

# Solving combinatorial equations via computer algebra

Combinatorics and Arithmetic for Physics, 15-17 November 2023

---

Hadrien Notarantonio (Inria Saclay – Sorbonne Université)

*Based on joint works with:*

*Alin Bostan (Inria Saclay)*

*Mohab Safey El Din (Sorbonne Université)*

*Sergey Yurkevich (University of Vienna)*



# Which **type of equations** are we looking at?

rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

**fixed-point** in  $F \rightsquigarrow$  **unique** solution in  $\mathbb{Q}[u][[t]]$

# Which **type of equations** are we looking at?

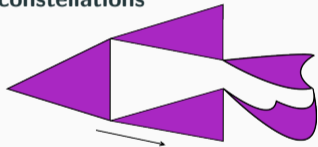
rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

fixed-point in  $F \rightsquigarrow$  **unique** solution in  $\mathbb{Q}[u][[t]]$

3-constellations



$$F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$$

# Which type of equations are we looking at?

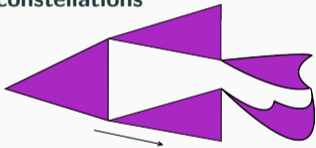
rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

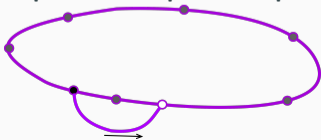
fixed-point in  $F \rightsquigarrow$  unique solution in  $\mathbb{Q}[u][[t]]$

3-constellations



$$F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$$

hard particles on planar maps



$$\begin{cases} F(t, u) = x - y + G(t, u) + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right) \\ G(t, u) = y + tsu \left( F(t, u)G(t, u) + \frac{G(t, u) - G(t, 1)}{u-1} \right) \end{cases}$$

## How to relate these **combinatorial objects** to such equations?

rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

[Tutte '68]

# How to relate these **combinatorial objects** to such equations?

rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

[Tutte '68]

$a_n := \# \{ \text{planar maps with } n \text{ edges} \}$

↓ refinement

$a_{n,d} := \# \{ \text{planar maps with } n \text{ edges,} \\ d \text{ of them on the external face} \}$

$$\sum_{n=0}^{\infty} a_n t^n$$

generating function

↓ refinement

$$F(t, u) := \sum_{n=0}^{\infty} \sum_{d=0}^n a_{n,d} u^d t^n \quad \text{complete generating function}$$

# How to relate these **combinatorial objects** to such equations?

rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

[Tutte '68]

$a_n := \# \{ \text{planar maps with } n \text{ edges} \}$

↓ refinement

$a_{n,d} := \# \{ \text{planar maps with } n \text{ edges,} \\ d \text{ of them on the external face} \}$

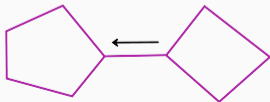
$$\sum_{n=0}^{\infty} a_n t^n$$

generating function

↓ refinement

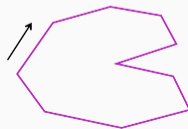
$$F(t, u) := \sum_{n=0}^{\infty} \sum_{d=0}^n a_{n,d} u^d t^n \quad \text{complete generating function}$$

•



1

$$tu^2 F(t, u)^2$$



$$tu \frac{uF(t, u) - F(t, 1)}{u-1}$$

# How to relate these **combinatorial objects** to such equations?

rooted planar maps



$$F(t, u) = 1 + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right)$$

[Tutte '68]

$a_n := \# \{ \text{planar maps with } n \text{ edges} \}$

↓ refinement

$a_{n,d} := \# \{ \text{planar maps with } n \text{ edges,} \\ d \text{ of them on the external face} \}$

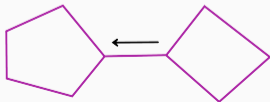
$$\sum_{n=0}^{\infty} a_n t^n$$

generating function

↓ refinement

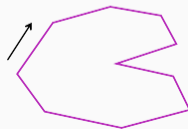
$$F(t, u) := \sum_{n=0}^{\infty} \sum_{d=0}^n a_{n,d} u^d t^n \quad \text{complete generating function}$$

•



1

$$tu^2 F(t, u)^2$$



$$tu \frac{uF(t, u) - F(t, 1)}{u-1}$$

$$F(t, 1) = \sum_{n=0}^{\infty} a_n t^n$$



## Solving functional equations

D-finite

Algebraic

Rational

$$\frac{1+6t}{1-2t+5t^2}$$

$$(1-t)^{\frac{1}{3}} - (1+2t)^{\frac{4}{5}}$$

$\exp(t)$

# Solving functional equations

D-finite

Algebraic

Rational

$$\frac{1+6t}{1-2t+5t^2}$$

$$(1-t)^{\frac{1}{3}} - (1+2t)^{\frac{4}{5}}$$

$\exp(t)$

In this talk

**Solving** = **Classifying** the initial series  $F(t, 1)$

+ **Computing** a **witness** of this classification

(e.g.  $R \in \mathbb{Q}[z, t]$  s.t.  $R(F(t, 1), t) = 0$ )

# Solving functional equations

D-finite

Algebraic

Rational

$$\frac{1+6t}{1-2t+5t^2}$$

$$(1-t)^{\frac{1}{3}} - (1+2t)^{\frac{4}{5}}$$

$\exp(t)$

In this talk

**Solving** = **Classifying** the initial series  $F(t, 1)$   
+ **Computing** a **witness** of this classification  
(e.g.  $R \in \mathbb{Q}[z, t]$  s.t.  $R(F(t, 1), t) = 0$ )

Going back to our planar maps...

$F(t, 1) = 1 + 2t + 9t^2 + 54t^3 + 378t^4 + \dots \in \mathbb{Q}[[t]]$   
annihilated by  $R = 27t^2z^2 + (1 - 18t)z + 16t - 1 \in \mathbb{Q}[z, t]$

# Solving functional equations

D-finite

Algebraic

Rational

$$\frac{1+6t}{1-2t+5t^2}$$

$$(1-t)^{\frac{1}{3}} - (1+2t)^{\frac{4}{5}}$$

$\exp(t)$

In this talk

**Solving** = **Classifying** the initial series  $F(t, 1)$   
+ **Computing** a **witness** of this classification  
(e.g.  $R \in \mathbb{Q}[z, t]$  s.t.  $R(F(t, 1), t) = 0$ )

Going back to our planar maps...

$F(t, 1) = 1 + 2t + 9t^2 + 54t^3 + 378t^4 + \dots \in \mathbb{Q}[[t]]$   
annihilated by  $R = 27t^2z^2 + (1 - 18t)z + 16t - 1 \in \mathbb{Q}[z, t]$

From  $R$ :

- (Recurrence)  $a_0 = 1$  and  $(n+3)a_{n+1} - 6(2n+1)a_n = 0$ ,
- (Closed-form)  $a_n = 2 \frac{3^n (2n)!}{n(n+2)!}$ ,
- (Asymptotics)  $a_n \sim 2 \frac{12^n}{\sqrt{\pi n^5}}$ , when  $n \rightarrow +\infty$ .

## Objectives

- **Introduce** so-called Discrete Differential Equations (DDEs),
- **Determine the nature** of the solutions of DDEs,
- Provide an **efficient algorithm** for computing a witness,
- Implementation in action  $\rightsquigarrow$  Solving a problem **previously out of reach**.

## Objectives

- **Introduce** so-called Discrete Differential Equations (DDEs),
- **Determine the nature** of the solutions of DDEs,
- Provide an **efficient algorithm** for computing a witness,
- Implementation in action  $\rightsquigarrow$  Solving a problem **previously out of reach**.

