

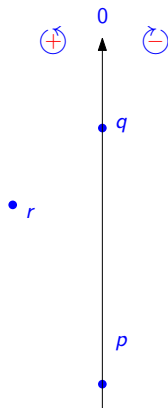
A canonical tree decomposition for chirotopes

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Joint work with Mathilde Bouvel, Valentin Féray, Xavier Goaoc

CALIN Seminar, LIPN, 2024
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Orientation, chirotope of labelled points



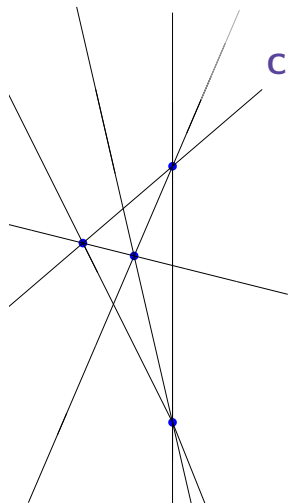
Orientation of three labelled points

$$\chi(p, q, r) = \begin{cases} +1 & \text{if } p_p, p_q, p_r \text{ oriented CCW,} \\ -1 & \text{if } p_p, p_q, p_r \text{ oriented CW,} \\ 0 & \text{if } p_p, p_q, p_r \text{ aligned.} \end{cases}$$

$$\chi(p, q, r) = \text{sign} \begin{vmatrix} x_p & x_q & x_r \\ y_p & y_q & y_r \\ 1 & 1 & 1 \end{vmatrix}$$

Remark: $\chi(p, q, r) = \chi(q, r, p) = \chi(r, p, q)$

Orientation, chirotope of labelled points

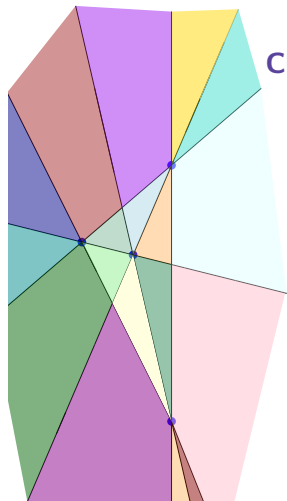


Chirotope of labelled points set

- Point set $\mathcal{P} = \{p_\ell\}_{\ell \in X}$ labelled by X
- Points in **general position** (no three aligned, no parallel lines)
- $\chi_{\mathcal{P}} : (X)_3 \rightarrow \{-1, +1\}$

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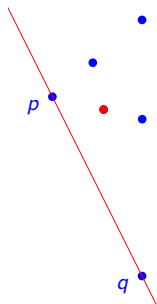
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- Chirotopes allow to abstract from coordinates

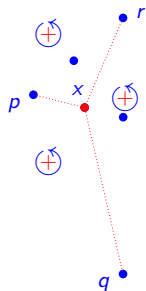
Chirotopes encode many useful properties

Example of properties

- $[p, q] \in \text{Conv}(\mathcal{P})$
 $\Leftrightarrow \forall x \in X \setminus \{p, q\}, \chi(p, q, x) = \text{cst}$



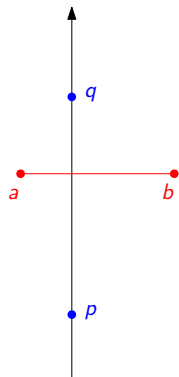
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- (p, q) separates $[a, b]$
 $\Leftrightarrow \chi(p, q, a) = -\chi(p, q, b)$
- $[p, q]$ and $[a, b]$ are intersecting $\Leftrightarrow \dots$

Chirotopes

Chirotopes are a **useful** combinatorial and geometric object...

- Finite number t_n of chirotopes on n elements



- Useful for **exact algorithms**, not depending on coordinates
- Can be used for benchmarking algorithms

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...but also quite **complex** to understand!

- Decide whether $f : (X)_3 \rightarrow \{-1, 1\}$ is **realizable** is **NP-hard** [Shor 91]
- Number of chirotopes t_n exactly known only for $n \leq 11$ (up to relabelling) [Aichholzer et al 2002]
- Best known asymptotics $t_n = n^{4n + \Theta(n/\log n)}$ [Goodman and Pollak 93]

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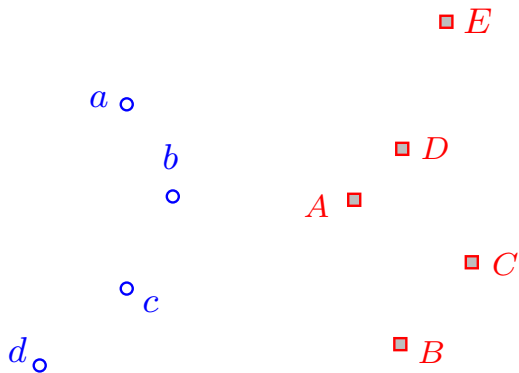
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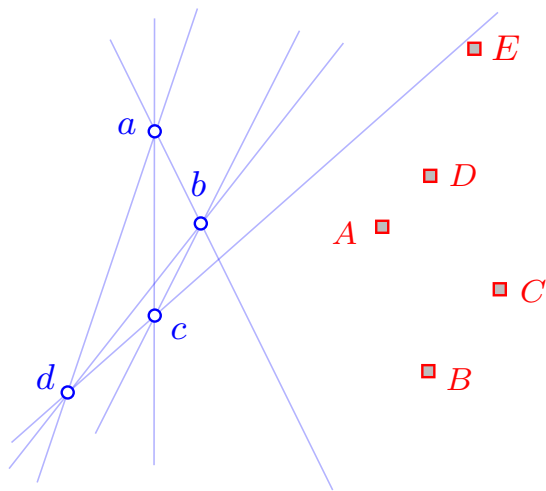
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The space of chirotopes is hard to explore!

