

Size of maximum b -matchings in sparse random graphs

Mathieu Leconte

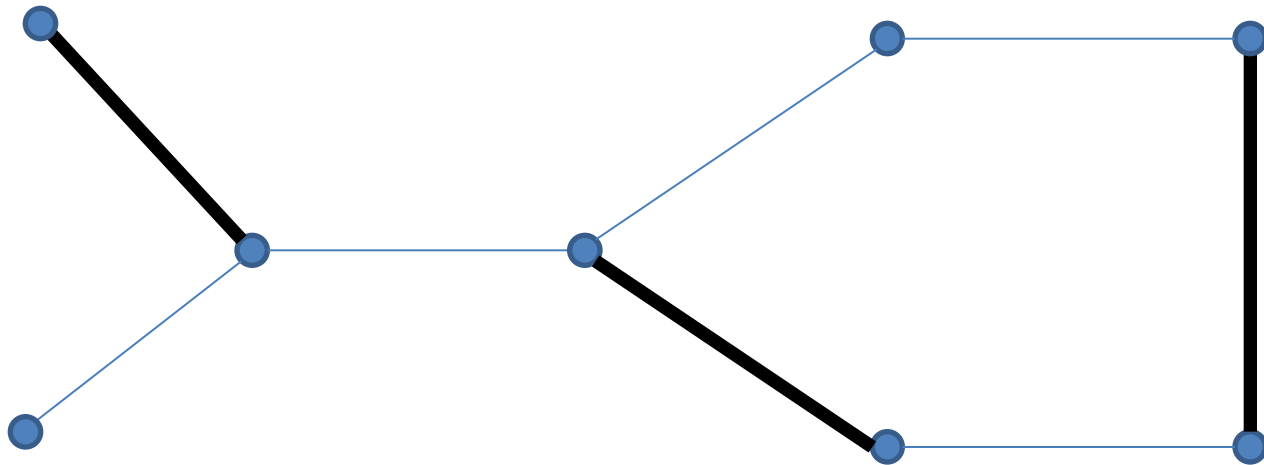
Joint work with Marc Lelarge and
Laurent Massoulié

Goal?

- Computing the asymptotic density of a b -matching as the size of graphs tend to infinity
 - Can we compute other quantities in a similar way?
- Sequences of sparse random graphs with a known local structure
 - Graphs that converge (in the local weak sense) towards **Galton-Watson trees**
 - Ex: Erdos-Renyi, configuration model

Matchings

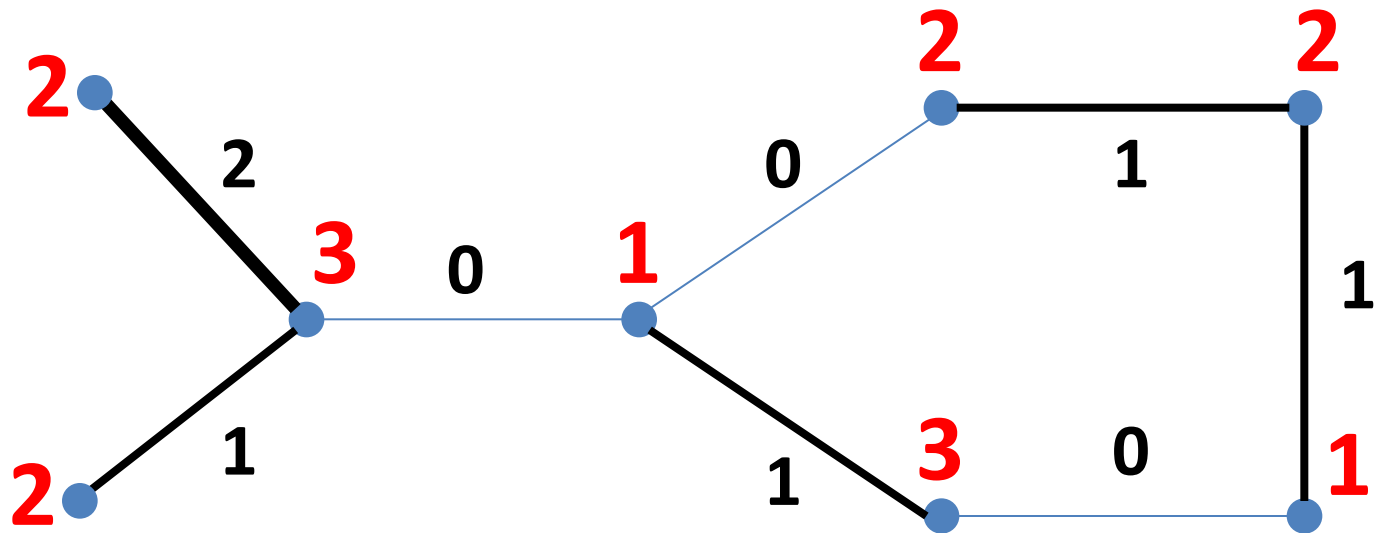
- Graph $G=(V,E)$
- Matching= subset of edges E' such that each vertex is adjacent to at most one element in E'



studied by Zdeborova and Mézard (2006)
Bordenave, Lelarge, Salez (2011)

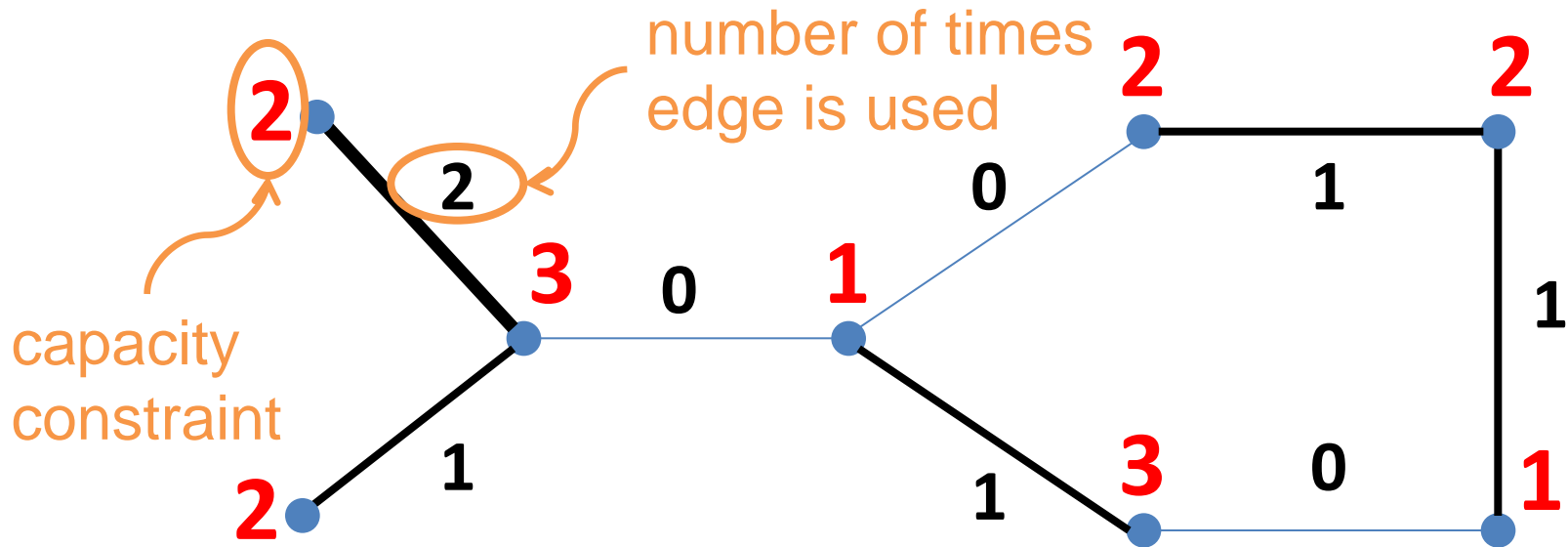
b-matchings

- Capacity constraint b_v at vertex v
- Each edge may be used more than once
- Total usage of edges adjacent to v must not exceed b_v