## Plan

- I Perform the above points for DDEs [Bostan, N., Safey El Din '23]
- II Perform the above points for *systems* of DDEs [N., Yurkevich '23]

## Objects of interest: Discrete Differential Equations

### Definition

Given  $f \in \mathbb{Q}[u]$ ,  $k \geq 1$ , and  $Q \in \mathbb{Q}[y_0, \dots, y_k, t, u]$ ,

$$F = f + t \cdot Q(F, \Delta F, \dots, \Delta^k F, t, u) \quad (\text{DDE})$$

is a *Discrete Differential Equation*, where  $\Delta : F \in \mathbb{Q}[u][[t]] \mapsto \frac{F(t,u) - F(t,1)}{u-1} \in \mathbb{Q}[u][[t]]$ , and where for  $\ell \geq 1$  we define  $\Delta^{\ell+1} = \Delta^\ell \circ \Delta$ .

## Objects of interest: Discrete Differential Equations

### Definition

Given  $f \in \mathbb{Q}[u]$ ,  $k \geq 1$ , and  $Q \in \mathbb{Q}[y_0, \dots, y_k, t, u]$ ,

$$F = f + t \cdot Q(F, \Delta F, \dots, \Delta^k F, t, u) \quad (\text{DDE})$$

is a *Discrete Differential Equation*, where  $\Delta : F \in \mathbb{Q}[u][[t]] \mapsto \frac{F(t,u) - F(t,1)}{u-1} \in \mathbb{Q}[u][[t]]$ , and where for  $\ell \geq 1$  we define  $\Delta^{\ell+1} = \Delta^\ell \circ \Delta$ .

Going back to our 3-constellations...

$$F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$$



# Objects of interest: Discrete Differential Equations

## Definition

Given  $f \in \mathbb{Q}[u]$ ,  $k \geq 1$ , and  $Q \in \mathbb{Q}[y_0, \dots, y_k, t, u]$ ,

$$F = f + t \cdot Q(F, \Delta F, \dots, \Delta^k F, t, u) \quad (\text{DDE})$$

is a *Discrete Differential Equation*, where  $\Delta : F \in \mathbb{Q}[u][[t]] \mapsto \frac{F(t,u) - F(t,1)}{u-1} \in \mathbb{Q}[u][[t]]$ , and where for  $\ell \geq 1$  we define  $\Delta^{\ell+1} = \Delta^\ell \circ \Delta$ .

Going back to our 3-constellations...

$$F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$$

## Theorem

[Bousquet-Mélou, Jehanne '06]

The unique solution in  $\mathbb{Q}[u][[t]]$  of (DDE) is **algebraic** over  $\mathbb{Q}(t, u)$ .

↪ Constructive proof  $\implies$  **algorithm**

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right),$

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0.$

• **Compute**  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$

• **Consider**

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

- Consider

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

- Show that there exist distinct  $U_1, U_2 \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]]$  s.t.  $\partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0$ ,

## Bousquet-Mélou and Jehanne's algorithm

$$\text{Input: } F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right),$$

$$\text{Output: } 81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0.$$

• Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Consider

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

• Show that there exist distinct  $U_1, U_2 \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]]$  s.t.  $\partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Set up

$$\text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_u P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ m \cdot (U_1 - U_2) - 1 = 0. \end{cases}$$

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

• Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Consider

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

• Show that there exist distinct  $U_1, U_2 \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]]$  s.t.  $\partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Set up

$$\text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_u P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ m \cdot (U_1 - U_2) - 1 = 0. \end{cases}$$

### Elimination theory

- Eliminate all series but  $F(t, 1)$

## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

• Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Consider

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

• Show that there exist distinct  $U_1, U_2 \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]]$  s.t.  $\partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Set up

$$\text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_u P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ m \cdot (U_1 - U_2) - 1 = 0. \end{cases}$$

### Elimination theory

- Eliminate all series but  $F(t, 1)$   
→ **Resultants**



## Bousquet-Mélou and Jehanne's algorithm

**Input:**  $F(t, u) = 1 + tu \left( F(t, u)^3 + (2F(t, u) + F(t, 1)) \frac{F(t, u) - F(t, 1)}{u-1} + \frac{F(t, u) - F(t, 1) - (u-1)\partial_u F(t, 1)}{(u-1)^2} \right)$ ,

**Output:**  $81t^2 F(t, 1)^3 - 9t(9t - 2)F(t, 1)^2 + (27t^2 - 66t + 1)F(t, 1) - 3t^2 + 47t - 1 = 0$ .

• Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Consider

$$\partial_u F(t, u) \cdot \partial_x P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) + \partial_u P(F(t, u), u, F(t, 1), \partial_u F(t, 1)) = 0,$$

• Show that there exist distinct  $U_1, U_2 \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]]$  s.t.  $\partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0$ ,

• Set up

$$\text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_x P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ \partial_u P(F(t, U_i), U_i, F(t, 1), \partial_u F(t, 1)) = 0, \\ m \cdot (U_1 - U_2) - 1 = 0. \end{cases}$$

### Elimination theory

- Eliminate all series but  $F(t, 1)$   
→ **Resultants**  
→ **Gröbner bases**

$$\mathcal{S} : \quad \text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_x P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_u P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \end{cases} \quad U_1 - U_2 \neq 0.$$

$$S : \quad \text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_x P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_u P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \end{cases} \quad U_1 - U_2 \neq 0.$$

### Assumptions

- $U_1, U_2$  are **distinct series**,
- $S$  has **finitely many solutions** in  $\overline{\mathbb{Q}(t)}$ ,
- $S$  generates a **radical ideal** over  $\mathbb{Q}(t)$ .

$$\mathcal{S} : \quad \text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_x P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_u P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \end{cases} \quad U_1 - U_2 \neq 0.$$

### Assumptions

- $U_1, U_2$  are **distinct series**,
- $\mathcal{S}$  has **finitely many solutions** in  $\overline{\mathbb{Q}(t)}$ ,
- $\mathcal{S}$  generates a **radical ideal** over  $\mathbb{Q}(t)$ .

### Useful properties

- $\mathfrak{S}_2$  acts on  $V(\mathcal{S})$  by **permuting**  $U_1, U_2$ ,
- $\#V(\mathcal{S}) \leq$  **Bézout bound** associated with  $\mathcal{S}$ ,
- Allows to forget  $U_1 - U_2 \neq 0$  in the Bézout bound.

$$S : \quad \text{For } 1 \leq i \leq 2, \begin{cases} P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_x P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \\ \partial_u P(F(t, U_i), F(t, 1), \partial_u F(t, 1), t, U_i) = 0, \end{cases} \quad U_1 - U_2 \neq 0.$$

## Assumptions

- $U_1, U_2$  are **distinct series**,
- $S$  has **finitely many solutions** in  $\overline{\mathbb{Q}(t)}^6$ ,
- $S$  generates a **radical ideal** over  $\mathbb{Q}(t)$ .

## Useful properties

- $\mathfrak{S}_2$  acts on  $V(S)$  by **permuting**  $U_1, U_2$ ,
- $\#V(S) \leq$  **Bézout bound** associated with  $S$ ,
- Allows to forget  $U_1 - U_2 \neq 0$  in the Bézout bound.

[Bostan, N., Safey El Din '23]

Under the above assumptions:

$\delta := \deg(P)$

- There exists some nonzero polynomial  $R \in \mathbb{Q}[z_0, t]$  whose partial degrees are upper bounded by  $\delta^2(\delta - 1)^4/2$ , such that  $R(F(t, 1), t) = 0$ .
- There exists an algorithm computing this  $R$  in  $O_{\log}(\delta^{17})$  ops. in  $\mathbb{Q}$ .

(We proved a **general version** of this result)

## Some preliminaries on Gröbner bases

$\mathcal{A} := \mathbb{Q}[x, \mathbf{y}]$  polynomial ring, where  $\mathbf{y} = y_1, \dots, y_s$ .