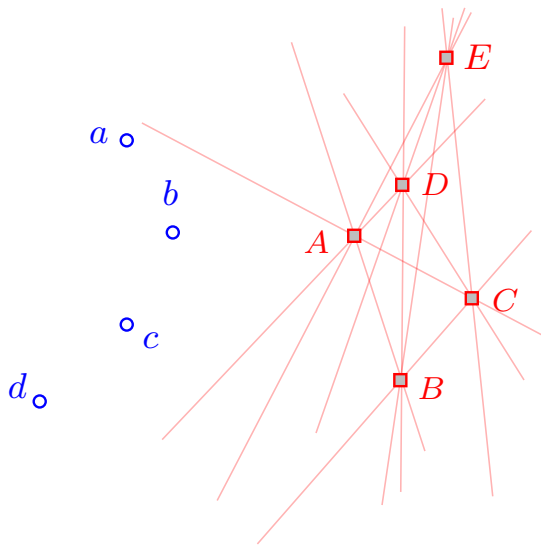
Mutually avoiding sets, modular decomposition



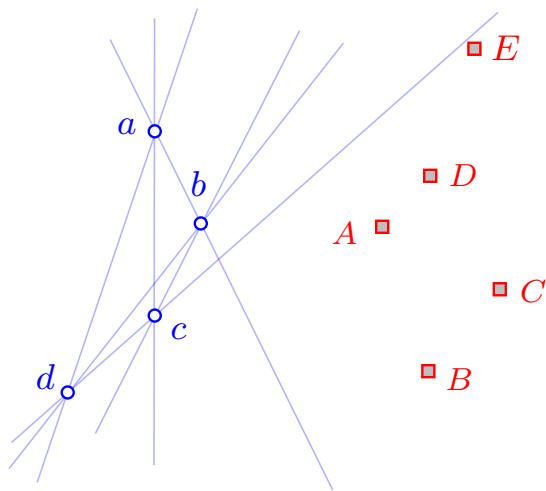
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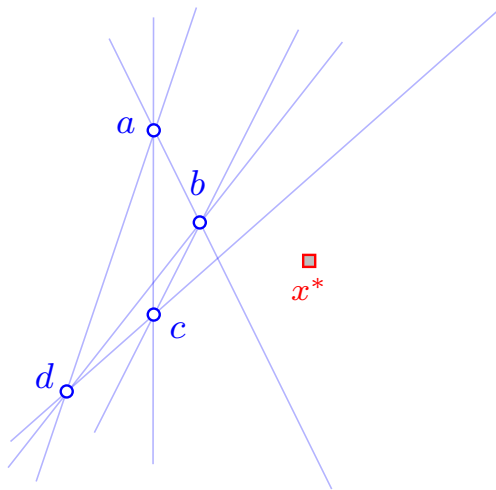
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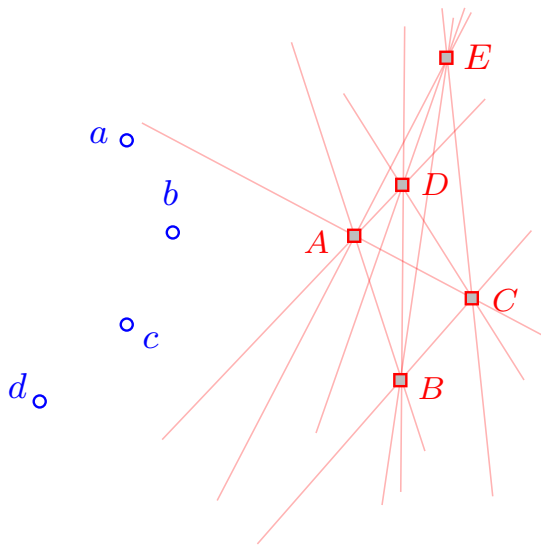
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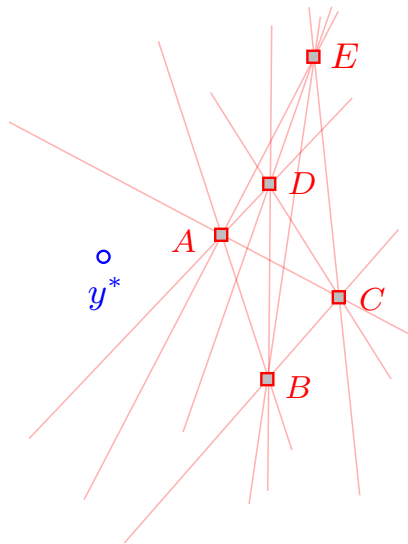
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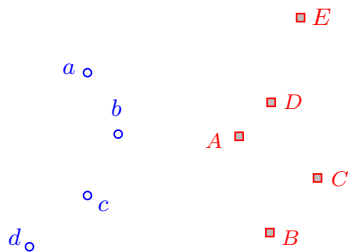
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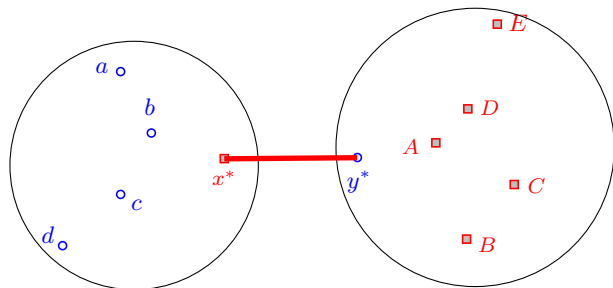
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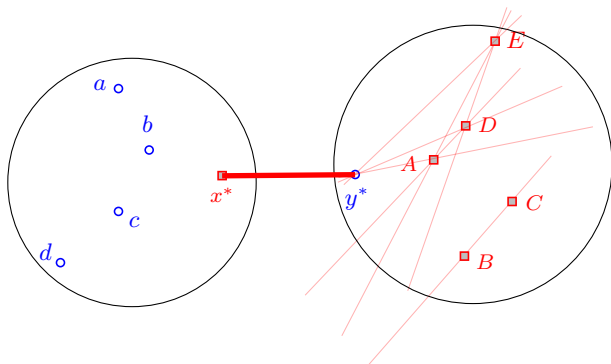
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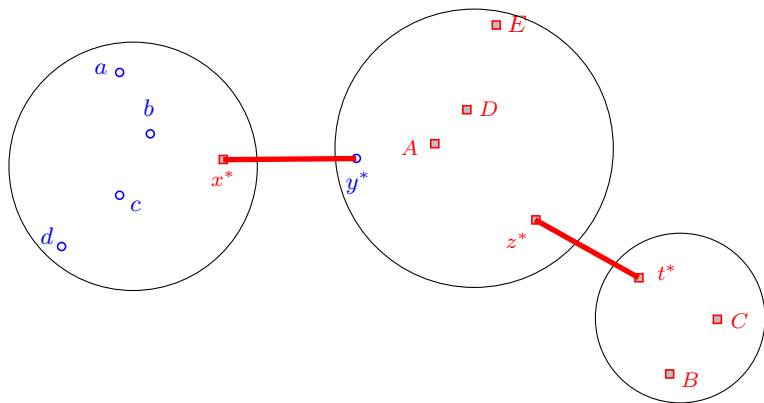
Chirotope is read following proxies!