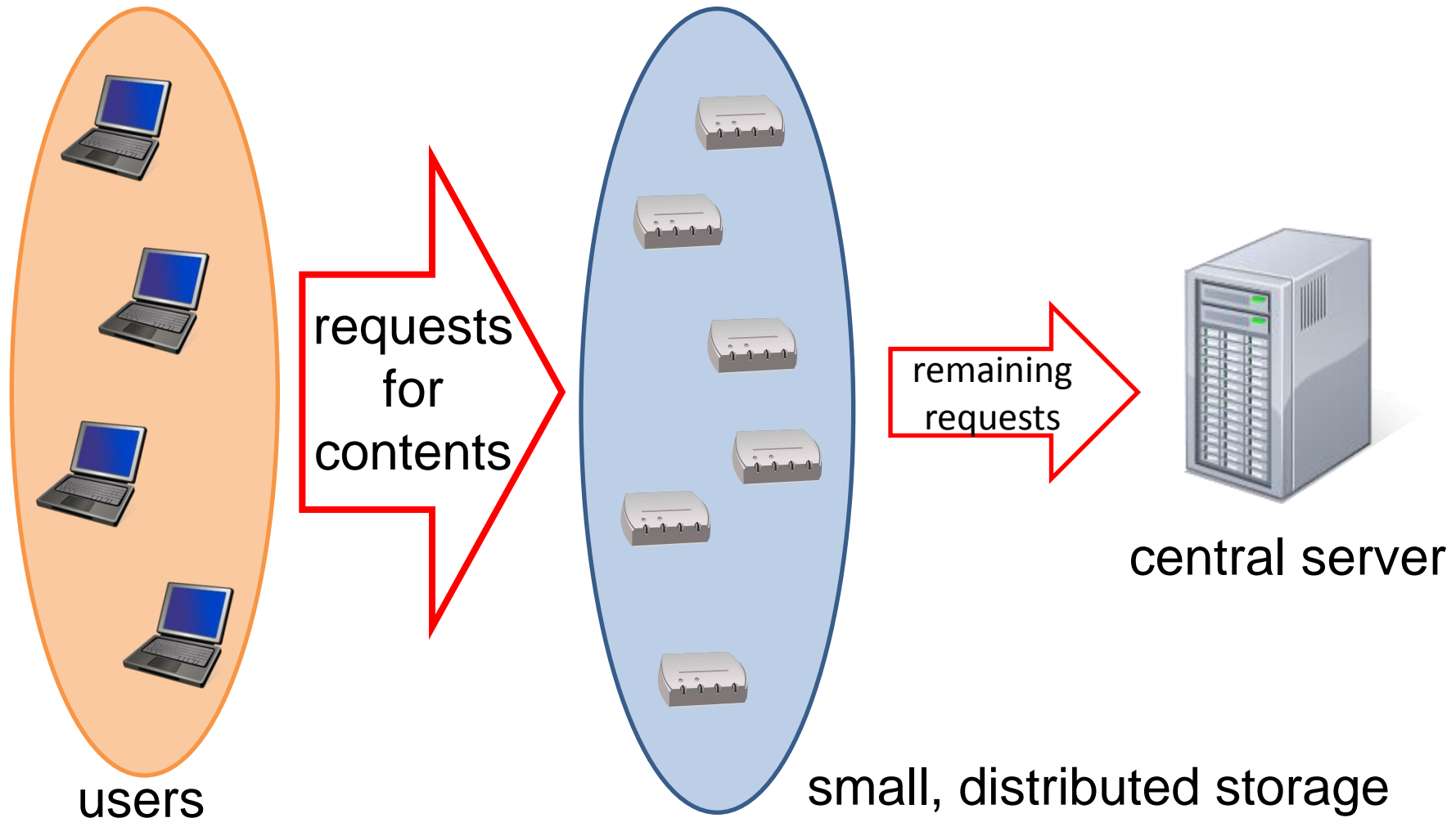


b-matchings

- Capacity constraint b_v at vertex v
- Each edge may be used more than once
- Total usage of edges adjacent to v must not exceed b_v

