

# Dynamic combinatorics, complex systems and applications to physics

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Jordan

Mathematics

Chaos Theory

Continuous & Discrete Modelisation

Image Processing

Computer Science

Complex Systems Complexity

Business Banking

Computation Techniques

Decision Making

Physics

Artificial Intelligence

Mechatronics

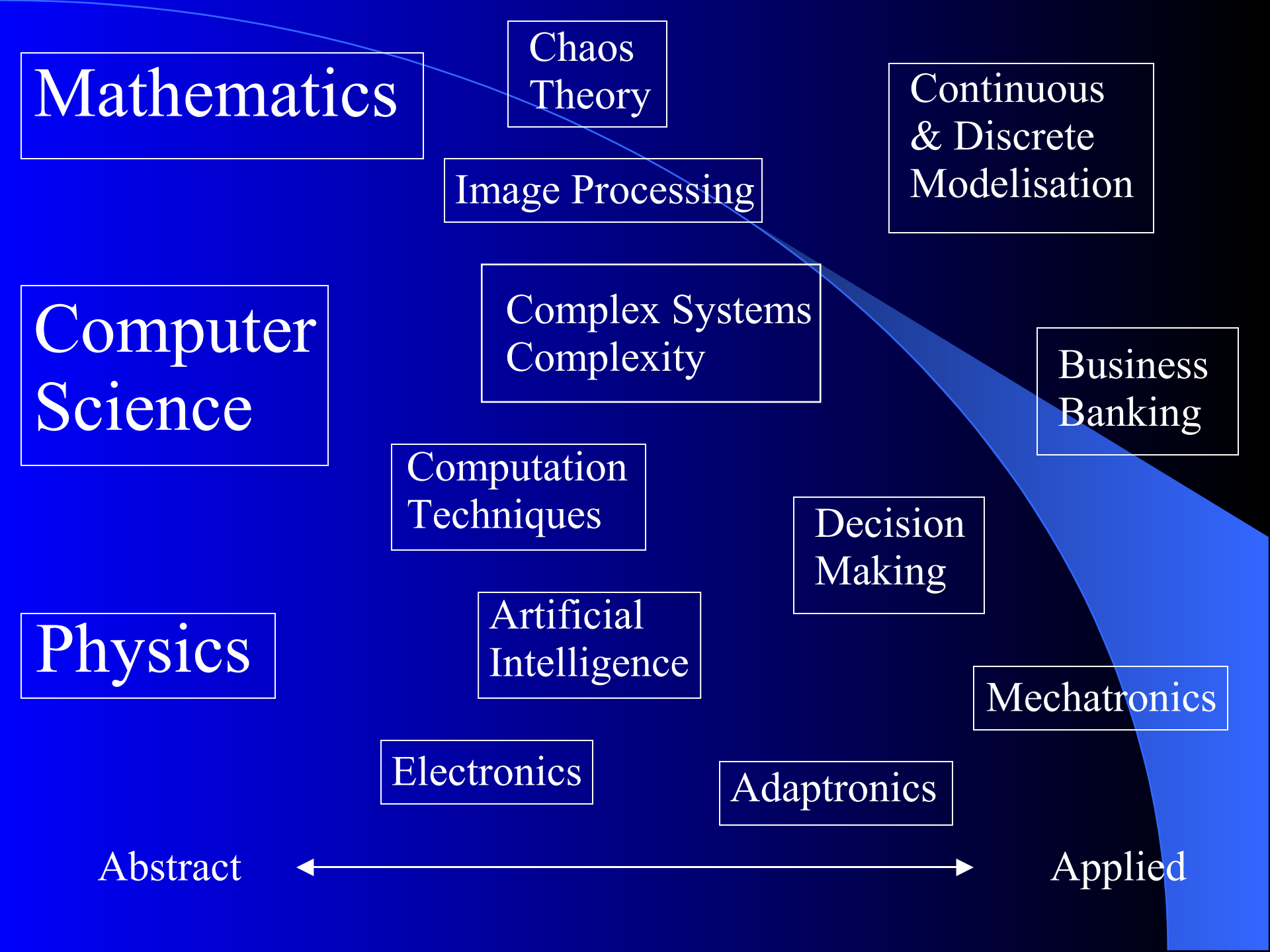
Electronics

Adaptronics

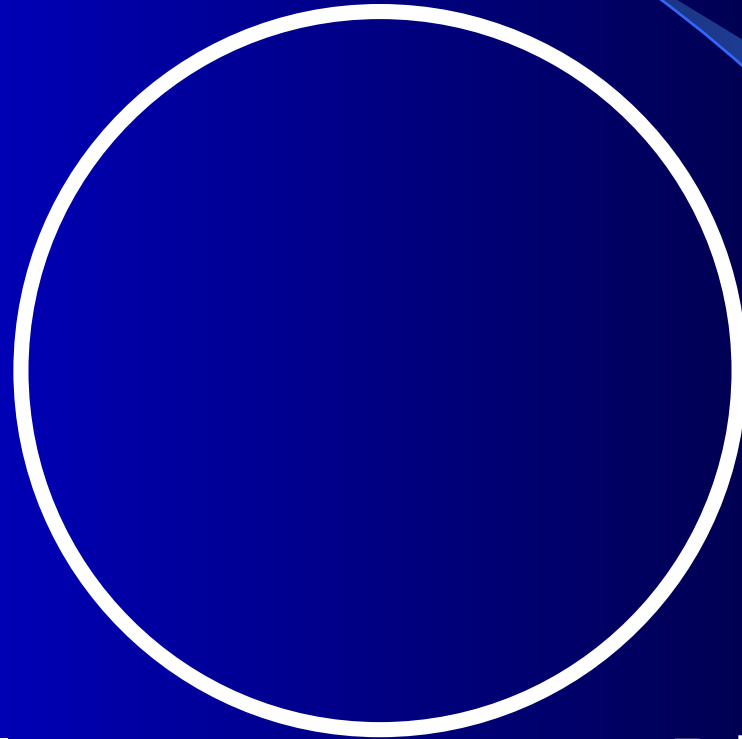
Abstract



Applied



# Combinatorics (discrete mathematics)



Information  
(computer sci.)

Physics  
(classical/quant.)

Combinatorics  
(discrete mathematics)