### Monomial orders

- $x^4 y_1^3 y_2^2 \succ_{lex} x^3 y_1^4 y_2^2$  for a **lexicographic order**,
- $x^4 y_1^2 y_2^3 \succ_{bmon} x^4 y_1^3 y_2$  for a **block monomial order**.

## Some preliminaries on Gröbner bases

$\mathcal{A} := \mathbb{Q}[x, \mathbf{y}]$  polynomial ring, where  $\mathbf{y} = y_1, \dots, y_s$ .

### Monomial orders

- $x^4 y_1^3 y_2^2 \succ_{lex} x^3 y_1^4 y_2^2$  for a **lexicographic order**,
- $x^4 y_1^2 y_2^3 \succ_{bmon} x^4 y_1^3 y_2$  for a **block monomial order**.

### Leading terms for some order $\succ$

For  $Q \in \mathcal{A}$ , the leading term  $\text{LT}_\succ(Q)$  of  $Q$  is the monomial of **highest weight** for  $\succ$ .

## Some preliminaries on Gröbner bases

$\mathcal{A} := \mathbb{Q}[x, \mathbf{y}]$  polynomial ring, where  $\mathbf{y} = y_1, \dots, y_s$ .

### Monomial orders

- $x^4 y_1^3 y_2^2 \succ_{lex} x^3 y_1^4 y_2^2$  for a **lexicographic order**,
- $x^4 y_1^2 y_2^3 \succ_{bmon} x^4 y_1^3 y_2$  for a **block monomial order**.

### Leading terms for some order $\succ$

For  $Q \in \mathcal{A}$ , the leading term  $\text{LT}_\succ(Q)$  of  $Q$  is the monomial of **highest weight** for  $\succ$ .

### Definition

Fix a monomial order  $\succ$  on  $\mathcal{A}$ . A finite subset  $G = \{g_1, \dots, g_t\}$  of an ideal  $\mathcal{I} \subset \mathcal{A}$  different from 0 is said to be a **Gröbner basis** if  $\langle \text{LT}_\succ(g_1), \dots, \text{LT}_\succ(g_t) \rangle = \langle \text{LT}_\succ(\mathcal{I}) \rangle$ .



# Some preliminaries on Gröbner bases

$\mathcal{A} := \mathbb{Q}[x, \mathbf{y}]$  polynomial ring, where  $\mathbf{y} = y_1, \dots, y_s$ .

## Monomial orders

- $x^4 y_1^3 y_2^2 \succ_{lex} x^3 y_1^4 y_2^2$  for a **lexicographic order**,
- $x^4 y_1^2 y_2^3 \succ_{bmon} x^4 y_1^3 y_2$  for a **block monomial order**.

## Leading terms for some order $\succ$

For  $Q \in \mathcal{A}$ , the leading term  $\text{LT}_\succ(Q)$  of  $Q$  is the monomial of **highest weight** for  $\succ$ .

## Definition

Fix a monomial order  $\succ$  on  $\mathcal{A}$ . A finite subset  $G = \{g_1, \dots, g_t\}$  of an ideal  $\mathcal{I} \subset \mathcal{A}$  different from 0 is said to be a **Gröbner basis** if  $\langle \text{LT}_\succ(g_1), \dots, \text{LT}_\succ(g_t) \rangle = \langle \text{LT}_\succ(\mathcal{I}) \rangle$ .

## Properties

- Such bases always **exist** and **generate**  $\mathcal{I}$ ,
- Computing Gröbner bases is **NP-hard**,
- Gröbner bases are a **powerful tool** in elimination theory.

## New geometric modelling of the problem with A. Bostan and M. Safey El Din

There exist 2 solutions  $(x, \mathbf{u}) \in \overline{\mathbb{Q}(t)}^2$  with **distinct**  $\mathbf{u}$ -coordinates to

$$\begin{cases} P(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = 0, \\ \partial_x P(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = 0, \quad \mathbf{u} \neq \mathbf{0}, \\ \partial_{\mathbf{u}} P(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = 0. \end{cases}$$

## New geometric modelling of the problem with A. Bostan and M. Safey El Din

There exist 2 solutions  $(x, \mathbf{u}) \in \overline{\mathbb{Q}(t)}^2$  with **distinct**  $\mathbf{u}$ -coordinates to

$$\begin{cases} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \\ \partial_x \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \quad \mathbf{u} \neq \mathbf{0}, \\ \partial_{\mathbf{u}} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}. \end{cases}$$

$$\pi_x : (x, \mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^4 \mapsto (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3,$$

$$\mathbf{W} := \pi_x(V(\mathbf{P}, \partial_x \mathbf{P}, \partial_{\mathbf{u}} \mathbf{P}) \setminus V(\mathbf{u}))$$

$$\pi_u : (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3 \mapsto (z_0, z_1) \in \overline{\mathbb{Q}(t)}^2,$$

## New geometric modelling of the problem with A. Bostan and M. Safey El Din

There exist 2 solutions  $(x, \mathbf{u}) \in \overline{\mathbb{Q}(t)}^2$  with **distinct**  $\mathbf{u}$ -coordinates to

$$\begin{cases} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \\ \partial_x \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \quad \mathbf{u} \neq \mathbf{0}, \\ \partial_{\mathbf{u}} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}. \end{cases}$$

$$\pi_x : (x, \mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^4 \mapsto (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3,$$

$$\mathbf{W} := \pi_x(V(\mathbf{P}, \partial_x \mathbf{P}, \partial_{\mathbf{u}} \mathbf{P}) \setminus V(\mathbf{u}))$$

$$\pi_u : (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3 \mapsto (z_0, z_1) \in \overline{\mathbb{Q}(t)}^2,$$

**Characterize** with polynomial constraints

$$\mathcal{F}_2 := \{\alpha_z \in \overline{\mathbb{Q}(t)}^2 \mid \# \pi_u^{-1}(\alpha_z) \cap \mathbf{W} \geq 2\}$$

# New geometric modelling of the problem with A. Bostan and M. Safey El Din

There exist 2 solutions  $(x, \mathbf{u}) \in \overline{\mathbb{Q}(t)}^2$  with **distinct**  $\mathbf{u}$ -coordinates to

$$\begin{cases} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \\ \partial_x \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}, \quad \mathbf{u} \neq \mathbf{0}, \\ \partial_{\mathbf{u}} \mathbf{P}(x, \mathbf{u}, \mathbf{F}(t, \mathbf{0}), \partial_{\mathbf{u}}\mathbf{F}(t, \mathbf{0})) = \mathbf{0}. \end{cases}$$

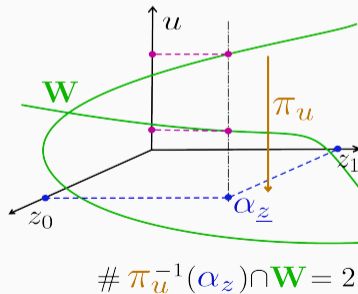
$$\pi_x : (x, \mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^4 \mapsto (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3,$$

$$\mathbf{W} := \pi_x(V(\mathbf{P}, \partial_x \mathbf{P}, \partial_{\mathbf{u}} \mathbf{P}) \setminus V(\mathbf{u}))$$

$$\pi_u : (\mathbf{u}, z_0, z_1) \in \overline{\mathbb{Q}(t)}^3 \mapsto (z_0, z_1) \in \overline{\mathbb{Q}(t)}^2,$$

**Characterize** with polynomial constraints

$$\mathcal{F}_2 := \{\alpha_z \in \overline{\mathbb{Q}(t)}^2 \mid \# \pi_u^{-1}(\alpha_z) \cap \mathbf{W} \geq 2\}$$