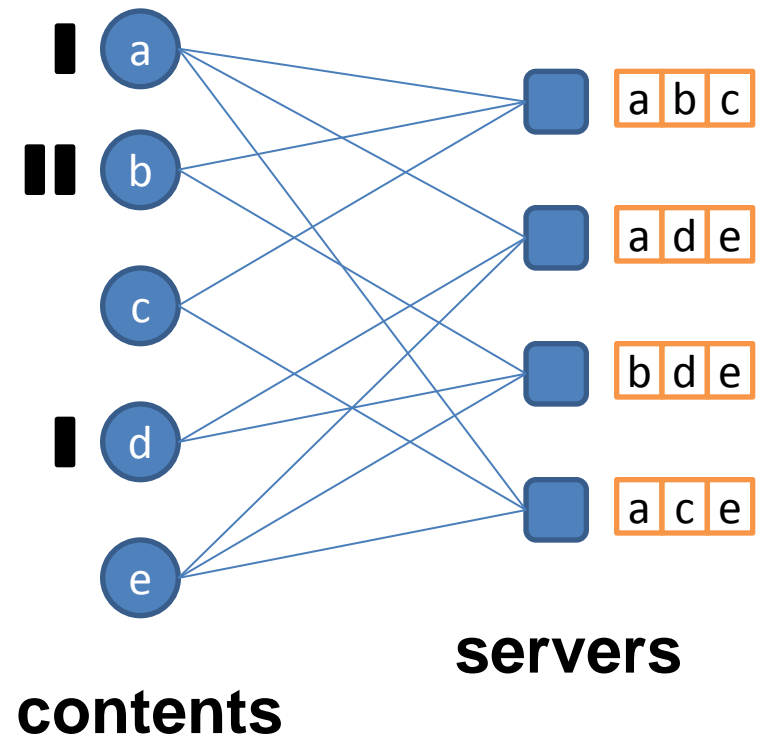


An application: distributed content distribution system



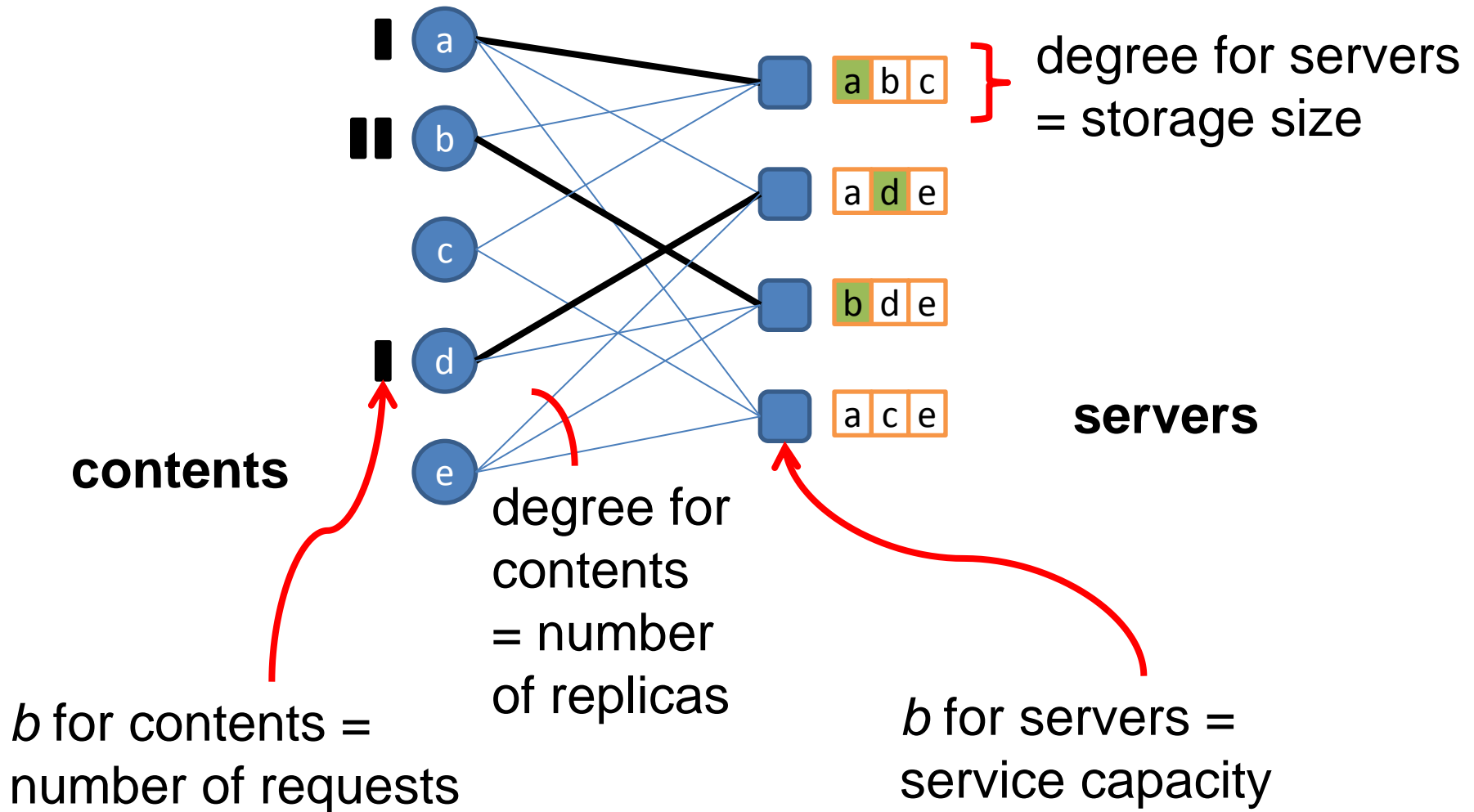
Distributed content distribution system

- n contents, m servers
- Contents stored by each server determined **independently** (but **not uniformly**)
- Number of requests and number of replicas for each content **jointly determined**
- $n \rightarrow \infty, n/m \rightarrow \beta$

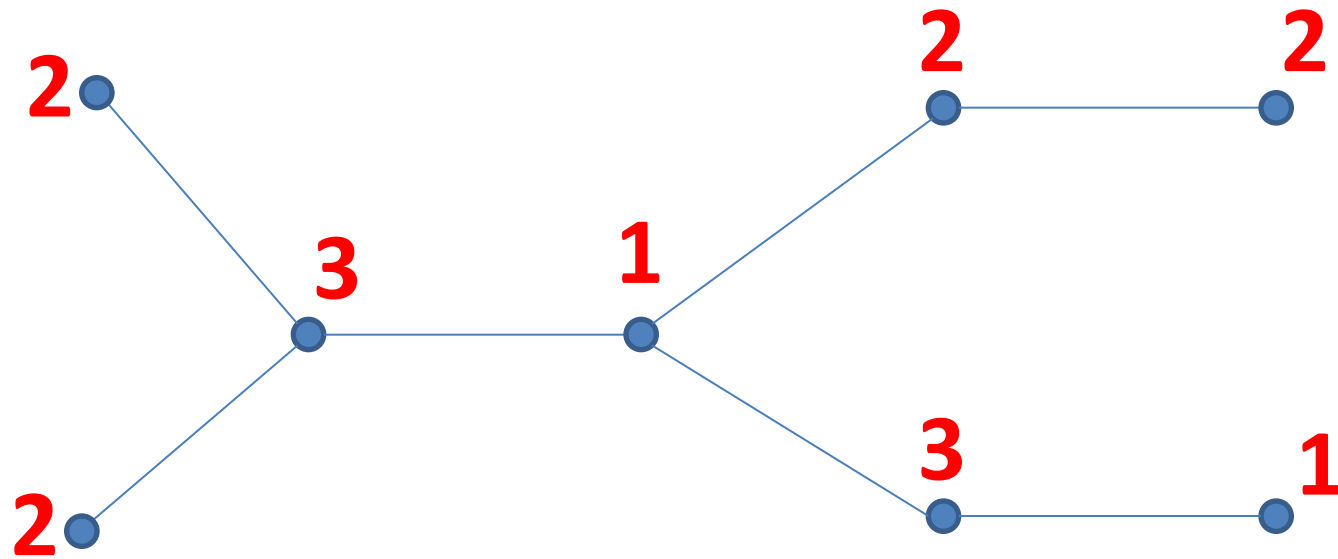


What fraction of the requests can be served?

