

# Periods in action

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IHES

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## What is a period?

A **period** is the integral on a closed path of a rational function in one or several variables with *rational* coefficients.

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- coefficients in  $\mathbb{Q}$
- coefficients in  $\mathbb{C}(t)$ , **the period is a function of  $t$ .**

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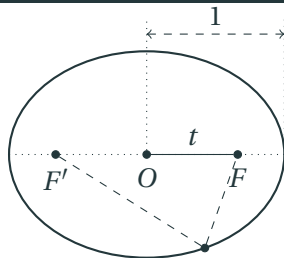




## An ellipse

eccentricity  $t$ 

major radius 1

perimeter  $E(t)$ 

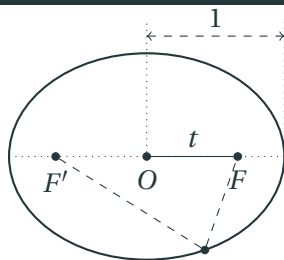
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$$E(t) = 2 \int_{-1}^1 \sqrt{\frac{1-t^2x^2}{1-x^2}} dx$$

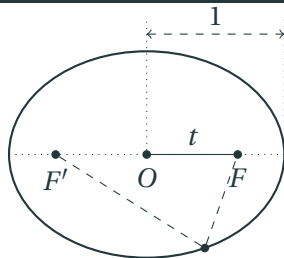


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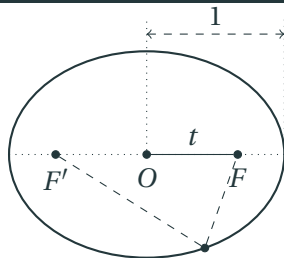


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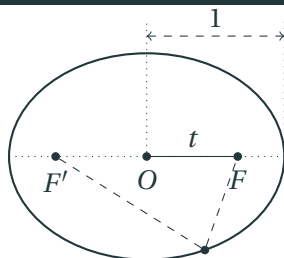
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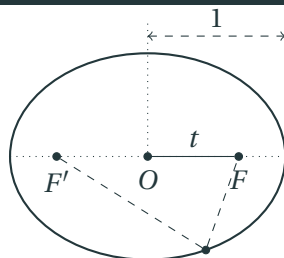
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**since then** Many applications in algebraic geometry (Gauß-Manin connection)  
 geometry of the cycles  $\leftrightarrow$  analytic properties of the periods

## Content

- ① Computing periods
- ② Multiple binomial sums
- ③ Volume of semialgebraic sets

## **Computing periods**

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**implicit**  $t(t-1)(64t-1)(3t-2)(6t+1)y''' + (4608t^4 - 6372t^3 + 813t^2 + 514t - 4)y'' + 4(576t^3 - 801t^2 - 108t + 74)y' = 0$  (+ init. cond.)



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- **equality testing**, given differential equations and initial conditions
- **numerical analytic continuation** with certified precision

(D. V. Chudnovsky and G. V. Chudnovsky 1990; van der Hoeven 1999; Mezzarobba 2010)

```
sage: from ore_algebra import *
sage: dop = (z^2+1)*Dz^2 + 2*z*Dz
sage: dop.numerical_solution(ini=[0,1], path=[0,1])
           [0.78539816339744831 +/- 1.08e-18]
sage: dop.numerical_solution(ini=[0,1], path=[0,i+1,2*i,i-1,0,1])
           [3.9269908169872415 +/- 4.81e-17] + [+/- 4.63e-21]*I
```

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One equation fits all cycles, the **Picard-Fuchs equation**.

$$\text{recall } E(t) = \oint \sqrt{\frac{1-t^2x^2}{1-x^2}} dx = \frac{1}{2\pi i} \oint \overbrace{\frac{1}{1 - \frac{1-t^2x^2}{(1-x^2)y^2}}}^{R(t,x,y)} dx dy$$

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### Computational proof

$$(t-t^3) \frac{\partial^2 R}{\partial t^2} + (1-t^2) \frac{\partial R}{\partial t} + tR =$$

$$\frac{\partial}{\partial x} \left( -\frac{t(-1-x+x^2+x^3)y^2(-3+2x+y^2+x^2(-2+3t^2-y^2))}{(-1+y^2+x^2(t^2-y^2))^2} \right) + \frac{\partial}{\partial y} \left( \frac{2t(-1+t^2)x(1+x^3)y^3}{(-1+y^2+x^2(t^2-y^2))^2} \right)$$