**Input:**  $F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right),$

$k = 2$

**Output:**  $t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$

$$\text{Input: } F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right), \quad \mathbf{k} = 2$$

$$\text{Output: } t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$$

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 0), \partial_u F(t, 0)) = 0$ ,

$$\text{Input: } F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right), \quad \mathbf{k} = 2$$

$$\text{Output: } t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$$

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 0), \partial_u F(t, 0)) = 0$ ,
- Compute  $G_u$  **Gröbner basis** of  $\langle P, \partial_1 P, \partial_2 P, m \cdot u - 1 \rangle \cap \mathbb{Q}(t)[u, z_0, z_1]$  for  $\{u\} \succ_{lex} \{z_0, z_1\}$ :



$$\text{Input: } F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right),$$

$$k = 2$$

$$\text{Output: } t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$$

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 0), \partial_u F(t, 0)) = 0$ ,
- Compute  $G_u$  **Gröbner basis** of  $\langle P, \partial_1 P, \partial_2 P, m \cdot u - 1 \rangle \cap \mathbb{Q}(t)[u, z_0, z_1]$  for  $\{u\} \succ_{\text{lex}} \{z_0, z_1\}$ :

 $B_0 :$ 

$$\gamma_0$$

$$B_1 : \begin{cases} \beta_1 \cdot u + \gamma_1 \\ \vdots \\ \beta_r \cdot u + \gamma_r \end{cases}, \gamma_i, \beta_j \in \mathbb{Q}(t)[z_0, z_1]$$

$$B_2 : g_2 := u^2 + \beta_{r+1} \cdot u + \gamma_{r+1}$$

“At  $\alpha \in \pi_u(V(G_u)) \subset \overline{\mathbb{Q}(t)}^2$ ,  
there exist two **distinct** solutions in  $u$ ”

$$\text{Input: } F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right),$$

$$k = 2$$

$$\text{Output: } t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$$

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 0), \partial_u F(t, 0)) = 0$ ,
- Compute  $G_u$  **Gröbner basis** of  $\langle P, \partial_1 P, \partial_2 P, m \cdot u - 1 \rangle \cap \mathbb{Q}(t)[u, z_0, z_1]$  for  $\{u\} \succ_{\text{lex}} \{z_0, z_1\}$ :

 $B_0 :$ 

$$\gamma_0$$

$$B_1 : \begin{cases} \beta_1 \cdot u + \gamma_1 \\ \vdots \\ \beta_r \cdot u + \gamma_r \end{cases}, \gamma_i, \beta_j \in \mathbb{Q}(t)[z_0, z_1]$$

$B_2 : g_2 := u^2 + \beta_{r+1} \cdot u + \gamma_{r+1}$

“At  $\alpha \in \pi_u(V(G_u)) \subset \overline{\mathbb{Q}(t)}^2$ ,  
there exist two **distinct** solutions in  $u$ ”

At  $\alpha \in V(G_u \cap \mathbb{K}[t, z_0, z_1])$  fixed,  
there exist two solutions in  $u$

$$\implies \beta_i, \gamma_j = 0 \quad (\text{equations})$$

$$\text{Input: } F(t, u) = 1 + t \left( uF(t, u) + \frac{F(t, u) - F(t, 0) - u\partial_u F(t, 0)}{u^2} \right),$$

$$k = 2$$

$$\text{Output: } t^3 F(t, 0)^3 - F(t, 0) + 1 = 0.$$

- Compute  $P \in \mathbb{Q}(t)[x, u, z_0, z_1]$  such that  $P(F(t, u), u, F(t, 0), \partial_u F(t, 0)) = 0$ ,
- Compute  $G_u$  **Gröbner basis** of  $\langle P, \partial_1 P, \partial_2 P, m \cdot u - 1 \rangle \cap \mathbb{Q}(t)[u, z_0, z_1]$  for  $\{u\} \succ_{\text{lex}} \{z_0, z_1\}$ :

$$B_0 : \quad \gamma_0$$

$$B_1 : \begin{cases} \beta_1 \cdot u + \gamma_1 \\ \vdots \\ \beta_r \cdot u + \gamma_r \end{cases}, \gamma_i, \beta_j \in \mathbb{Q}(t)[z_0, z_1]$$

“At  $\alpha \in \pi_u(V(G_u)) \subset \overline{\mathbb{Q}(t)}^2$ ,  
there exist two **distinct** solutions in  $u$ ”

$$B_2 : \quad g_2 := u^2 + \beta_{r+1} \cdot u + \gamma_{r+1}$$

At  $\alpha \in V(G_u \cap \mathbb{K}[t, z_0, z_1])$  fixed,  
there exist two solutions in  $u$   
 $\implies \beta_i, \gamma_j = 0$  (**equations**)

### [Extension theorem]

$\alpha \in \pi_u(V(G_u)) \implies \text{LeadingCoeff}_u(g_2) \neq 0$   
Distinct solutions in  $u \implies \text{disc}_u(g_2) \neq 0$  (**inequations**)

Projecting  $\implies$  Elimination theorem

Lifting points of the projections  $\implies$  Extension theorem

Projecting  $\implies$  Elimination theorem

Lifting points of the projections  $\implies$  Extension theorem

**[Proposition]** Let  $g \in (\mathbb{Q}(t)[z_0, z_1])[u]$ . Then  $g$  has at least  $i$  distinct solutions at  $\alpha \in \overline{\mathbb{Q}(t)}^2$  if and only if the  $(i \times i)$ -minors of the **Hermite quadratic form** associated with  $g$  **do not vanish simultaneously** at  $\alpha$ .

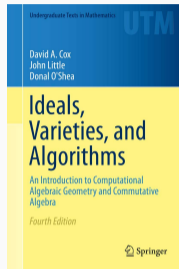
$\rightsquigarrow$  Reduces to studying the **multiplication maps**  $(M_{u^\ell} : q \mapsto q \cdot u^\ell)_{\ell \geq 1}$  in  $(\mathbb{Q}[t, z_0, z_1])[u]/\langle g \rangle$

Projecting  $\implies$  Elimination theorem

Lifting points of the projections  $\implies$  Extension theorem

**[Proposition]** Let  $g \in (\mathbb{Q}(t)[z_0, z_1])[u]$ . Then  $g$  has at least  $i$  distinct solutions at  $\alpha \in \overline{\mathbb{Q}(t)}^2$  if and only if the  $(i \times i)$ -minors of the **Hermite quadratic form** associated with  $g$  **do not vanish simultaneously** at  $\alpha$ .

$\rightsquigarrow$  Reduces to studying the **multiplication maps**  $(M_{u^\ell} : q \mapsto q \cdot u^\ell)_{\ell \geq 1}$  in  $(\mathbb{Q}[t, z_0, z_1])[u]/\langle g \rangle$



Projecting  $\implies$  Elimination theorem

Lifting points of the projections  $\implies$  Extension theorem

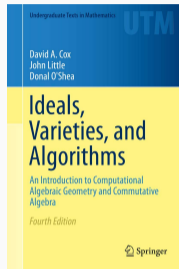
**[Proposition]** Let  $g \in (\mathbb{Q}(t)[z_0, z_1])[u]$ . Then  $g$  has at least  $i$  distinct solutions at  $\alpha \in \overline{\mathbb{Q}(t)}^2$  if and only if the  $(i \times i)$ -minors of the **Hermite quadratic form** associated with  $g$  **do not vanish simultaneously** at  $\alpha$ .