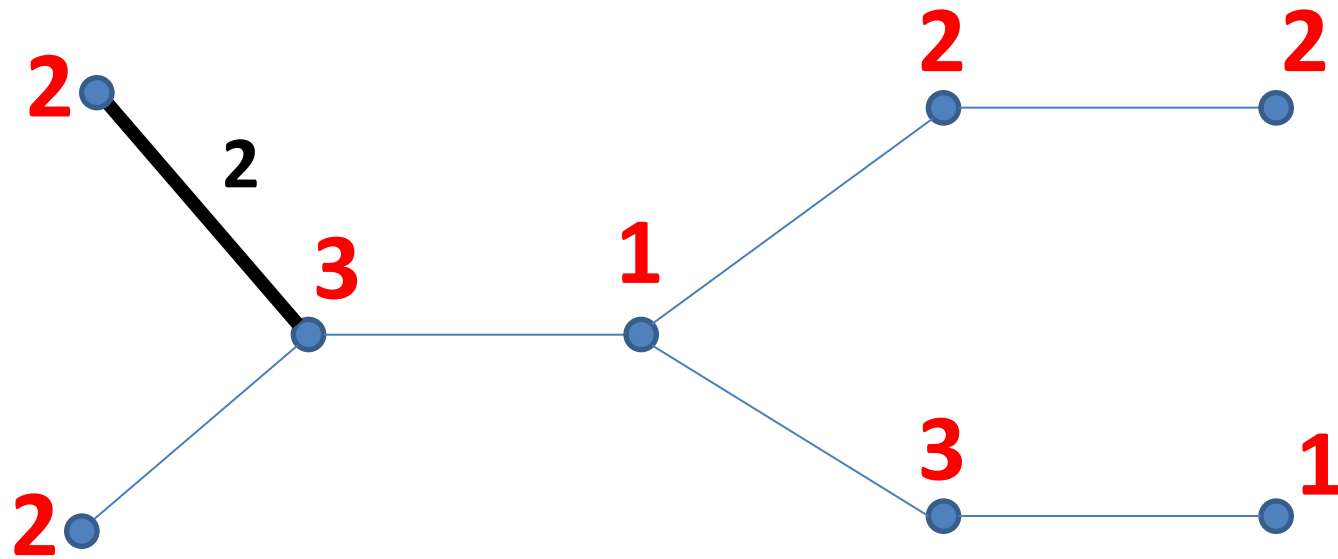
Distributed content distribution system



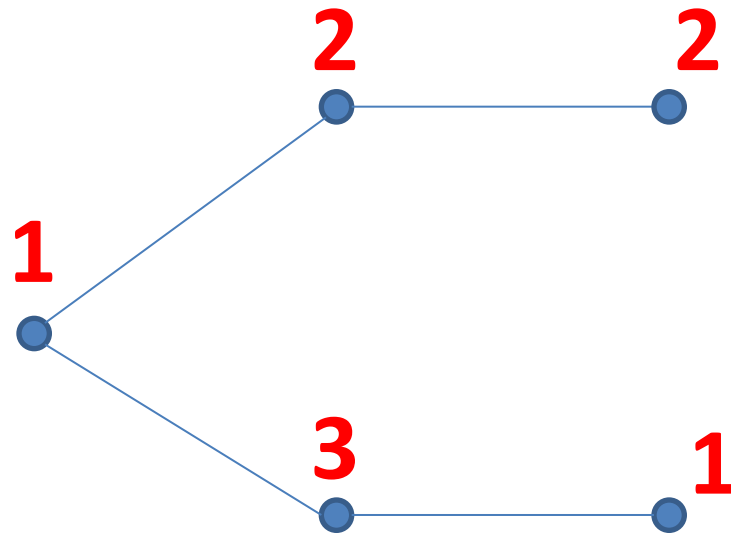
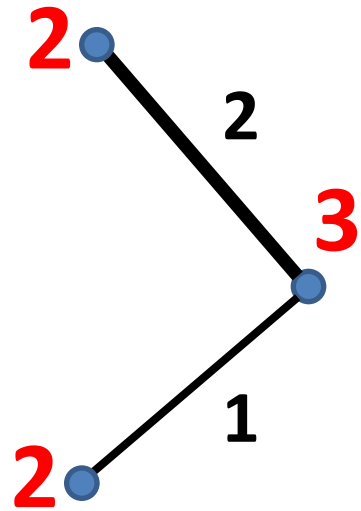
Greedy algorithm on trees for finding a maximum b-matching



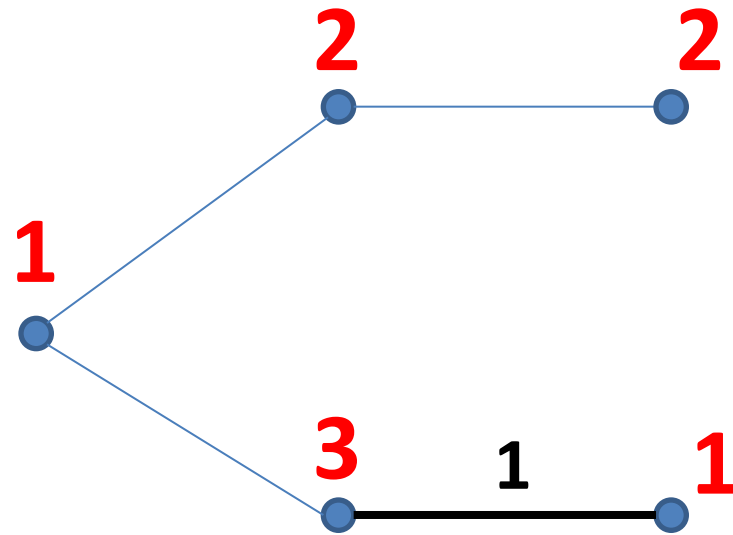
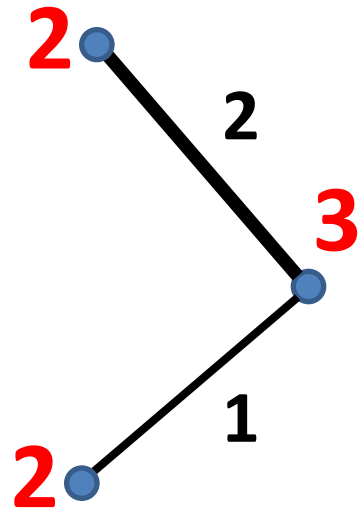
Greedy algorithm on trees



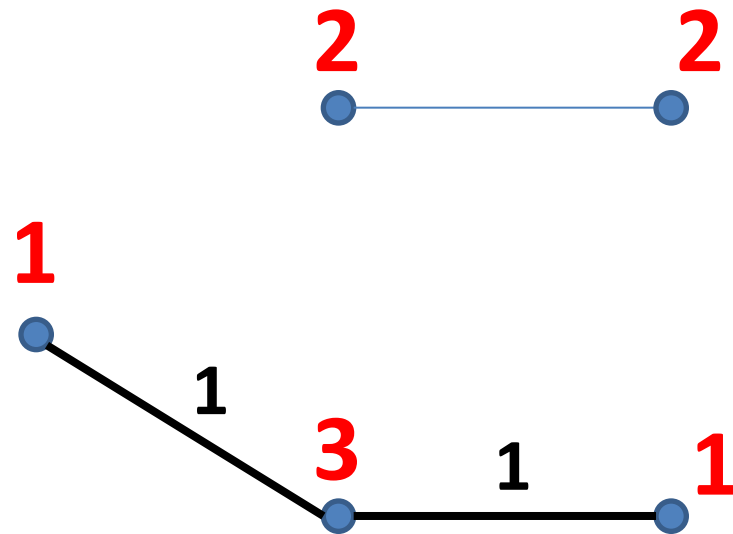
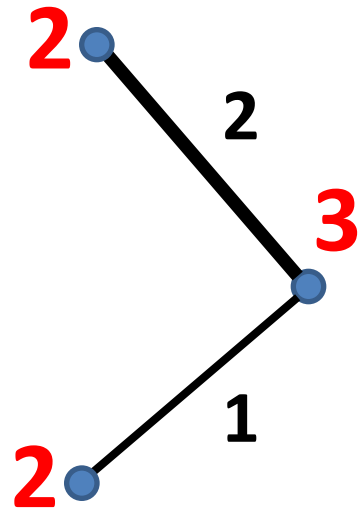
Greedy algorithm on trees