$$\chi(a, d, B) = \chi(a, d, x^*)$$

Mutually avoiding sets, modular decomposition



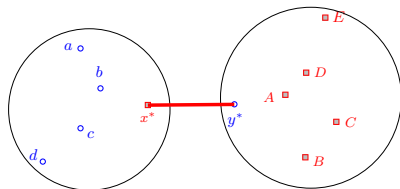
Mutually avoiding sets, modular decomposition



Chirotope is read following proxies!

$$\chi(a, d, B) = \chi(a, d, x^*) \quad \chi(a, E, B) = \chi(y^*, E, z^*)$$

Converse operation: Bowtie operation



- χ sign function on $X \cup \{x^*\}$, and ξ on $Y \cup \{y^*\}$

- the **bowtie** $\kappa \stackrel{\text{def}}{=} \chi_{x^*} \bowtie_{y^*} \xi$ is defined on $X \cup Y$ by:

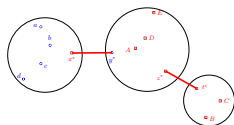
$$\left\{ \begin{array}{ll} \kappa(x_1, x_2, x_3) = \chi(x_1, x_2, x_3) & \text{if } x_1, x_2, x_3 \text{ are all in } X; \\ \kappa(x_1, x_2, y) = \chi(x_1, x_2, x^*) & \text{if } x_1, x_2 \text{ are in } X \text{ and } y \text{ is in } Y; \\ \kappa(x, y_2, y_3) = \xi(y^*, y_2, y_3) & \text{if } x \text{ is in } X \text{ and } y_2, y_3 \text{ are in } Y; \\ \kappa(y_1, y_2, y_3) = \xi(y_1, y_2, y_3) & \text{if } y_1, y_2, y_3 \text{ are all in } Y. \end{array} \right.$$

- $\chi_{x^*} \bowtie_{y^*} \xi$ is a realizable chirotope if and only if χ and ξ are realizable and x^* and y^* are **extreme** in χ and ξ .

[Bouvel, Féray, Goaoc, K.]

First properties of decomposition

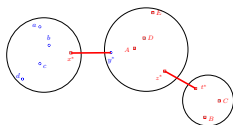
- κ is **indecomposable** if there is no nontrivial decomposition
 $\kappa = \chi_{x^*} \bowtie_{y^*} \xi$
- Every chirotope admits a decomposition built from indecomposable chirotopes
- Every decomposition can be represented by a (indecomposable) **chirotope tree**



- It allows a nice description of a **realizable** chirotope while avoiding providing a full realization

First properties of decomposition

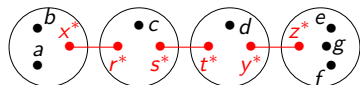
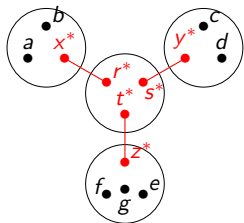
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- It allows a nice description of a **realizable** chirotope while avoiding providing a full realization
- **Is the chirotope tree canonical (unique)?**