$\rightsquigarrow$  Reduces to studying the **multiplication maps**  $(M_{u^\ell} : q \mapsto q \cdot u^\ell)_{\ell \geq 1}$  in  $(\mathbb{Q}[t, z_0, z_1])[u]/\langle g \rangle$

**[Bostan, N., Safey El Din '23]**

Disjunction of conjunctions of polynomial equations and inequations whose zero set is  $\mathcal{F}_2$

(Our strategy works in the **general case**)



Projecting  $\implies$  Elimination theorem

Lifting points of the projections  $\implies$  Extension theorem

**[Proposition]** Let  $g \in (\mathbb{Q}(t)[z_0, z_1])[u]$ . Then  $g$  has at least  $i$  distinct solutions at  $\alpha \in \overline{\mathbb{Q}(t)}^2$  if and only if the  $(i \times i)$ -minors of the **Hermite quadratic form** associated with  $g$  do not vanish simultaneously at  $\alpha$ .

$\rightsquigarrow$  Reduces to studying the multiplication maps  $(M_{u^\ell} : q \mapsto q \cdot u^\ell)_{\ell \geq 1}$  in  $(\mathbb{Q}[t, z_0, z_1])[u]/\langle g \rangle$

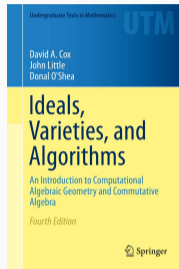
**[Bostan, N., Safey El Din '23]**

Disjunction of conjunctions of polynomial equations and inequations whose zero set is  $\mathcal{F}_2$

(Our strategy works in the **general case**)

**[5-constellations  $k = 4$ ]**

Strategy	Timing	$(d_{z_0}, d_t)$
Duplication	$> 5d$	?
Elimination	<b>2d21h</b>	<b>(9, 3)</b>





# Systems of Discrete Differential Equations

What could be extended to **systems**?

What could be extended to **systems**?

Modelling **special Eulerian planar orientations**:

$$\begin{cases} F(t, u) = 1 + t \cdot \left( u + 2uF(t, u)^2 + 2uG(t, 1) + u \frac{F(t, u) - uF(t, 1)}{u-1} \right), \\ G(t, u) = t \cdot \left( 2uF(t, u)G(t, u) + uF(t, u) + uG(t, 1) + u \frac{G(t, u) - uG(t, 1)}{u-1} \right). \end{cases}$$

[Bonichon, Bousquet-Mélou, Dorbec, Pennarun '17]

What could be extended to **systems**?

Modelling **special Eulerian planar orientations**:

$$\begin{cases} F(t, u) = 1 + t \cdot \left( u + 2uF(t, u)^2 + 2uG(t, 1) + u \frac{F(t, u) - uF(t, 1)}{u-1} \right), \\ G(t, u) = t \cdot \left( 2uF(t, u)G(t, u) + uF(t, u) + uG(t, 1) + u \frac{G(t, u) - uG(t, 1)}{u-1} \right). \end{cases}$$

[Bonichon, Bousquet-Mélou, Dorbec, Pennarun '17]

Modelling **hard particles on planar maps**:

$$\begin{cases} F(t, u) = x - y + G(t, u) + tu \left( uF(t, u)^2 + \frac{uF(t, u) - F(t, 1)}{u-1} \right) \\ G(t, u) = y + tsu \left( F(t, u)G(t, u) + \frac{G(t, u) - G(t, 1)}{u-1} \right) \end{cases}$$

[Bousquet-Mélou, Jehanne '06]

[Popescu '86, Swan '98]

(1.4) THEOREM. *Let  $k$  be a field,  $k\langle X \rangle$  the algebraic power series ring in  $X = (X_1, \dots, X_r)$  over  $k$ ,  $f$  a finite system of polynomial equations over  $k\langle X \rangle$  and  $\hat{y} = (\hat{y}_1, \dots, \hat{y}_n) \in k[[X]]^n$  a formal solution of  $f$  such that  $\hat{y}_i \in k[[X_1, \dots, X_{s_i}]]$ ,  $1 \leq i \leq n$  for some positive integers  $s_i \leq r$ . Then there exists a solution  $y = (y_1, \dots, y_n)$  of  $f$  in  $k\langle X \rangle$  such that  $y_i \in k\langle X_1, \dots, X_{s_i} \rangle$ ,  $1 \leq i \leq n$ .*

## Popescu's theorem yielding **algebraicity** of the solutions

[Popescu '86, Swan '98]

(1.4) THEOREM. *Let  $k$  be a field,  $k\langle X \rangle$  the algebraic power series ring in  $X = (X_1, \dots, X_r)$  over  $k$ ,  $f$  a finite system of polynomial equations over  $k\langle X \rangle$  and  $\hat{y} = (\hat{y}_1, \dots, \hat{y}_n) \in k[[X]]^n$  a formal solution of  $f$  such that  $\hat{y}_i \in k[[X_1, \dots, X_{s_i}]]$ ,  $1 \leq i \leq n$  for some positive integers  $s_i \leq r$ . Then there exists a solution  $y = (y_1, \dots, y_n)$  of  $f$  in  $k\langle X \rangle$  such that  $y_i \in k\langle X_1, \dots, X_{s_i} \rangle$ ,  $1 \leq i \leq n$ .*

- Solutions of systems of DDEs are **unique** with components in  $\mathbb{Q}[\mathbf{u}][[t]] \implies$  they are **algebraic**!

## Popescu's theorem yielding algebraicity of the solutions

[Popescu '86, Swan '98]

(1.4) THEOREM. Let  $k$  be a field,  $k\langle X \rangle$  the algebraic power series ring in  $X = (X_1, \dots, X_r)$  over  $k$ ,  $f$  a finite system of polynomial equations over  $k\langle X \rangle$  and  $\hat{y} = (\hat{y}_1, \dots, \hat{y}_n) \in k[[X]]^n$  a formal solution of  $f$  such that  $\hat{y}_i \in k[[X_1, \dots, X_{s_i}]]$ ,  $1 \leq i \leq n$  for some positive integers  $s_i \leq r$ . Then there exists a solution  $y = (y_1, \dots, y_n)$  of  $f$  in  $k\langle X \rangle$  such that  $y_i \in k\langle X_1, \dots, X_{s_i} \rangle$ ,  $1 \leq i \leq n$ .

- Solutions of systems of DDEs are **unique** with components in  $\mathbb{Q}[u][[t]] \implies$  they are **algebraic!**

[planar maps]

$$H(t, u) = 1 + t \left( u^2 H(t, u)^2 + u \frac{uH(t, u) - G(t, u)}{u-1} \right)$$

- There exists a solution  $(H, G) = (F, F(t, 1))$ , where  $F \in \mathbb{Q}[u][[t]]$ ,
- The involved series are  $F(t, 1)$  and  $F(t, u)$ , and  $\{t\} \subset \{t, u\}$ .

## Popescu's theorem yielding **algebraicity** of the solutions

[Popescu '86, Swan '98]

(1.4) THEOREM. Let  $k$  be a field,  $k\langle X \rangle$  the algebraic power series ring in  $X = (X_1, \dots, X_r)$  over  $k$ ,  $f$  a finite system of polynomial equations over  $k\langle X \rangle$  and  $\hat{y} = (\hat{y}_1, \dots, \hat{y}_n) \in k[[X]]^n$  a formal solution of  $f$  such that  $\hat{y}_i \in k[[X_1, \dots, X_{s_i}]]$ ,  $1 \leq i \leq n$  for some positive integers  $s_i \leq r$ . Then there exists a solution  $y = (y_1, \dots, y_n)$  of  $f$  in  $k\langle X \rangle$  such that  $y_i \in k\langle X_1, \dots, X_{s_i} \rangle$ ,  $1 \leq i \leq n$ .

- Solutions of systems of DDEs are **unique** with components in  $\mathbb{Q}[\mathbf{u}][[t]] \implies$  they are **algebraic**!