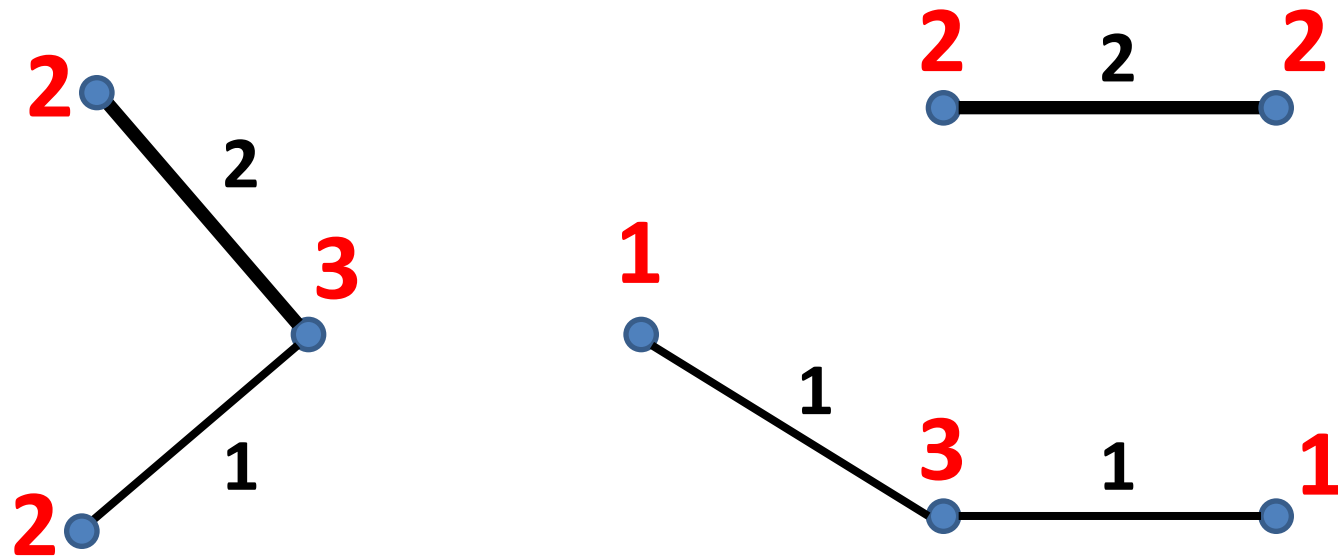
Greedy algorithm on trees



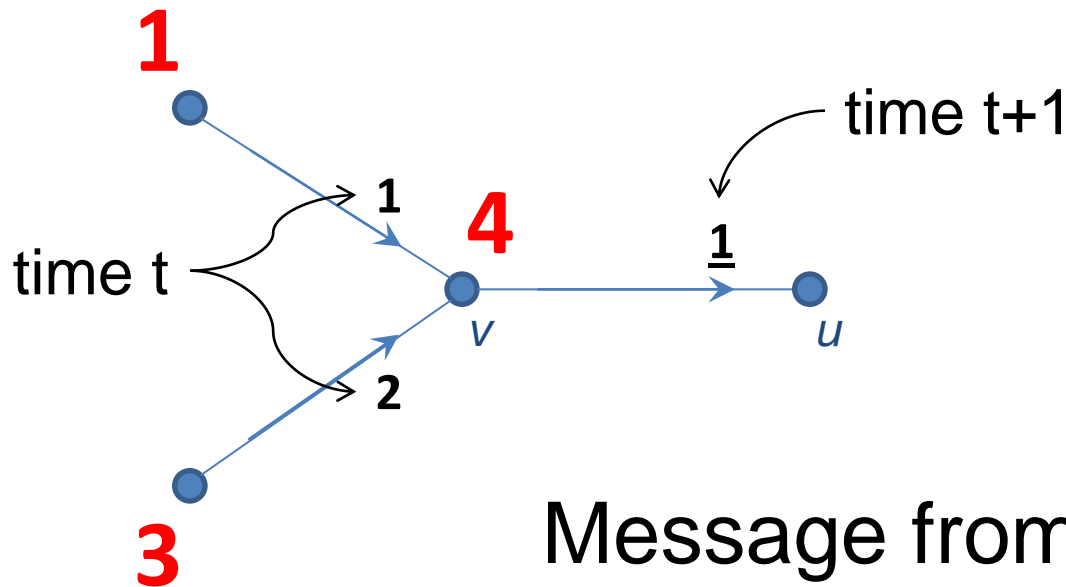
Greedy algorithm on trees



Greedy algorithm on trees



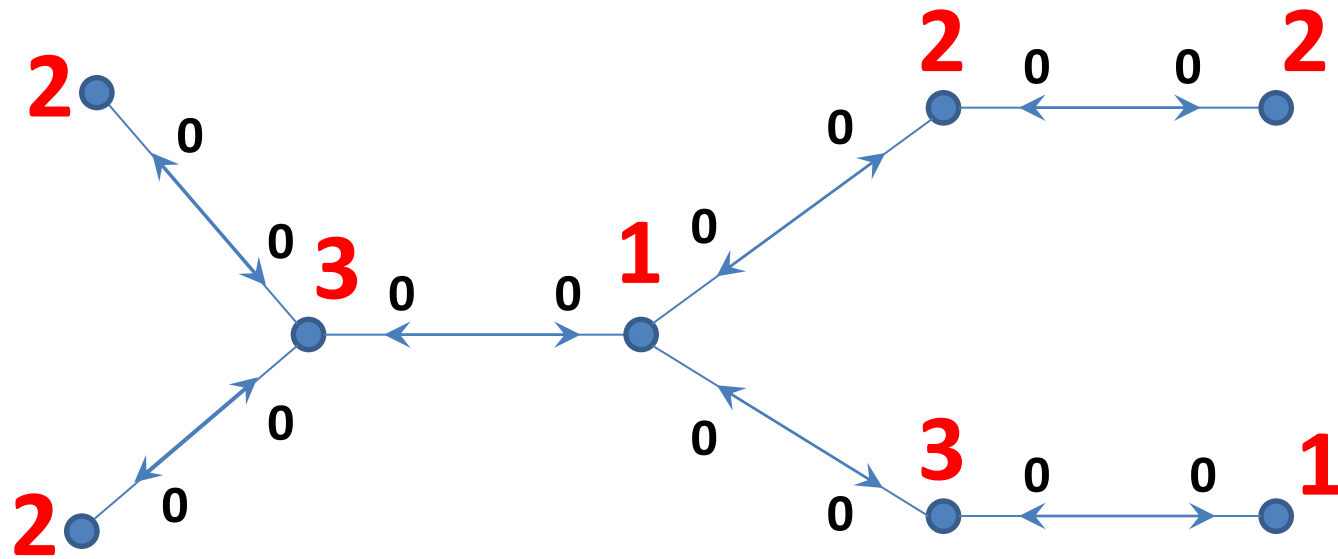
Message-passing version: BP at 0 temperature