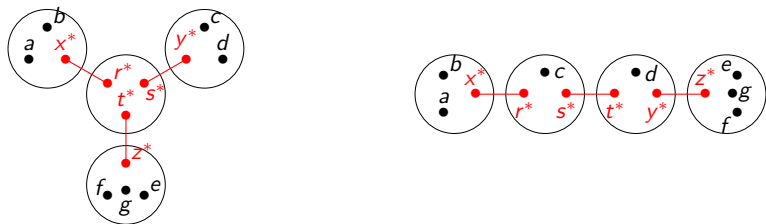
Convex issue

Two trees with the **same** associated chirotope:



Convex issue

Two trees with the **same** associated chirotope:



But if we merge / recompose adjacent convex nodes, same tree:



Unicity of decomposition

A chirotope tree is **canonical**:

- If every node is either **convex** or **indecomposable**
- There is **no edge** between two convex nodes

Proposition: Every chirotope admits a **unique canonical tree** (up to relabelling the proxies) [BFGK.]

Proof sketch

Define two transformations on chirotope trees:

- \diamondrightarrow that merges two adjacent convex nodes
- \boxtimesrightarrow that decomposes a nonconvex node in two

Proposition: [BFGK.] The transformation $\Rightarrow := \diamondrightarrow \cup \boxtimesrightarrow$ terminates and is locally confluent: if $T \Rightarrow T_1$ and $T \Rightarrow T_2$ then there exists T_3 such that $T_1 \Rightarrow^* T_3$ and $T_2 \Rightarrow^* T_3$.

The main **difficulty** resides in the case where a node can be decomposed in two manners.

Good/Bad news

Good news:

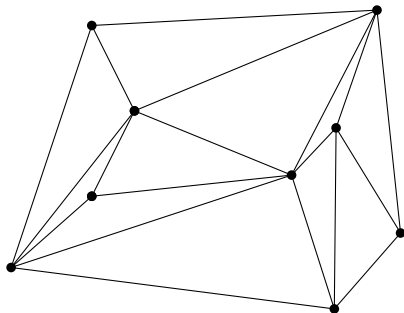
- Chirotope trees provide a nice way to build chirotopes from smaller ones
- Unicity gives a simple way to prove that two chirotopes built recursively are different

Bad news: For n large enough we have [BFGK.]

$$d_n/t_n = \mathcal{O}(n^{-3})$$

Triangulations

A **triangulation** of a point set \mathcal{P} is a maximal **crossing-free** set of edges between elements of \mathcal{P} .



The set of triangulations of \mathcal{P} only depends on its chirotope.

We write \mathcal{T}_κ the set of triangulations of a chirotope κ .

Triangulations

Many questions are still **open**:

- For every κ on n elements,

$$|\mathcal{T}_\kappa| = \mathcal{O}(30^n) \text{ [Sharir and Sheffer, 2011]}$$

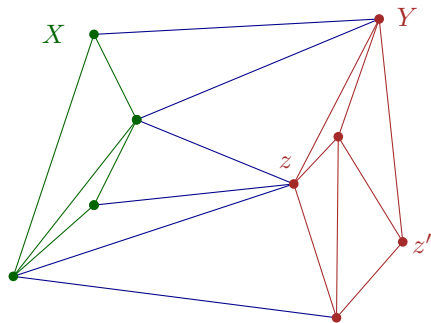
$$|\mathcal{T}_\kappa| = \Omega(2.63^n) \text{ [Aichholzer et al, 2016]}$$

- It is conjectured that the minimal is $\mathcal{O}(3.47^n)$ [Hurtado and Noy 97]
- Max known: Koch chains has $\approx 9.08^n$ triangulations [Rutschmann and Wettstein 2022]

Algorithmically: Compute the number of triangulations of a given point set:

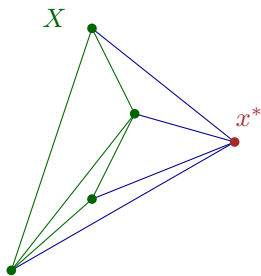
- $\mathcal{O}(n^2 2^n)$ [Alvarez and Seidel 2013]
- $\mathcal{O}(n^{(11+o(1))\sqrt{n}})$ [Marx and Miltzow 2016]

Triangulations of bowties



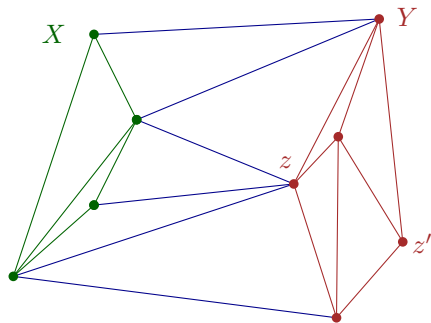
$$T \in \mathcal{T}_K$$

Triangulations of bowties



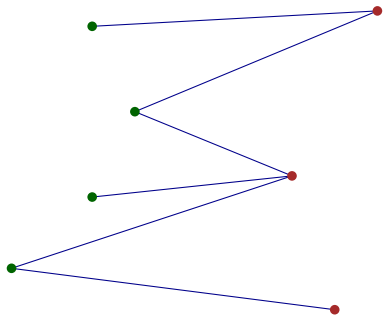
$$T' = \pi_{Y \rightarrow x^*}(T)$$

Triangulations of bowties



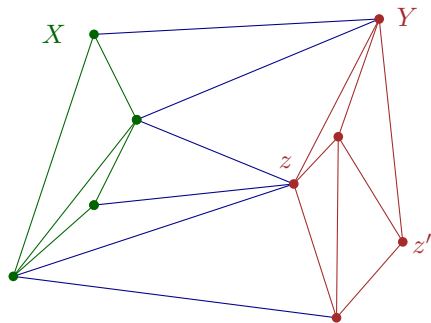
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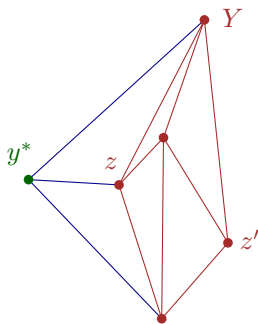
T_{XY}