[planar maps]

$$H(t, u) = 1 + t \left( u^2 H(t, u)^2 + u \frac{uH(t, u) - G(t, u)}{u-1} \right)$$

- There exists a solution  $(H, G) = (F, F(t, 1))$ , where  $F \in \mathbb{Q}[\mathbf{u}][[t]]$ ,
- The involved series are  $F(t, 1)$  and  $F(t, u)$ , and  $\{t\} \subset \{t, u\}$ .

The proof is **highly not constructive**... How to compute witnesses?



## Constructive algebraicity theorem for solutions of systems of DDEs (FPSAC'23)

[N., Yurkevich '23]

Let  $n, k \geq 1$  be integers and  $f_1, \dots, f_n \in \mathbb{Q}[u]$ ,  $Q_1, \dots, Q_n \in \mathbb{Q}[y_1, \dots, y_{n(k+1)}, t, u]$  be polynomials. Denote  $\nabla^k F := F, \Delta F, \dots, \Delta^k F$ . Then the **system** of DDEs

$$\begin{cases} (\mathbf{E}_{F_1}): F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots \\ (\mathbf{E}_{F_n}): F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u). \end{cases} \quad (\text{SDDEs})$$

admits a **unique** vector of solutions  $(F_1, \dots, F_n) \in \mathbb{Q}[u][[t]]^n$ , and all its components are **algebraic** over  $\mathbb{Q}(t, u)$ .



# Constructive algebraicity theorem for solutions of systems of DDEs (FPSAC'23)

[N., Yurkevich '23]

Let  $n, k \geq 1$  be integers and  $f_1, \dots, f_n \in \mathbb{Q}[u]$ ,  $Q_1, \dots, Q_n \in \mathbb{Q}[y_1, \dots, y_{n(k+1)}, t, u]$  be polynomials. Denote  $\nabla^k F := F, \Delta F, \dots, \Delta^k F$ . Then the **system** of DDEs

$$\begin{cases} (\mathbf{E}_{F_1}): F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots \\ (\mathbf{E}_{F_n}): F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u). \end{cases} \quad (\text{SDDEs})$$

admits a **unique** vector of solutions  $(F_1, \dots, F_n) \in \mathbb{Q}[u][[t]]^n$ , and all its components are **algebraic** over  $\mathbb{Q}(t, u)$ .

[Proof sketch]

- There exists a **polynomial system**  $\mathcal{S}$  defined over  $\mathbb{Q}(t)$  in  $nk(n+2)$  equations and unknowns, that admits a solution  $\mathcal{P}$  with  $F_1(t, 1)$  as one of its coordinates,
- The **Jacobian** of  $\mathcal{S}$  is **invertible** at  $\mathcal{P} \implies F_1(t, 1)$  is **algebraic** over  $\mathbb{Q}(t)$ .

## Identifying more polynomial equations

Consider

$$\rightsquigarrow F_1, F_2 \equiv F_1(t, u), F_2(t, u) \in \mathbb{Q}[u][[t]]$$

$$\begin{cases} 0 = (1 - F_1) \cdot (u - 1) + tu \cdot (2uF_1^2 - uF_1(t, 1) + 2uF_2(t, 1) - 2F_1^2 + u + F_1 - 2F_2(t, 1) - 1), \\ 0 = F_2 \cdot (1 - u) + tu \cdot (2uF_1F_2 + uF_1 - 2F_1F_2 - F_1 + F_2 - F_2(t, 1)). \end{cases}$$

Denote by  $E_1, E_2 \in \mathbb{Q}(t)[x_1, x_2, u, z_0, z_1]$  polynomials such that

$$\text{for } i \in \{1, 2\}, \quad E_i(F_1(t, u), F_2(t, u), u, F_1(t, 1), F_2(t, 1)) = 0. \quad (\equiv E_i(u))$$

## Identifying more polynomial equations

Consider

$$\rightsquigarrow F_1, F_2 \equiv F_1(t, u), F_2(t, u) \in \mathbb{Q}[u][[t]]$$

$$\begin{cases} 0 = (1 - F_1) \cdot (u - 1) + tu \cdot (2uF_1^2 - uF_1(t, 1) + 2uF_2(t, 1) - 2F_1^2 + u + F_1 - 2F_2(t, 1) - 1), \\ 0 = F_2 \cdot (1 - u) + tu \cdot (2uF_1F_2 + uF_1 - 2F_1F_2 - F_1 + F_2 - F_2(t, 1)). \end{cases}$$

Denote by  $E_1, E_2 \in \mathbb{Q}(t)[x_1, x_2, u, z_0, z_1]$  polynomials such that

$$\text{for } i \in \{1, 2\}, \quad E_i(F_1(t, u), F_2(t, u), u, F_1(t, 1), F_2(t, 1)) = 0. \quad (\equiv E_i(u))$$

**Differentiating** with respect to  $u$  yields

$$\begin{pmatrix} (\partial_{x_1} E_1)(u) & (\partial_{x_2} E_1)(u) \\ (\partial_{x_1} E_2)(u) & (\partial_{x_2} E_2)(u) \end{pmatrix} \cdot \begin{pmatrix} \partial_u F_1 \\ \partial_u F_2 \end{pmatrix} + \begin{pmatrix} (\partial_u E_1)(u) \\ (\partial_u E_2)(u) \end{pmatrix} = 0.$$

$$\text{For } \mathbf{U}(t) \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]], \begin{cases} \text{if } (\partial_{x_1} E_1 \cdot \partial_{x_2} E_2 - \partial_{x_1} E_2 \cdot \partial_{x_2} E_1)(\mathbf{U}(t)) = 0, \\ \text{then } (\partial_{x_1} E_1 \cdot \partial_u E_2 - \partial_{x_1} E_2 \cdot \partial_u E_1)(\mathbf{U}(t)) = 0. \end{cases}$$

## Identifying more polynomial equations

Consider

$$\rightsquigarrow F_1, F_2 \equiv F_1(t, u), F_2(t, u) \in \mathbb{Q}[u][[t]]$$

$$\begin{cases} 0 = (1 - F_1) \cdot (\mathbf{u} - 1) + t\mathbf{u} \cdot (2\mathbf{u}F_1^2 - \mathbf{u}F_1(t, 1) + 2\mathbf{u}F_2(t, 1) - 2F_1^2 + \mathbf{u} + F_1 - 2F_2(t, 1) - 1), \\ 0 = F_2 \cdot (1 - \mathbf{u}) + t\mathbf{u} \cdot (2\mathbf{u}F_1F_2 + \mathbf{u}F_1 - 2F_1F_2 - F_1 + F_2 - F_2(t, 1)). \end{cases}$$

Denote by  $E_1, E_2 \in \mathbb{Q}(t)[x_1, x_2, \mathbf{u}, z_0, z_1]$  polynomials such that

$$\text{for } i \in \{1, 2\}, \quad E_i(F_1(t, \mathbf{u}), F_2(t, \mathbf{u}), \mathbf{u}, F_1(t, 1), F_2(t, 1)) = 0. \quad (\equiv E_i(\mathbf{u}))$$

**Differentiating** with respect to  $\mathbf{u}$  yields

$$\begin{pmatrix} (\partial_{x_1} E_1)(\mathbf{u}) & (\partial_{x_2} E_1)(\mathbf{u}) \\ (\partial_{x_1} E_2)(\mathbf{u}) & (\partial_{x_2} E_2)(\mathbf{u}) \end{pmatrix} \cdot \begin{pmatrix} \partial_{\mathbf{u}} F_1 \\ \partial_{\mathbf{u}} F_2 \end{pmatrix} + \begin{pmatrix} (\partial_{\mathbf{u}} E_1)(\mathbf{u}) \\ (\partial_{\mathbf{u}} E_2)(\mathbf{u}) \end{pmatrix} = 0.$$

$$\text{For } \mathbf{U}(t) \in \bigcup_{d \geq 1} \overline{\mathbb{Q}}[[t^{\frac{1}{d}}]], \quad \begin{cases} \text{if } (\partial_{x_1} E_1 \cdot \partial_{x_2} E_2 - \partial_{x_1} E_2 \cdot \partial_{x_2} E_1)(\mathbf{U}(t)) = 0, \\ \text{then } (\partial_{x_1} E_1 \cdot \partial_{\mathbf{u}} E_2 - \partial_{x_1} E_2 \cdot \partial_{\mathbf{u}} E_1)(\mathbf{U}(t)) = 0. \end{cases}$$

Does this yield an **elimination procedure**?