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*Problem (mostly) solved!*

# Multiple binomial sums

joint work with Alin Bostan and Bruno Salvy

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$$\sum_{k=0}^{2n} (-1)^k \binom{2n}{k}^3 = (-1)^n \frac{(3n)!}{(n!)^3} \quad (\text{Dixon})$$

$$\sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2 = \sum_{k=0}^n \binom{n}{k} \binom{n+k}{k} \sum_{j=0}^k \binom{k}{j}^3 \quad (\text{Strehl})$$

$$\sum_{i=0}^n \sum_{j=0}^n \binom{i+j}{i}^2 \binom{4n-2i-2j}{2n-2i} = (2n+1) \binom{2n}{n}^2$$

$$\sum_{r \geq 0} \sum_{s \geq 0} (-1)^{n+r+s} \binom{n}{r} \binom{n}{s} \binom{n+s}{s} \binom{n+r}{r} \binom{2n-r-s}{n} = \sum_{k \geq 0} \binom{n}{k}^4$$

$$\sum_i \sum_j \binom{2n}{n+i} \binom{2n}{n+j} |i^3 j^3 (i^2 - j^2)| = \frac{2n^4(n-1)(3n^2 - 6n + 2)}{(2n-3)(2n-1)} \binom{2n}{n}^2$$

Conjectured by Brent, Ohtsuka, Osborn, and Prodinger (2014)

$$1 + F_n^{-1,-1} + 2F_n^{0,0} - F_n^{0,1} + F_n^{1,0} - 3F_n^{1,1} + F_n^{1,2} - F_n^{3,1} + 3F_n^{3,2} - F_n^{3,3} - 2F_n^{4,2} + F_n^{4,3} - F_n^{5,2} = \sum_{m=0}^n \frac{\binom{n+2}{m} \binom{n+2}{m+1} \binom{n+2}{m+2}}{\binom{n+2}{1} \binom{n+2}{2}},$$

$$\text{where } F_n^{a,b} = \sum_{d=0}^{n-1} \sum_{c=0}^{d-a} \binom{d-a-c}{c} \binom{n}{d-a-c} \left( \binom{n+d+1-2a-2c+2b}{n-a-c+b} - \binom{n+d+1-2a-2c+2b}{n+1-a-c+b} \right)$$

Conjectured by Le Borgne

*Both proved using periods!*

## The not so formal grammar of binomial sums

○ → integer linear combination of the variables

$$\square \rightarrow \begin{pmatrix} \circ \\ \circ \end{pmatrix}$$

$$\square \rightarrow \text{Cst} \circ$$

$$\square \rightarrow \square + \square$$

$$\square \rightarrow \square \cdot \square$$

$$\square \rightarrow \sum_{n=\circ} \square$$

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integration  $t(27t+1)y'' + (54t+1)y' + 6y = 0$ , i.e.  $3(3n+2)(3n+1)u_n + (n+1)^2 u_{n+1} = 0$

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integration  $t(27t+1)y'' + (54t+1)y' + 6y = 0$ , i.e.  $3(3n+2)(3n+1)u_n + (n+1)^2 u_{n+1} = 0$

**conclusion** *Generating functions of binomial sums are periods!*

## Computing binomial sums with periods

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- Excellent running times, thanks to **simplification** and better algorithms for **integration**

## Binomial sums are diagonals of rational functions

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The converse does not hold, but...

- If  $(u_n)_{n \geq 0}$  is a binomial sum, then  $\sum_n u_n t^n$  is algebraic modulo  $p$  for all prime  $p$  (but finitely many).

$$y(t) \triangleq \sum_n \sum_{k=0}^n \binom{n+k}{k}^2 \binom{n}{k}^2 t^n$$
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**and of course**  $t^2(t^2 - 34t + 1)y''' + 3t(2t^2 - 51t + 1)y'' + (7t^2 - 112t + 1)y' + (t - 5)y = 0$ .

