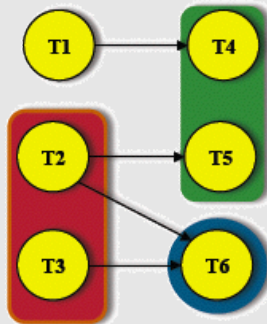


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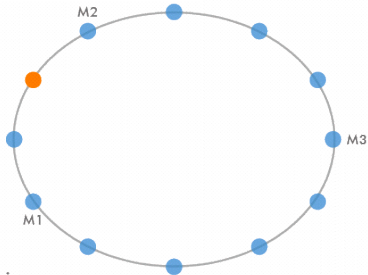
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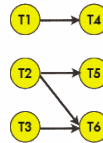
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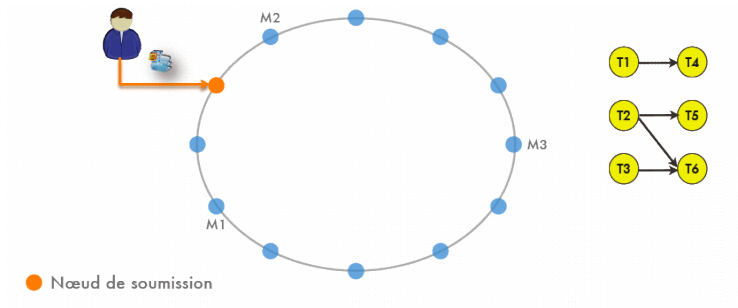
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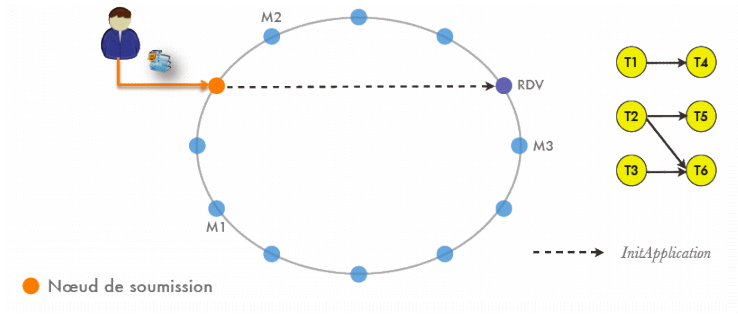
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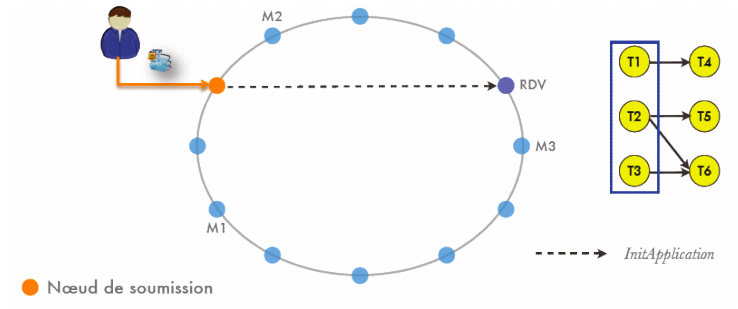
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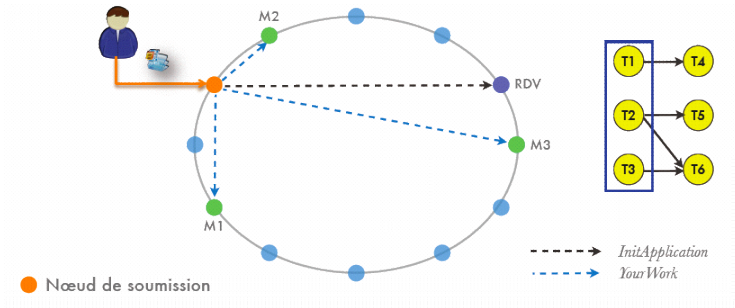
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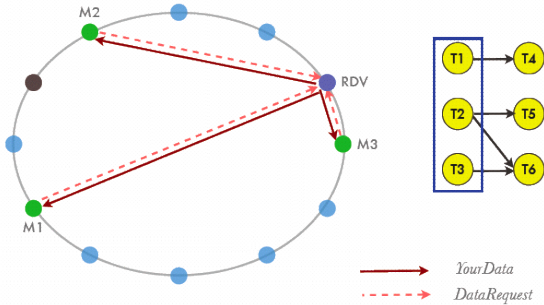


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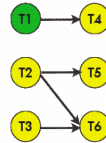
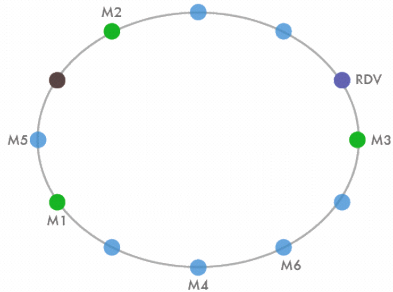