Triangulations of bowties



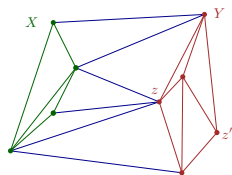
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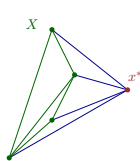


$$T'' = \pi_{X \rightarrow y^*}(T)$$

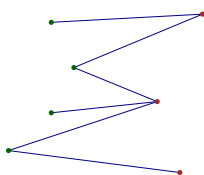
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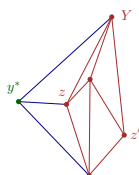
$T \in \mathcal{T}_\kappa$



$T' = \pi_{Y \rightarrow x^*}(T)$



T_{XY}



$T'' = \pi_{X \rightarrow y^*}(T)$

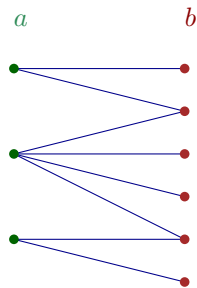
Bijection between:

- Triangulations of $\kappa = \chi_{x^*} \bowtie_{y^*} \xi$
- Triplets of triangulations of χ , ξ , and maximal crossing-free families of edged between the neighbors of x^* and y^*

$$|\mathcal{T}_\kappa| = \sum_{a,b \geq 2} \binom{a+b-2}{a-1} [x^a] P_{\chi, x^*}(x) [y^b] P_{\xi, y^*}(y).$$

Triangulations of bowties

Two ingredients:



$|T_{XY}|$ only depends on a and b :

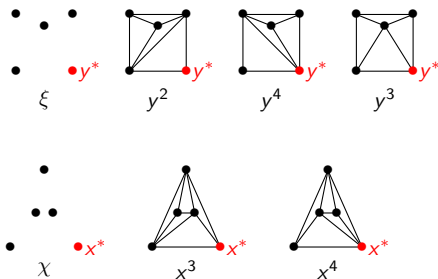
$$|T_{XY}| = \left(\binom{b}{a-1} \right) = \binom{a+b-2}{a-1}$$

Triangulations of bowties

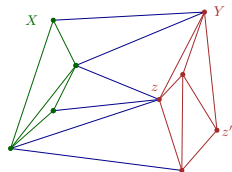
Two ingredients: Triangulation polynomial

$$P_{\chi, x^*}(x) = \sum_{T \in \mathcal{T}_\chi} x^{\deg_T(x^*)} \quad P_{\chi, x^*}(1) = |\mathcal{T}_\chi|$$

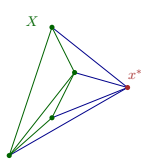
Example: $P_{\chi, x^*}(x) = x^3(x + 1)$ and $P_{\xi, y^*}(y) = y^2(1 + y + y^2)$.



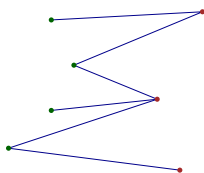
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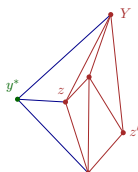
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T_{XY}



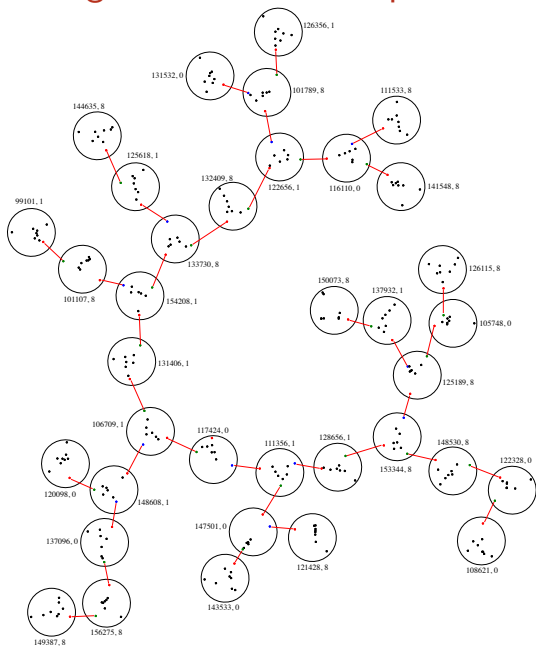
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Triangulations of chirotope trees



36 nodes, each decorated with a chirotope of size 9, adding up to 254 elements.

Number of triangulations?