Message from v to u :

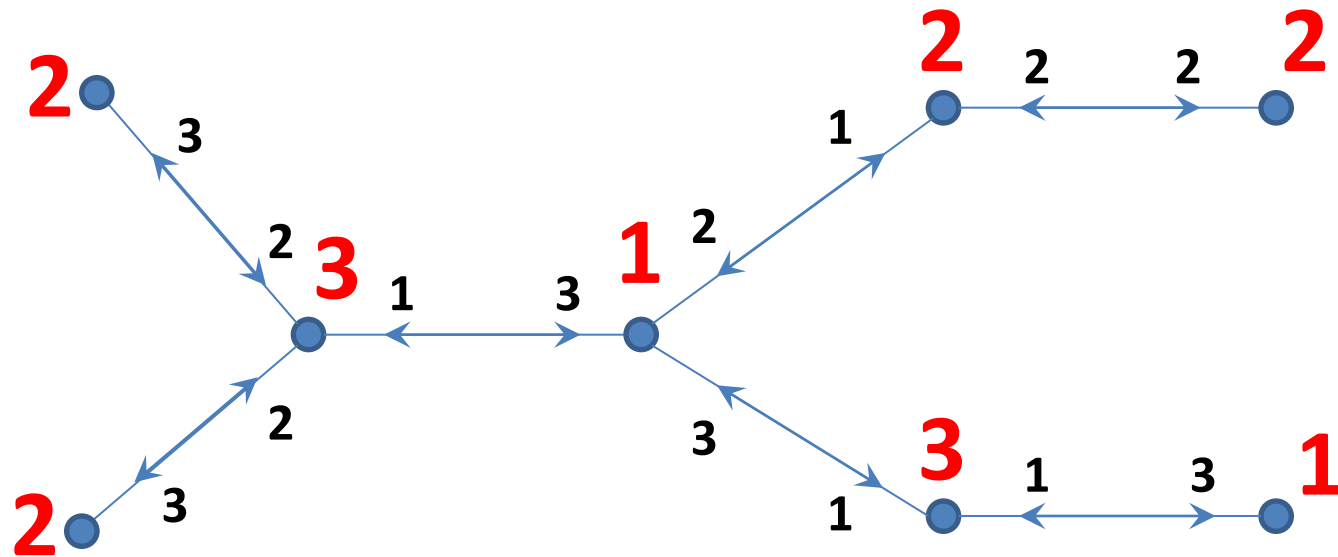
$$I_{v \rightarrow u}^{t+1} = \left(b_v - \sum_{w \in \partial v \setminus u} I_{w \rightarrow v}^t \right)^+ = P_{v \rightarrow u} \left(\mathbf{I}^t \right)$$

BP at 0 temperature: initial messages



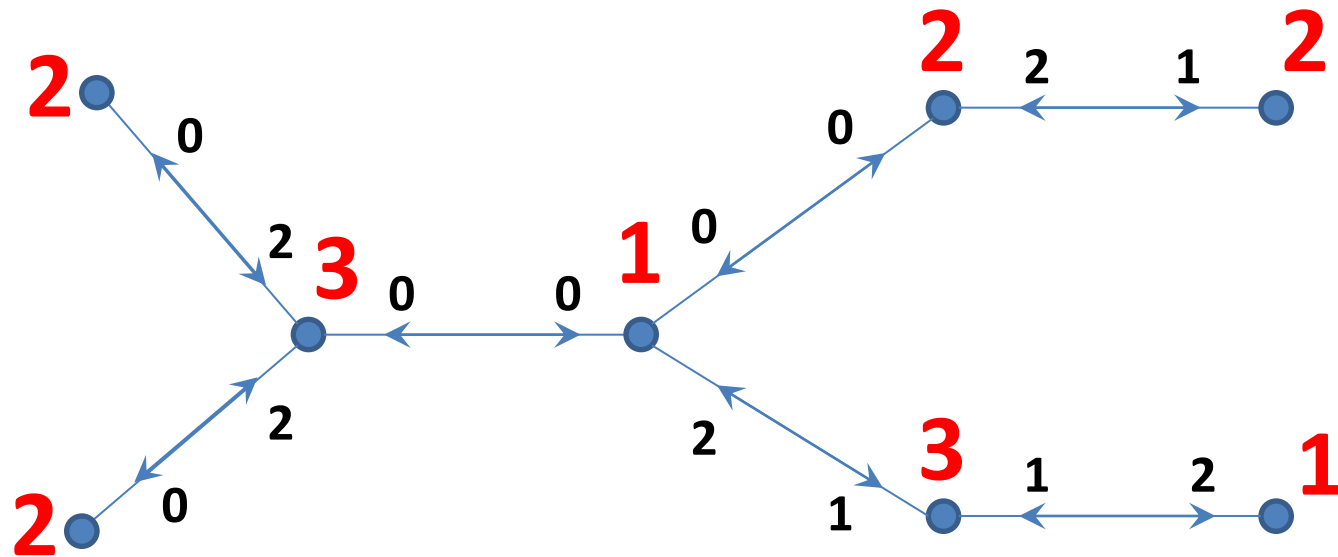
$$I_0=0$$

BP at 0 temperature: after 1 iteration



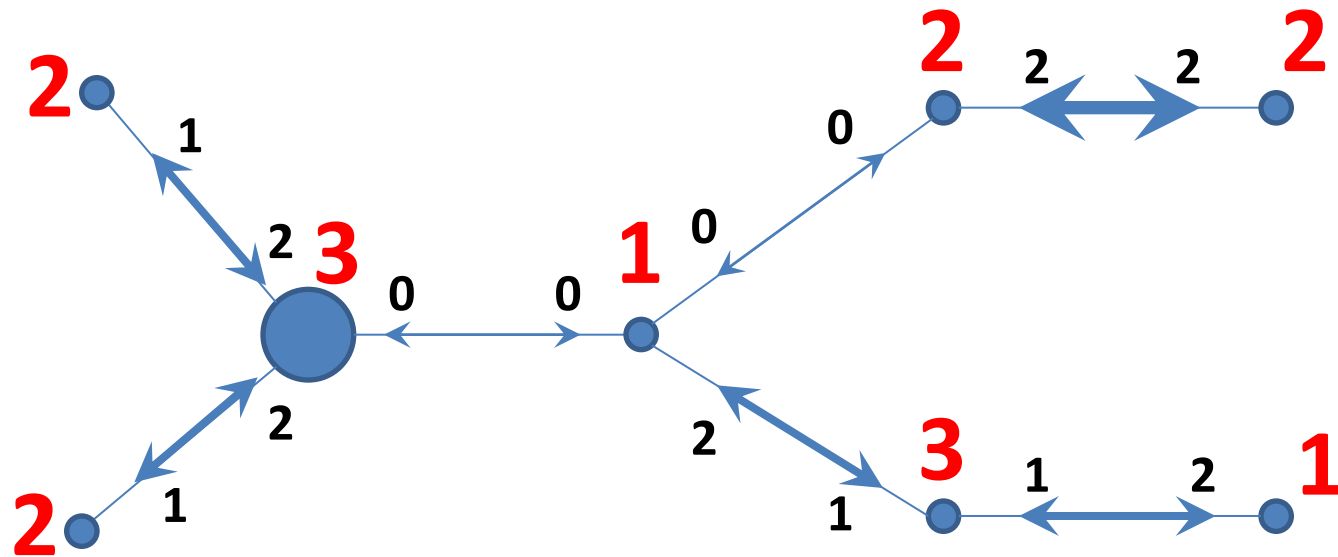
$$I_1 = P_G(I_0)$$

BP at 0 temperature:
keep iterating...