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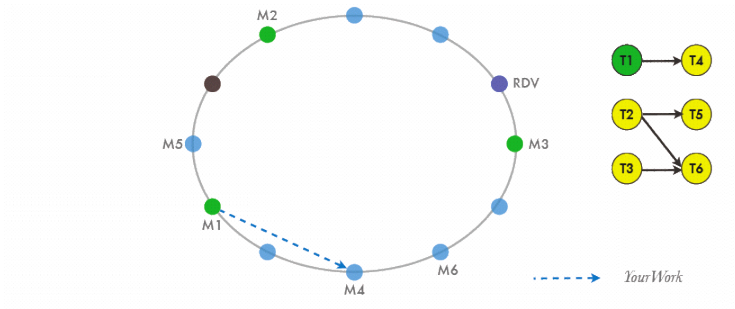




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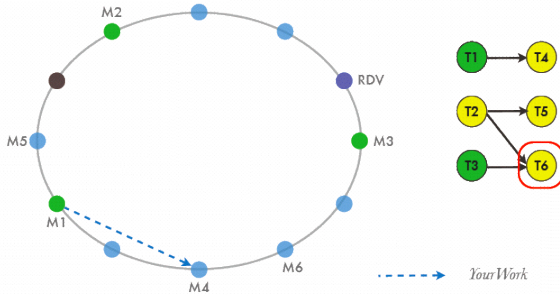


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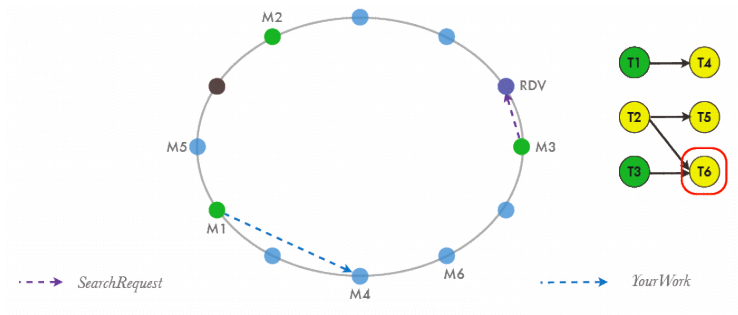


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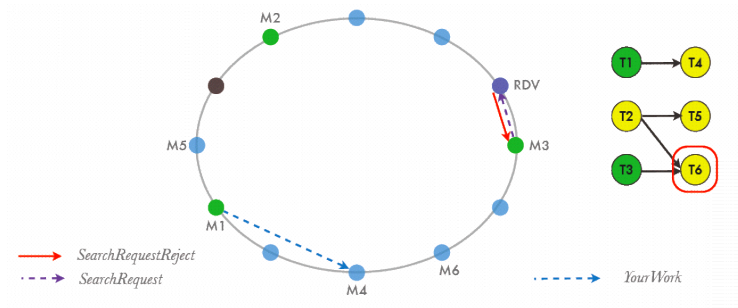




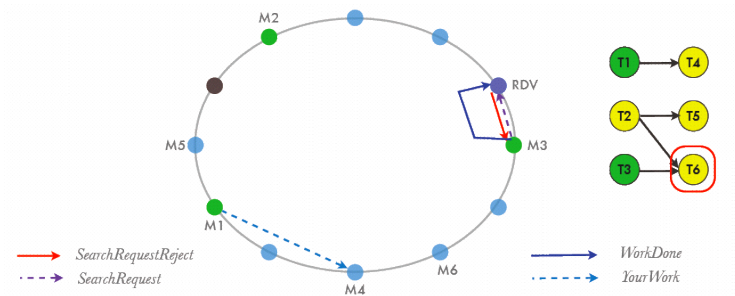
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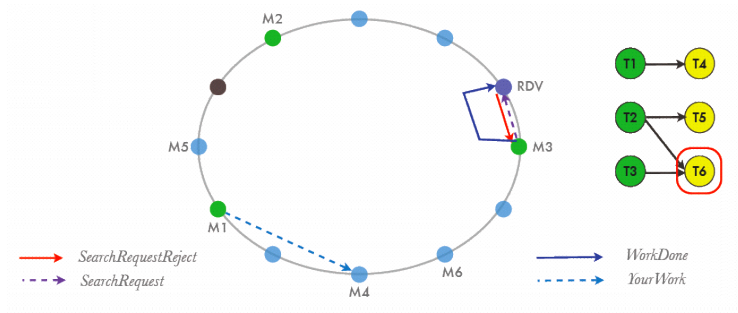
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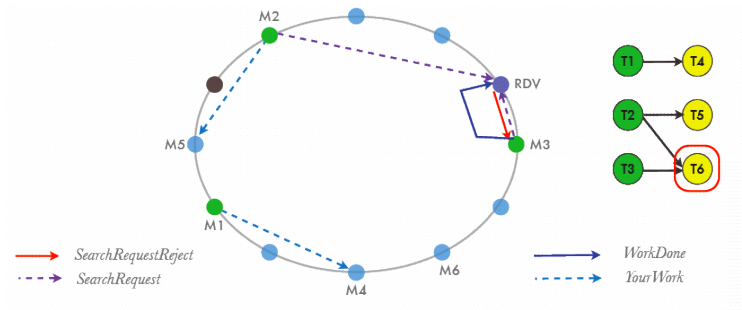