Triangulations of chirotope trees

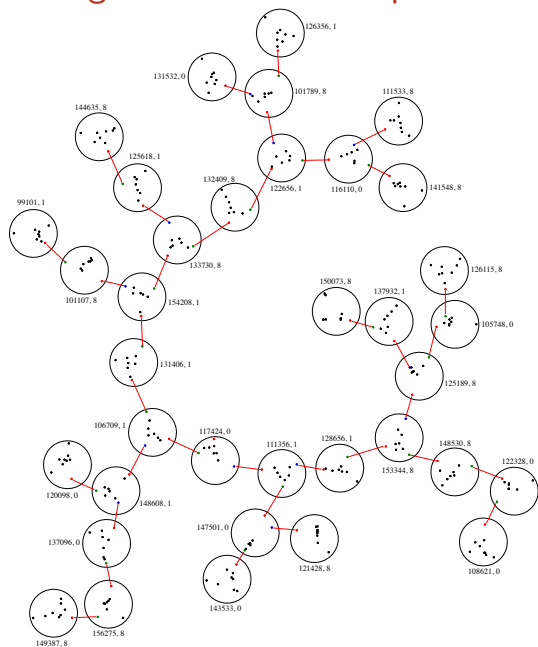
- Same ideas with the same bijection.
- However the full triangulation polynomial of every node is needed:

$$Q_{\xi, \{x_i^*\}}(x_1, \dots, x_k, \{y_{i,j}\}_{1 \leq i < j \leq k}) = \sum_{T \in \mathcal{T}_\xi} \prod_{i=1}^k x_i^{\deg_T(x_i^*)} \cdot \prod_{\substack{x_i^* x_j^* \in T \\ i < j}} y_{i,j}.$$

- and the formula are more complex, as we compute multivariate polynomials instead of just a number.

Conclusion: **[BFGK.]** the number of triangulation of a chirotope tree can be computed in **polynomial time** from the **full triangulation polynomials** of its nodes.

Triangulations of chirotope trees



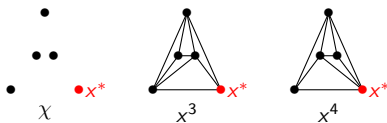
36 nodes, each decorated with a chirotope of size 9, adding up to **254** elements.

Number of triangulations:

$$|T_{\mathcal{K}}| \approx 5.92966751 \cdot 10^{180}$$

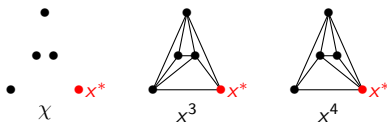
computed **exactly** in a few seconds using sage.

Computing triangulation polynomials

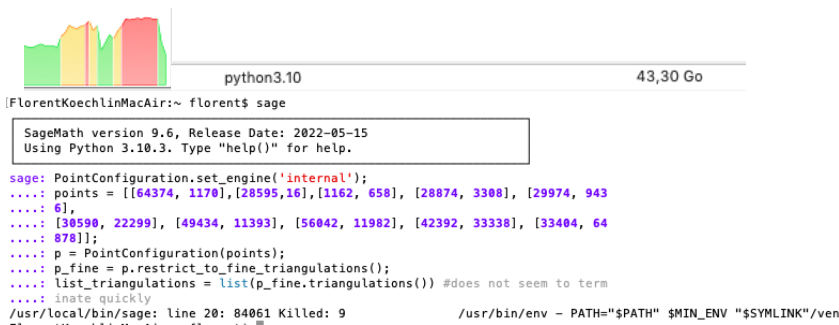


- First idea: enumerate all triangulations with Sage and deduce the polynomial

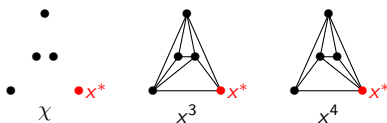
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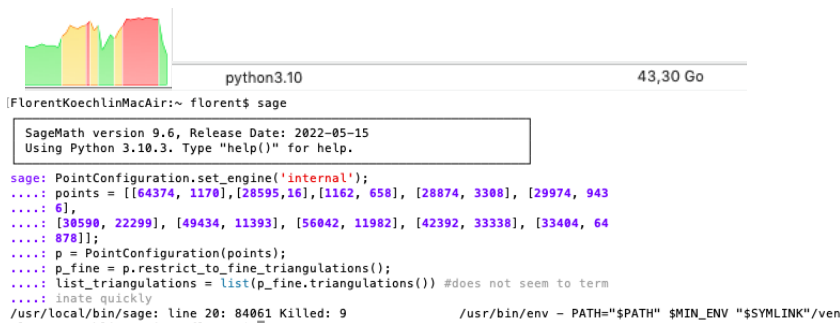
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Computing triangulation polynomials

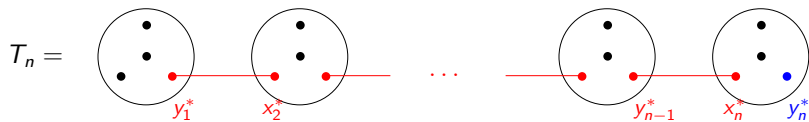


- First idea: enumerate all triangulations with Sage and deduce the polynomial → **bug in Sage!**

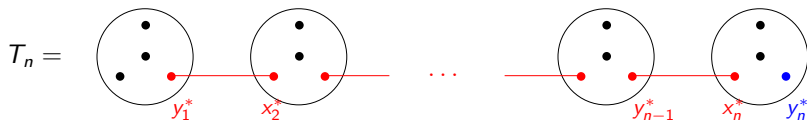


- Better idea: adapt the $\mathcal{O}(n^2 2^n)$ algorithm of [Alvarez and Seidel 2013] to compute the polynomial