## A polynomial system for systems of 2 DDEs of order 1

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(F_1, \Delta F_1, F_2, \Delta F_2, t, u), \\ F_2 = f_2(u) + t \cdot Q_2(F_1, \Delta F_1, F_2, \Delta F_2, t, u). \end{cases} \quad (\text{SDDEs})$$

## A polynomial system for systems of 2 DDEs of order 1

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(F_1, \Delta F_1, F_2, \Delta F_2, t, u), \\ F_2 = f_2(u) + t \cdot Q_2(F_1, \Delta F_1, F_2, \Delta F_2, t, u). \end{cases} \quad (\text{SDDEs})$$

**Define** the “numerators”  $E_1, E_2$  and the polynomials

$$\text{Det} := \det \begin{pmatrix} \partial_{x_1} E_1 & \partial_{x_2} E_1 \\ \partial_{x_1} E_2 & \partial_{x_2} E_2 \end{pmatrix} \quad \text{and} \quad P := \det \begin{pmatrix} \partial_{x_1} E_1 & \partial_u E_1 \\ \partial_{x_1} E_2 & \partial_u E_2 \end{pmatrix}.$$

**Set up** the duplicated polynomial system  $\mathcal{S}$ , consisting in the 2 duplications of the polynomials  $(E_1, E_2, \text{Det}, P)$ : it has 8 equations and unknowns.

Moreover, one of its solutions in  $\overline{\mathbb{Q}(t)}^8$  is

$$\mathcal{P} := (F_1(t, U_1), F_2(t, U_1), F_1(t, U_2), F_2(t, U_2), U_1, U_2, F_1(t, 1), F_2(t, 1)).$$

**Compute** an element of  $\langle \mathcal{S}, m \cdot (U_1 - U_2) - 1 \rangle \cap \mathbb{Q}[z_0, t]$ .

## A polynomial system for systems of 2 DDEs of order 1

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(F_1, \Delta F_1, F_2, \Delta F_2, t, u), \\ F_2 = f_2(u) + t \cdot Q_2(F_1, \Delta F_1, F_2, \Delta F_2, t, u). \end{cases} \quad (\text{SDDEs})$$

**Define** the “numerators”  $E_1, E_2$  and the polynomials

$$\text{Det} := \det \begin{pmatrix} \partial_{x_1} E_1 & \partial_{x_2} E_1 \\ \partial_{x_1} E_2 & \partial_{x_2} E_2 \end{pmatrix} \quad \text{and} \quad P := \det \begin{pmatrix} \partial_{x_1} E_1 & \partial_u E_1 \\ \partial_{x_1} E_2 & \partial_u E_2 \end{pmatrix}.$$

**Set up** the duplicated polynomial system  $\mathcal{S}$ , consisting in the 2 duplications of the polynomials  $(E_1, E_2, \text{Det}, P)$ : it has 8 equations and unknowns.

Moreover, one of its solutions in  $\overline{\mathbb{Q}(t)}^8$  is

$$\mathcal{P} := (F_1(t, U_1), F_2(t, U_1), F_1(t, U_2), F_2(t, U_2), U_1, U_2, F_1(t, 1), F_2(t, 1)).$$

**Compute** an element of  $\langle \mathcal{S}, m \cdot (U_1 - U_2) - 1 \rangle \cap \mathbb{Q}[z_0, t]$ .

The **geometric strategy** previously described for solving one DDE can be extended here!

## Conclusion and perspectives

- **Decidability**: geometry-driven algorithm computing  $R \in \mathbb{Q}[z, t] \setminus \{0\}$  s.t.  $R(F_1(t, 1), t) = 0$ ,
- **Resolution** of the DDE of 5-constellations in an **automatic fashion**,
- **Constructive proof** of algebraicity of solutions of SDDEs.



## Conclusion and perspectives

- **Decidability**: geometry-driven algorithm computing  $R \in \mathbb{Q}[z, t] \setminus \{0\}$  s.t.  $R(F_1(t, 1), t) = 0$ ,
- **Resolution** of the DDE of 5-constellations in an **automatic fashion**,
- **Constructive proof** of algebraicity of solutions of SDDEs.

- **Implementing** the algorithm in a *Maple* package?  
(Work in progress)
- **Expanded algorithmic comparison** in the system case?  
(Work in progress with [S. Yurkevich](#))
- More **nested** catalytic variables in the direction of Popescu's theorem?  
(Work in progress with [M. Bousquet-Mélou](#))

# RTCA: Computer Algebra for Functional Equations in Combinatorics & Physics

September 18 to December 11, 2023

Organizers:

**Alin Bostan** (Inria, Saclay), **Mark Giesbrecht** (University of Waterloo)  
**Christoph Koutschan** (RICAM, Linz), **Marni Mishna** (SFU, Burnaby)  
**Mohab Safey El Din** (Sorbonne Université), **Bruno Salvy** (Inria, Lyon)  
**Gilles Villard** (CNRS, Lyon)



## Recent Trends in Computer Algebra

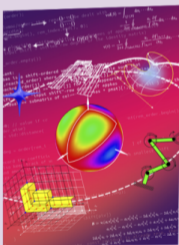
Thematic program with courses, workshops and topical days

**CIRM Preparatory School**  
March 6-10, 2023

**Fundamental Algorithms and  
Algorithmic Complexity**  
Special week: September 18-22  
**Workshop: September 25-29**

**Geometry of Polynomial System  
Solving, Optimization and Topology**  
Special week: October 9-13  
**Workshop: October 16-20**  
Topical days: October 23-24

**Computer Algebra for Functional  
Equations in Combinatorics & Physics**  
Special week: Nov. 27-Dec. 1  
**Workshop: December 4-8**  
Topical day: December 11



Program coordinated by the Centre Emile Borel (CEB) at IHP (Paris) and also accessible online  
Participation of postdocs and PhD students is strongly encouraged  
Registration is free however mandatory

Scientific program and registration on: <https://indico.math.cnrs.fr/category/588>  
Deadline for financial support: March 15<sup>th</sup>, 2023  
Contact: RTCA2023-Paris@ihp.fr

CEB organization assistant: Sofia Minasian  
CEB manager: Sylvie Lhermitte  
[www.ihp.fr](http://www.ihp.fr)



**November 27th → December 11th**

- **Special week:**  
November 27th → December 1st,
- **Workshop:**  
December 4th → December 8th,
- **Topical day:**  
December 11th.