$$I_2 = P_G(I_1)$$

BP at 0 temperature: fixed-point



$$I_3 = P_G(I_2) = P_G(I_3) \Rightarrow I^* = I_3$$

Problem solved?

- On trees, can recover the size of a maximum b-matching from I^*
- However, we are interested in sequences of graphs that are asymptotically tree-like
 - Message passing may not converge on those
 - Possibly many fixed-points on an infinite tree
 - Some of them will not yield a maximum b-matching...
- BP at positive temperature T and then $T \rightarrow 0$

BP at positive temperature

- Gibbs measure on set of b-matchings

$$\mu_G^\lambda(\boldsymbol{\sigma}) = \frac{1}{Z_G(\lambda)} \lambda^{\sum_e \sigma_e} \prod_{v \in V} 1\left(\sum_{e \sim v} \sigma_e\right)$$

- At positive temperature, **BP messages are distributions over the integers**

– Initialization: $m_{v \rightarrow u}^{(0)} = (1, 0, \dots)$

– BP update: $m_{v \rightarrow u}^{(k+1)}(x) = R_{v \rightarrow u}^\lambda[\mathbf{m}^{(k)}](x)$

$$R_{v \rightarrow u}^\lambda[\mathbf{m}](x) = C \lambda^x \sum_{\boldsymbol{\sigma} \in \mathbb{N}^{\partial v \setminus u}, |\boldsymbol{\sigma}| \leq b_v - x} \prod_{w \in \partial v \setminus u} m_{w \rightarrow v}(\sigma_{vw})$$

where C is a normalization factor

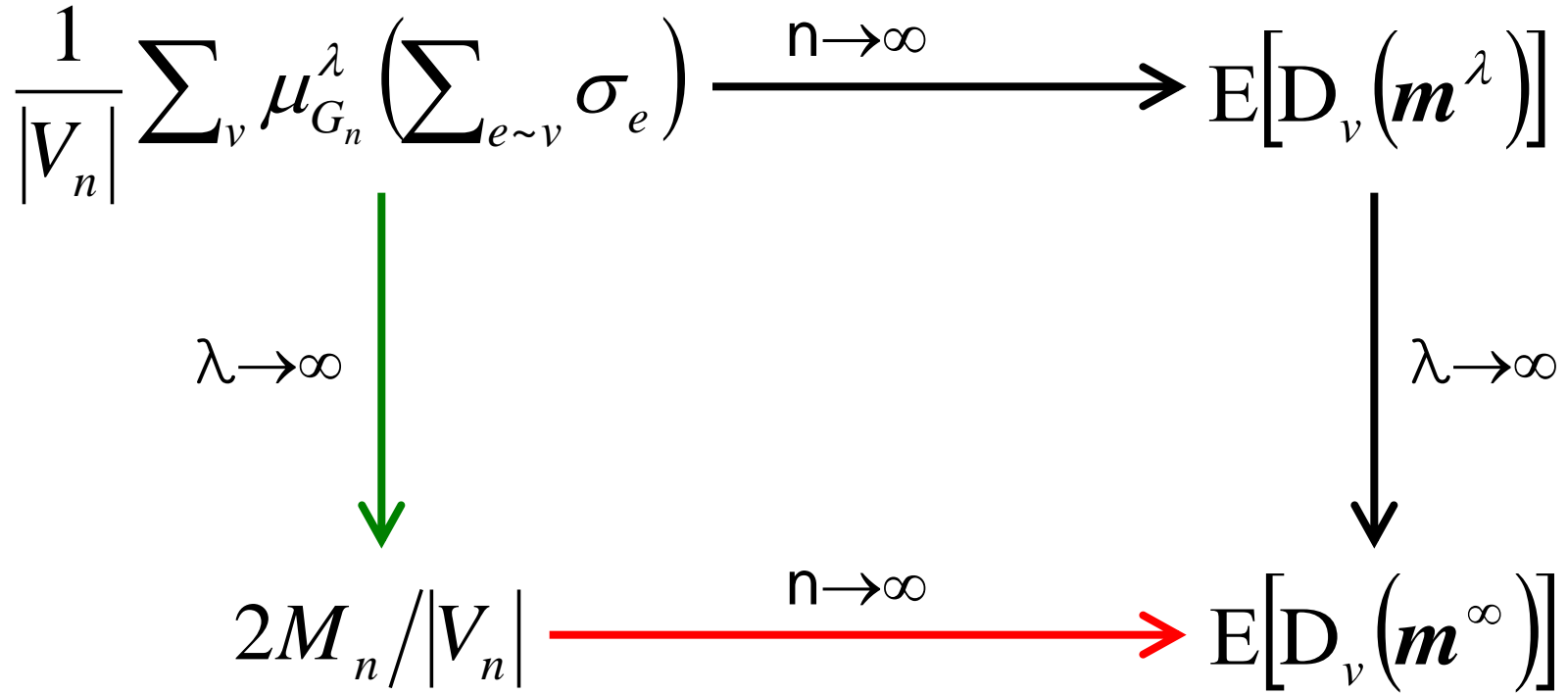
BP estimate

- Define $D_v(\mathbf{m}^\lambda)$ = BP estimate of total usage of edges adjacent to v , where \mathbf{m}^λ fixed-point of BP at temperature λ

$$D_v(\mathbf{m}^\lambda) = \frac{\sum_{\sigma \in \mathcal{N}^{\partial v}, |\sigma| \leq b_v} |\sigma| \prod_{u \in \partial v} m_{u \rightarrow v}^\lambda(\sigma_{uv})}{\sum_{\sigma \in \mathcal{N}^{\partial v}, |\sigma| \leq b_v} \prod_{u \in \partial v} m_{u \rightarrow v}^\lambda(\sigma_{uv})}$$

- $D_v(\mathbf{m}^\lambda) = \mu_G^\lambda \left(\sum_{e \sim v} \sigma_e \right)$ when G is a finite tree

Global plan



Global plan

$$\frac{1}{|V_n|} \sum_v \mu_{G_n}^\lambda \left(\sum_{e \sim v} \sigma_e \right) \xrightarrow{n \rightarrow \infty} \mathbb{E}[D_v(m^\lambda)]$$