Chain Triangulations: an analyzable case

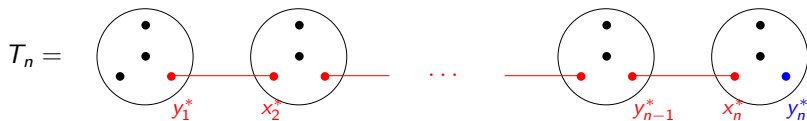


Chain Triangulations: an analyzable case



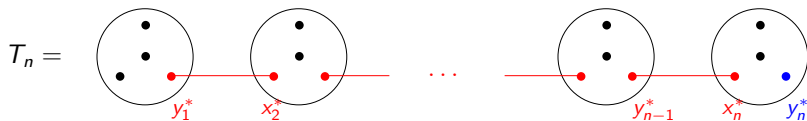
- $\chi_{n+1} = \chi_n y_n \bowtie_{x_{n+1}} \chi_1$
- We introduce $P_n(y) \stackrel{\text{def}}{=} P_{\chi_n, y_n^*}(y)$ and $Q_{\chi_1, x_n^*, y_n^*}(x, y) = x^3 y^3$

Chain Triangulations: an analyzable case

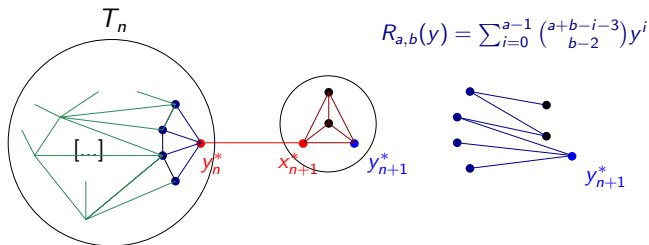


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Chain Triangulations: an analyzable case



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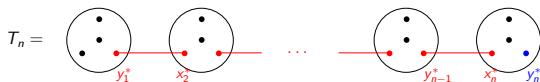
Chain Triangulations: an analyzable case



- $P_{n+1}(y) - \frac{y^4}{(1-y)^2} P_n(y) = -\frac{y^4}{(1-y)^2} P_n(1) + \frac{y^3}{1-y} P'_n(1)$
- Let us introduce $F(y, u) \stackrel{\text{def}}{=} \sum_{k \geq 1} P_k(y) u^k$

$$\left(1 - \frac{uy^4}{(1-y)^2}\right) F(y, u) = uy^3 \left(1 - \frac{y}{(1-y)^2} F(1, u) + \frac{1}{1-y} \partial_y F(1, u)\right)$$

Chain Triangulations: an analyzable case



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- **Kernel method:** find $y(u)$ analytic canceling the kernel.
- We obtain $F(1, u) = \sum_{n \geq 1} |\mathcal{T}_{\chi_n}| u^n$ (A066357 !) and

$$|\mathcal{T}_{\chi_n}| \sim_{n \rightarrow \infty} \frac{3 - 2\sqrt{2}}{\sqrt{2\pi}} \frac{16^n}{n^{3/2}}.$$

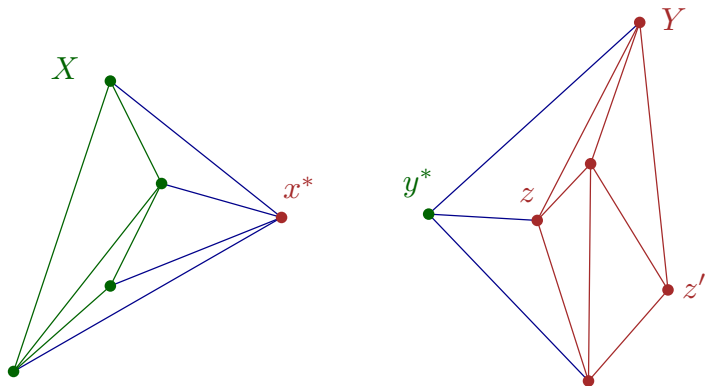
Conclusion

- We have seen a canonical chirotope decomposition
- Natural sense of "factorizing" chirotope
- The decomposition can be used to compute the number of triangulations of complex chirotopes

Many open questions:

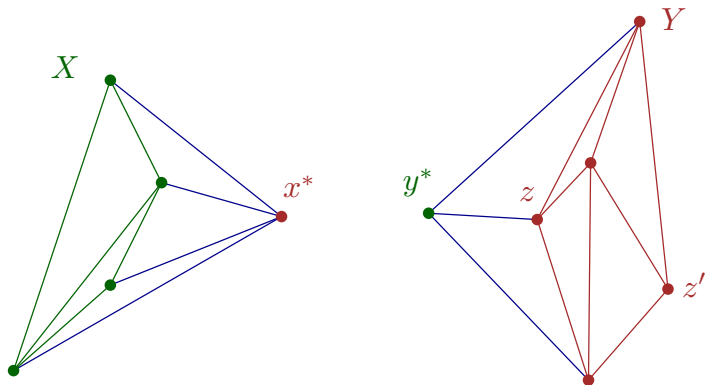
- What is the complexity of computing the canonical decomposition?
- Can more complex configurations be analyzed analytically?
- Can we unify it with other classical constructions?