**Thank you for your attention!**

## Solving 5-constellations using a Hybrid Guess-and-Prove strategy

**Input:** (The rather big DDE associated with the enumeration of 5-constellations)

**Output:**  $15625t^2F(t, 1)^5 - 31250t^2F(t, 1)^4 + (25000t^2 - 1000t)F(t, 1)^3 - (10000t^2 - 8700t)F(t, 1)^2 + (2000t^2 - 15855t + 16)F(t, 1) - 160t^2 + 8139t - 16 = 0$

## Solving 5-constellations using a Hybrid Guess-and-Prove strategy

**Input:** (The rather big DDE associated with the enumeration of 5-constellations)

**Output:**  $15625t^2F(t, 1)^5 - 31250t^2F(t, 1)^4 + (25000t^2 - 1000t)F(t, 1)^3 - (10000t^2 - 8700t)F(t, 1)^2 + (2000t^2 - 15855t + 16)F(t, 1) - 160t^2 + 8139t - 16 = 0$

- **Draw at random** a prime number  $p$  and some  $c \in \mathbb{F}_p$ ,
- **Compute** upper bounds  $(9, 3)$  on the bidegree of  $M \in \mathbb{F}_p[z, t]$  annihilating  $F(t, 1)$  modulo  $p$ ,
- **Expand** the truncated series  $F(t, 1) \bmod t^{55}$ ,  $55 = 2 \cdot 9 \cdot 3 + 1$ ,
- **Guess**  $R \in \mathbb{Q}[z, t]$  such that  $R(F(t, 1), t) = O(t^{(9+1) \cdot (3+1) - 1})$ ,
- **Check** that  $R(t, F(t, 1)) = O(t^{55})$ . ( $\implies R$  is satisfied)

# Solving 5-constellations using a **Hybrid Guess-and-Prove** strategy

**Input:** (The rather big DDE associated with the enumeration of 5-constellations)

**Output:**  $15625t^2F(t, 1)^5 - 31250t^2F(t, 1)^4 + (25000t^2 - 1000t)F(t, 1)^3 - (10000t^2 - 8700t)F(t, 1)^2 + (2000t^2 - 15855t + 16)F(t, 1) - 160t^2 + 8139t - 16 = 0$

- **Draw at random** a prime number  $p$  and some  $c \in \mathbb{F}_p$ ,
- **Compute** upper bounds  $(9, 3)$  on the bidegree of  $M \in \mathbb{F}_p[z, t]$  annihilating  $F(t, 1)$  modulo  $p$ ,
- **Expand** the truncated series  $F(t, 1) \bmod t^{55}$ ,  $55 = 2 \cdot 9 \cdot 3 + 1$ ,
- **Guess**  $R \in \mathbb{Q}[z, t]$  such that  $R(F(t, 1), t) = O(t^{(9+1) \cdot (3+1) - 1})$ ,
- **Check** that  $R(t, F(t, 1)) = O(t^{55})$ . ( $\implies R$  is satisfied)

[Bostan, N., Safey El Din '23]

$\rightsquigarrow$  **elimination strategy**,

$\rightsquigarrow$  **Newton iteration**,

$\rightsquigarrow$  **Hermite Padé approximants**,

$\rightsquigarrow$  **multiplicity lemma**.

# Solving 5-constellations using a **Hybrid Guess-and-Prove** strategy

**Input:** (The rather big DDE associated with the enumeration of 5-constellations)

**Output:**  $15625t^2F(t, 1)^5 - 31250t^2F(t, 1)^4 + (25000t^2 - 1000t)F(t, 1)^3 - (10000t^2 - 8700t)F(t, 1)^2 + (2000t^2 - 15855t + 16)F(t, 1) - 160t^2 + 8139t - 16 = 0$

- **Draw at random** a prime number  $p$  and some  $c \in \mathbb{F}_p$ ,
- **Compute** upper bounds  $(9, 3)$  on the bidegree of  $M \in \mathbb{F}_p[z, t]$  annihilating  $F(t, 1)$  modulo  $p$ ,
- **Expand** the truncated series  $F(t, 1) \bmod t^{55}$ ,  $55 = 2 \cdot 9 \cdot 3 + 1$ ,
- **Guess**  $R \in \mathbb{Q}[z, t]$  such that  $R(F(t, 1), t) = O(t^{(9+1) \cdot (3+1) - 1})$ ,
- **Check** that  $R(t, F(t, 1)) = O(t^{55})$ . ( $\implies R$  is satisfied)

[Bostan, N., Safey El Din '23]

$\rightsquigarrow$  **elimination strategy**,

$\rightsquigarrow$  **Newton iteration**,

$\rightsquigarrow$  **Hermite Padé approximants**,

$\rightsquigarrow$  **multiplicity lemma**.

<u>Strategy</u>	<u>Timing</u>	<u><math>(d_z, d_t)</math></u>
Elimination	2d21h	<b>(9, 3)</b>
Hybrid G-P	<b>2h40min</b>	<b>(5, 2)</b>

## A polynomial system for systems of DDEs

Consider

$$\begin{cases} F_1 = f_1(\mathbf{u}) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, \mathbf{u}), \\ \vdots \\ F_n = f_n(\mathbf{u}) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, \mathbf{u}). \end{cases} \quad (\text{SDDEs})$$

## A polynomial system for systems of DDEs

Consider

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u). \end{cases} \quad (\text{SDDEs})$$

**Perturbe** (SDDEs) and **define** the “numerators”  $E_1, \dots, E_n$  and the polynomials

$$\text{Det} := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{x_1} E_n & \dots & \partial_{x_n} E_n \end{pmatrix} \quad \text{and} \quad P := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_{n-1}} E_1 & \partial_u E_1 \\ \vdots & \ddots & \vdots & \vdots \\ \partial_{x_1} E_{n-1} & \dots & \partial_{x_{n-1}} E_{n-1} & \partial_u E_{n-1} \\ \partial_{x_1} E_n & \dots & \partial_{x_{n-1}} E_n & \partial_u E_n \end{pmatrix},$$



## A polynomial system for systems of DDEs

Consider

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u). \end{cases} \quad (\text{SDDEs})$$

**Perturbe** (SDDEs) and **define** the “numerators”  $E_1, \dots, E_n$  and the polynomials

$$\text{Det} := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{x_1} E_n & \dots & \partial_{x_n} E_n \end{pmatrix} \quad \text{and} \quad P := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_{n-1}} E_1 & \partial_u E_1 \\ \vdots & \ddots & \vdots & \vdots \\ \partial_{x_1} E_{n-1} & \dots & \partial_{x_{n-1}} E_{n-1} & \partial_u E_{n-1} \\ \partial_{x_1} E_n & \dots & \partial_{x_{n-1}} E_n & \partial_u E_n \end{pmatrix},$$

**Set** up the duplicated polynomial system ( $\mathcal{S}_{\text{dup}}$ ), consisting in the  $nk$  duplications of the polynomials  $E_1, \dots, E_n, \text{Det}, P$ . It has  $nk(n+2)$  variables and equations.

## A polynomial system for systems of DDEs

Consider

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u). \end{cases} \quad (\text{SDDEs})$$

**Perturbe** (SDDEs) and **define** the “numerators”  $E_1, \dots, E_n$  and the polynomials

$$\text{Det} := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{x_1} E_n & \dots & \partial_{x_n} E_n \end{pmatrix} \quad \text{and} \quad P := \det \begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_{n-1}} E_1 & \partial_u E_1 \\ \vdots & \ddots & \vdots & \vdots \\ \partial_{x_1} E_{n-1} & \dots & \partial_{x_{n-1}} E_{n-1} & \partial_u E_{n-1} \\ \partial_{x_1} E_n & \dots & \partial_{x_{n-1}} E_n & \partial_u E_n \end{pmatrix},$$

**Set** up the duplicated polynomial system  $(\mathcal{S}_{\text{dup}})$ , consisting in the  $nk$  duplications of the polynomials  $E_1, \dots, E_n, \text{Det}, P$ . It has  $nk(n+2)$  variables and equations.

**Compute** a non-trivial element of  $(\langle \mathcal{S}_{\text{dup}} \rangle : \det(\text{Jac}_{\mathcal{S}_{\text{dup}}})^\infty) \cap \mathbb{K}[t, z_0, \epsilon]$ , then set  $\epsilon$  to 0.