$\lambda \rightarrow \infty$



$$2M_n / |V_n|$$

$n \rightarrow \infty$



$$\mathbb{E}[D_v(m^\infty)]$$

$\lambda \rightarrow \infty$



finite graphs

BP on infinite tree

At the end...an ugly formula

$$\lim_{n \rightarrow \infty} \frac{2M_n}{|V_n|} = \inf_{q=g \circ g(q)} F(q)$$

where $F(q) = \mathbb{E} \left[b_v 1(|\partial v| > 0) + b_v \wedge \sum_{i=1}^{|\partial v|} Y_i - b_v \wedge \sum_{i=1}^{|\partial v|} \left(b_v - \sum_{j \neq i} X_j \right)^+ \right]$

$$X_i \sim q \text{ and } Y_i \sim \mathbb{E}[P_{v \rightarrow u}(X)] \sim g(q)$$

At the end...an ugly formula

two-step fixed-point

$$\lim_{n \rightarrow \infty} \frac{2M_n}{|V_n|} = \inf_{q=g \circ g(q)} F(q)$$

where $F(q) = \mathbb{E} \left[\begin{array}{l} \boxed{b_v} \mathbb{1}(|\partial v| > 0) + b_v \wedge \boxed{\sum_{i=1}^{|\partial v|} Y_i} \\ - b_v \wedge \boxed{\sum_{i=1}^{|\partial v|} \left(b_v - \sum_{j \neq i} X_j \right)^+} \end{array} \right]$

- outgoing messages

$$X_i \sim q \text{ and } Y_i \sim \mathbb{E}[\mathbb{P}_{v \rightarrow u}(\mathbf{X})] \sim g(q)$$

Convergence of BP at $T > 0$

- **Does BP converge?**

- matching case: messages are distributions on $\{0,1\}$
→ can be encoded by one real number

- **Negative dependence** (Pemantle) \Rightarrow BP update operator non-increasing (Salez) \Rightarrow adjacent sequences \Rightarrow [...] \Rightarrow **unique fixed-point**

- **Is $R_{v \rightarrow u}^\lambda$ non-increasing?**

- For general messages and *st*-order, **no!**

- Restrict to log-concave distributions

- $m_{v \rightarrow u}(x)/m_{v \rightarrow u}(x-1)$ non-increasing in $x \in \mathbb{N}^*$

- **Totally Positive functions** (Karlin) $\Rightarrow R_{v \rightarrow u}^\lambda$ **non-increasing for *lr*-order** \Rightarrow [...] \Rightarrow **unique fixed-point**

Merci!

Questions?

Applications: cuckoo hashing

Introduced by Pagh & Rodler, ESA'01

- m balls and n bins (mapping of objects and keys)
- Each ball is proposed 2 bins at random
- How large can m be such that it is possible to put each ball into one of its proposed bin, with no collisions allowed?
- Generalizations:
 - More choices per ball → still a matching problem
 - Larger bin capacity → $b > 1$ on one side
 - Batches of balls with same choices → $b > 1$ on both sides \Rightarrow edges may be used multiple times!

Monotonicity as $T \rightarrow 0$

$$\begin{array}{ccc}
 \frac{1}{|V_n|} \sum_v \mu_{G_n}^\lambda \left(\sum_{e \sim v} \sigma_e \right) & \xrightarrow{n \rightarrow \infty} & \mathbb{E}[D_v(m^\lambda)] \\
 \downarrow \lambda \rightarrow \infty & & \downarrow \lambda \rightarrow \infty \\
 2M_n / |V_n| & \xrightarrow{n \rightarrow \infty} & \mathbb{E}[D_v(m^\infty)]
 \end{array}$$

- Limit of $\mu_{G_n}^\lambda$ as $\lambda \rightarrow \infty$ is the uniform law over the set of maximum b-matchings

Correlation decay at $T > 0$

$$\begin{array}{ccc}
 \frac{1}{|V_n|} \sum_v \mu_{G_n}^\lambda \left(\sum_{e \sim v} \sigma_e \right) & \xrightarrow{n \rightarrow \infty} & \mathbb{E} \left[D_v \left(m^\lambda \right) \right] \\
 \downarrow \lambda \rightarrow \infty & & \downarrow \lambda \rightarrow \infty \\
 2M_n / |V_n| & \xrightarrow{n \rightarrow \infty} & \mathbb{E} \left[D_v \left(m^\infty \right) \right]
 \end{array}$$

- Unicity of BP fixed-point at $T > 0$
- G_n converges locally weakly to G
- G is a tree, hence BP estimate is correct

Monotonicity + continuity

$$\begin{array}{ccc} \frac{1}{|V_n|} \sum_v \mu_{G_n}^\lambda \left(\sum_{e \sim v} \sigma_e \right) & \xrightarrow{n \rightarrow \infty} & \mathbb{E}[D_v(m^\lambda)] \\ \downarrow \lambda \rightarrow \infty & & \downarrow \lambda \rightarrow \infty \\ 2M_n / |V_n| & \xrightarrow{n \rightarrow \infty} & \mathbb{E}[D_v(m^\infty)] \end{array}$$

- $m^\infty = \lim_{\lambda \rightarrow \infty} \uparrow m^\lambda$
- D_v continuous

Convexity argument

$$\begin{array}{ccc}
 \frac{1}{|V_n|} \sum_v \mu_{G_n}^\lambda \left(\sum_{e \sim v} \sigma_e \right) & \xrightarrow{n \rightarrow \infty} & \mathbb{E} \left[D_v \left(m^\lambda \right) \right] \\
 \downarrow \lambda \rightarrow \infty & & \downarrow \lambda \rightarrow \infty \\
 2M_n / |V_n| & \xrightarrow{n \rightarrow \infty} & \mathbb{E} \left[D_v \left(m^\infty \right) \right]
 \end{array}$$

- Uniform control in n
- Entropy term in the free energy becomes negligible as $T \rightarrow 0$

Non-increasing operator?

$$\text{Ex: } b_v=3, \partial v = \{u_1, u_2, u_3\}, m_{u_1 \rightarrow v} = \left(\frac{1}{2}, \frac{1}{2}, 0, 0\right) = m'_{u_1 \rightarrow v}$$

$$\text{and } m_{u_2 \rightarrow v} = \left(\frac{1}{4}, \frac{1}{4}, \frac{1}{3}, \frac{1}{6}\right) <_{st} m'_{u_2 \rightarrow v} = \left(\frac{1}{4}, \frac{1}{4}, \frac{1}{6}, \frac{1}{3}\right)$$

$$\mathbf{R}_{v \rightarrow u_3}^1[m] = \frac{12}{25} \left(\frac{11}{12}, \frac{8}{12}, \frac{3}{8}, \frac{1}{8}\right) <_{st} \mathbf{R}_{v \rightarrow u_3}^1[m'] = \frac{12}{23} \left(\frac{10}{12}, \frac{7}{12}, \frac{3}{8}, \frac{1}{8}\right)$$

\Rightarrow no monotonicity for general messages