Further work: faster computation



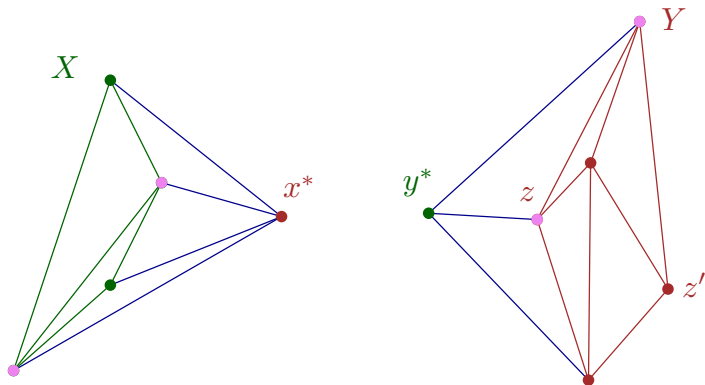
$$|\mathcal{T}_\kappa| = \sum_{a,b \geq 2} \binom{a+b-2}{a-1} [x^a] P_{\chi, x^*}(x) [y^b] P_{\xi, y^*}(y)$$

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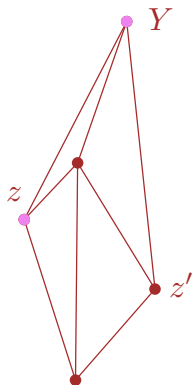
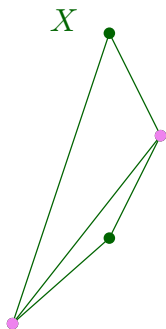
$$|\mathcal{T}_\kappa| = \left\langle \frac{P_{\chi, x^*}(x+1)}{x+1}, \frac{P_{\xi, y^*}(x+1)}{x+1} \right\rangle_x.$$

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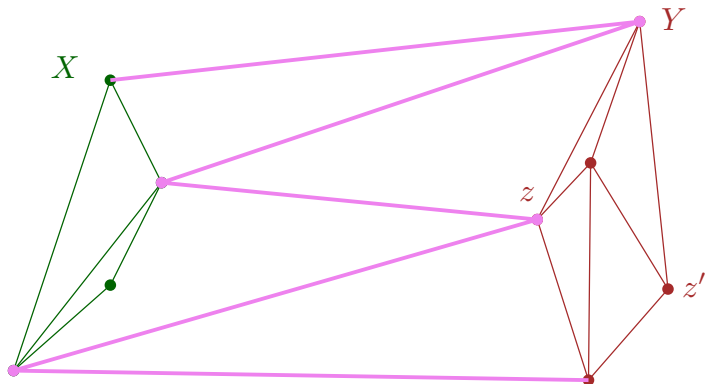
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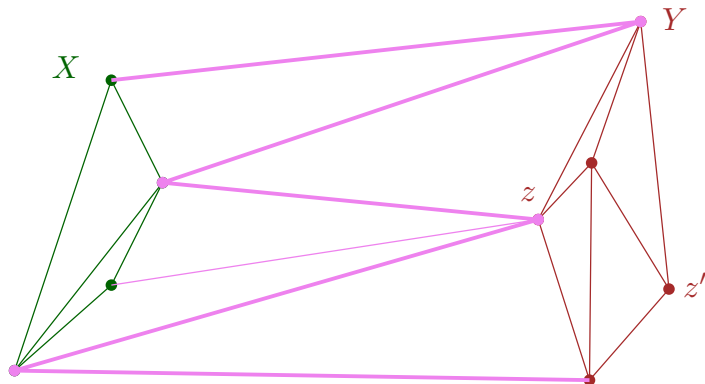
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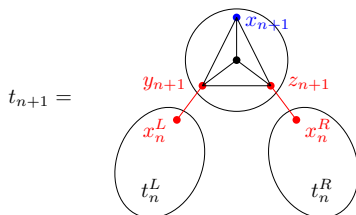
Further work: faster computation



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Further work: Binary chain

- Start with $\chi_1 =$ 
- Recursively build t_n :

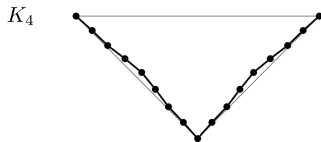


- Then if $Q_n(x) = (x + 1)P_n(x + 1)$:

$$Q_n(x) = \sum_{i=0}^s a_i x^i \Rightarrow Q_{n+1}(x) = (x + 1)^4 \sum_{k=2}^s \left(\sum_{i=k}^s a_i x^{i-k} \right)^2$$

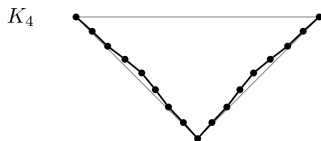
- Not analyzable for now, but greatly improves the possible number of iterations

Further work: Koch Chains [Rutschmann and Wettstein 2022]



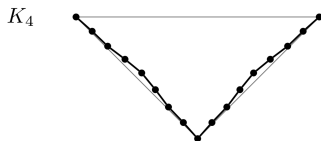
- **Chain:** x -increasing sequence of points x_1, \dots, x_n such that $[x_i x_{i+1}]$ is forced in every triangulation
- every known configuration with many triangulations is a chain
- Every chain can be decomposed with only two operators \wedge and \vee , and the basic chain of two points [Rutschmann and Wettstein 2022]

Further work: Koch Chains [Rutschmann and Wettstein 2022]



- **Koch chain**: best known configuration with maximal number of triangulations $\approx 9.08^n$
- Construction similar to ours (with "phantom" proxies)
- Similar techniques for counting triangulations
- We managed to **generalize** their construction beyond chains, but it seems that only chains reach the best number of triangulations

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Conclusion: still many things to look at!

Thank you!