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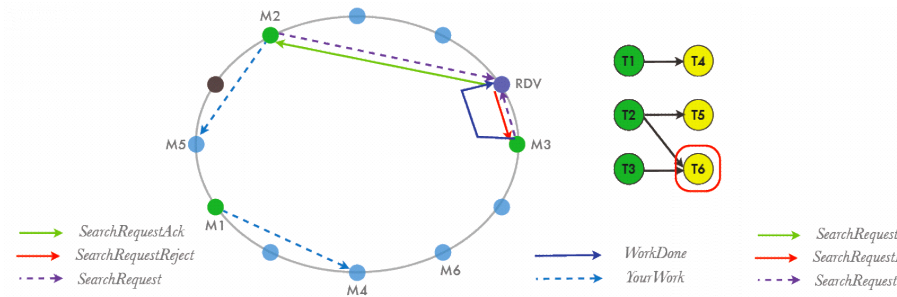


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- ⊕ Main observations:
 - ⊕ In all cases, PastryGrid terminates;
 - ⊕ The recovery time depends on the node type;
 - ⊕ The delay varies from 4:53s to 7:16:41s... but it works! The number of delayed applications varies from 44 to 98.



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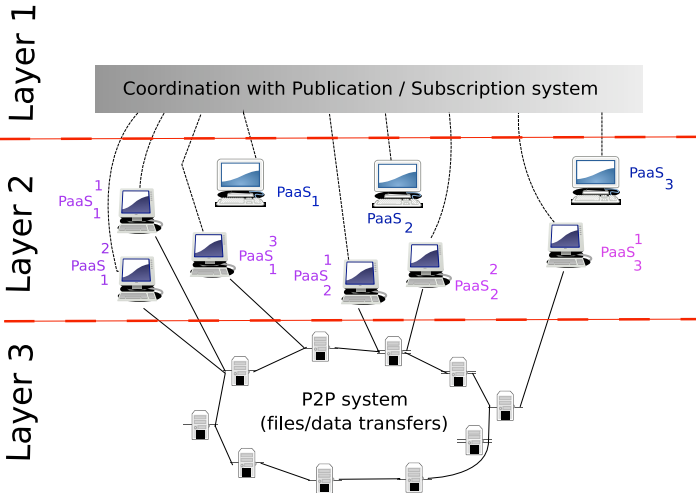
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- ➔ We have started an initiative for defining and designing a “general purpose PaaS” based on distributed protocols for coordination and data exchange.



Architecture overview of the PaaS Coordinated project





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➔ Some Challenges

- ➔ Where to insert the different connectors in the PaaS software stack to get an open infrastructure?

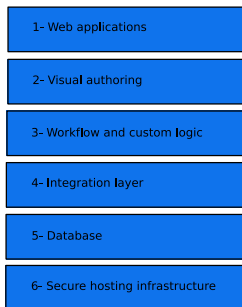


Figure: Software stack in PaaS (source: Coghead)



➔ Research opportunities

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- ➔ OSOE course: <http://www.osoe-project.org/lesson>
- ➔ TioLive tutorial:
<https://www.tiolive.com/documentation/tiolive-tutorial>
- ➔ Documentation for developers:
https://www.myerp5.com/kb/documentation_section/developer/
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- ➔ Main change: accept to manage redundancy, scaling the server (even for temporary needs), synchronisation ⇒ coordination with grid technology (BonjourGrid, PastryGrid?);



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- ➔ Ex: a company wants to install the Virtual Desktop EyeOS and the TioLive/ERP5 PaaS. During the night, the company rents different services:
 - ➔ one (company) to many – many (services) to many (companies) = new abilities, new business!
 - ➔ demonstrate that a single coordination protocol is better than configuring as many middlewares than we have software!



➔ **Towards PaaS and Clouds**
Our experience with
BonjourGrid and PastryGrid
– AOC Team –

Christophe Cérin¹

¹Université de Paris XIII, CNRS UMR 7030, France

Bi-lateral China-France workshop