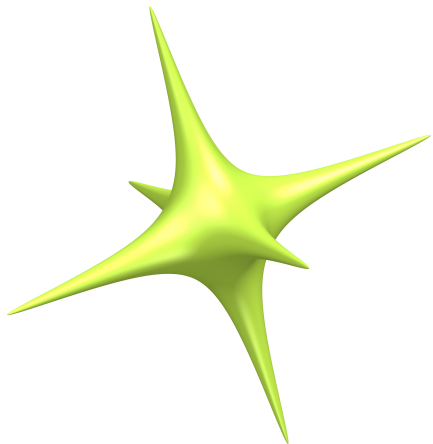
# **Volume of semialgebraic sets**

joint work with Mohab Safey El Din

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## A numeric integral

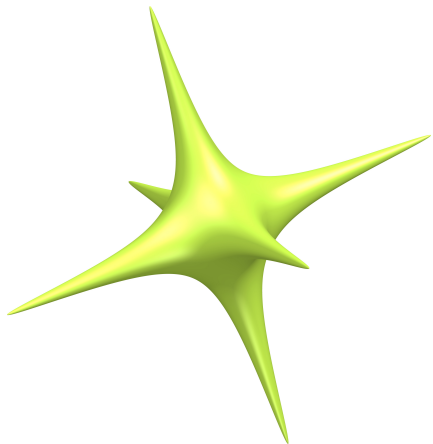
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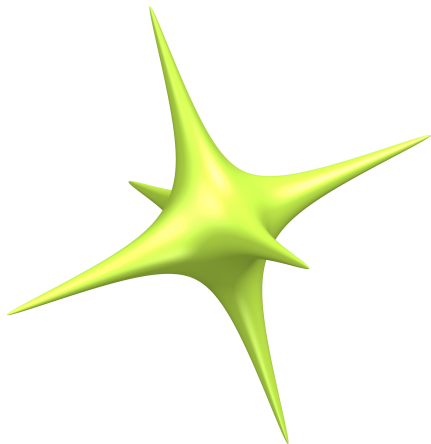


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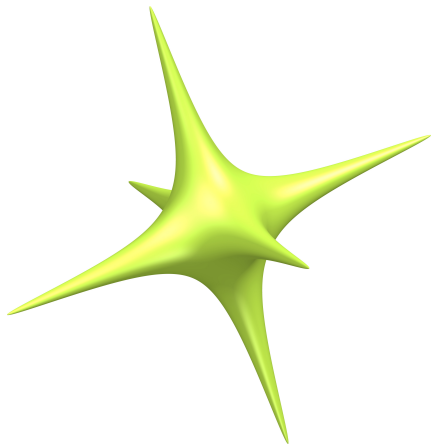
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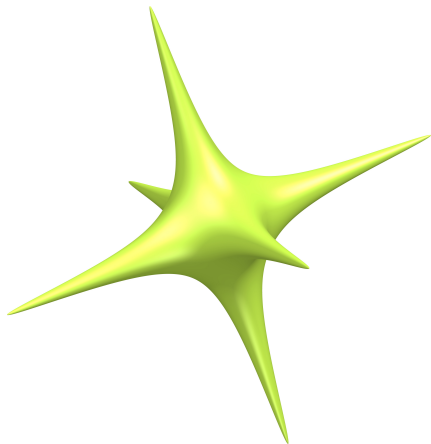
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- **No certification on precision**

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### Proposition

For any generic  $f \in \mathbb{R}[x_1, \dots, x_n]$ ,

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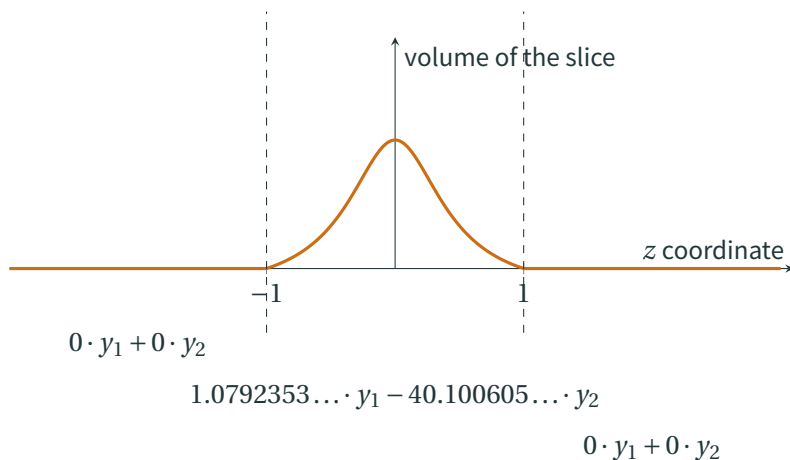
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## The “volume of a slice” function

$\{y_1, y_2\}$ , basis of the solution space of the Picard-Fuchs equation



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*The complexity is quasi-linear with respect to the precision!*

*(To get twice as many digits, you need only twice as much time.)*

## A hundred digits (within a minute)

$$\text{vol} \left( \text{star} \right) = 0.108575421460360937739503$$

395994207619810917874446  
607475444475822993285360  
673032928194943474414064  
066136624234627959808778  
1034932346781568...

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Questions?

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