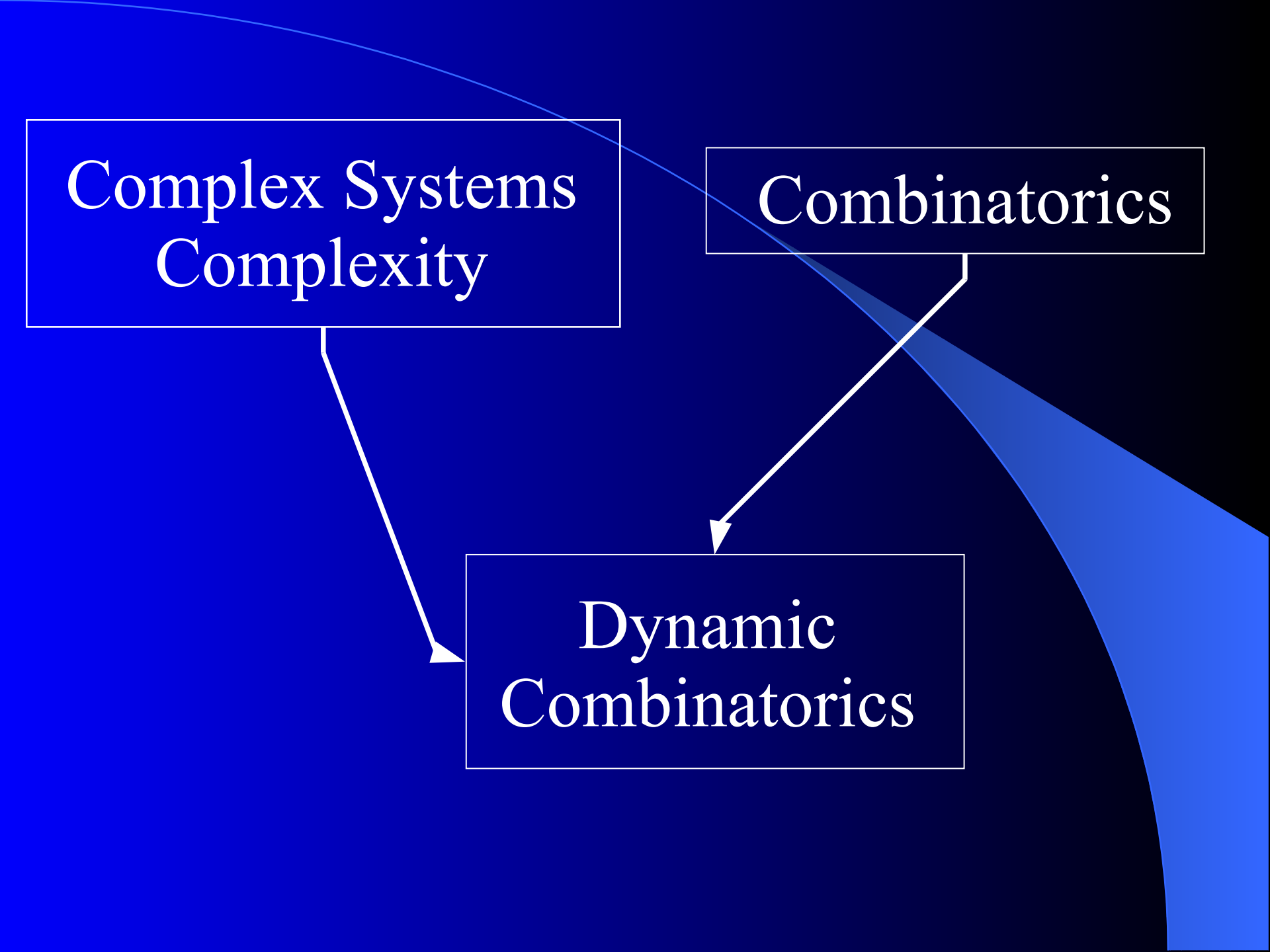
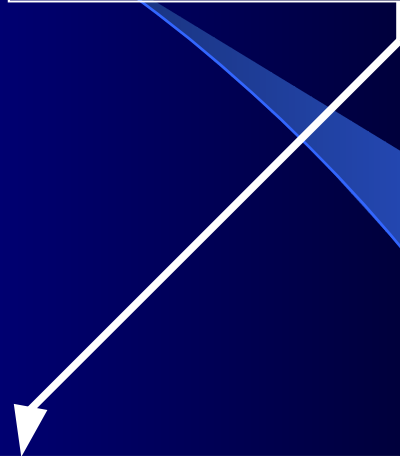
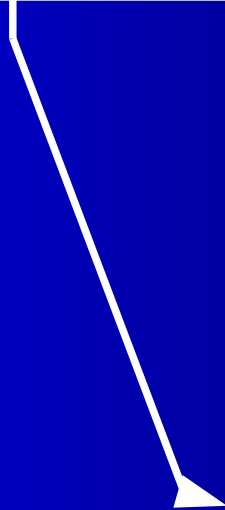
Information  
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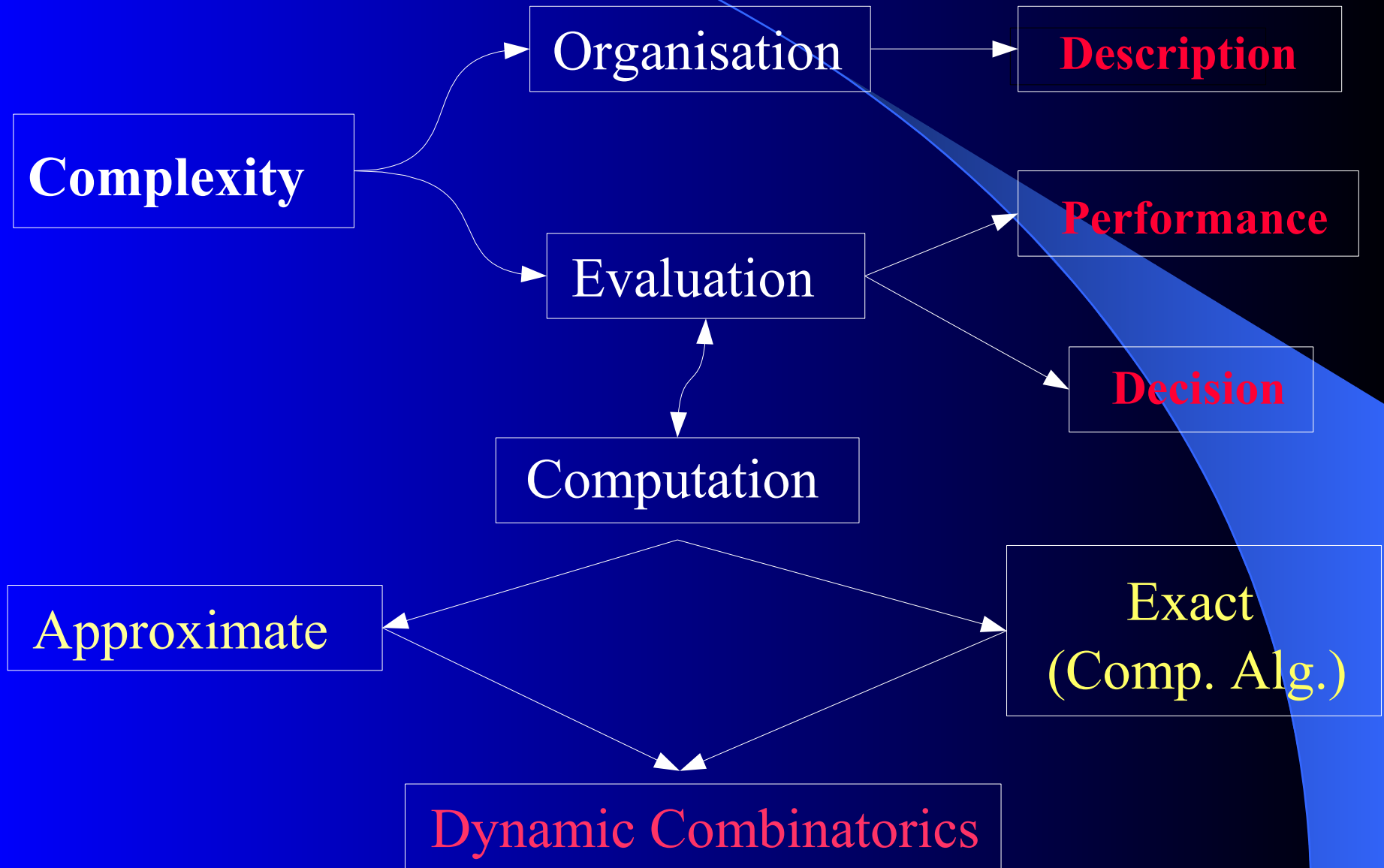
Complex Systems  
Complexity

Combinatorics

Dynamic  
Combinatorics



# Problematics of Dynamic Combinatorics



# What is the Legacy ?

Mainly:

- ✓ Data Structures
- ✓ Programs
- ✓ Theorems
- ✓ Computation rules
- ✓ Experiments

# What is the Legacy ? (Cont'd)

## Mathematics

- Noncommutative
- Representations
- Formulas, Universal Algebra
- Deformations

## Comp. Sciences

- Words
- Automata Transition Structures
- Trees with Operators
- q-analogues

## Physics

- Strings of operators
- Fields, Flows, Dynamic Systems (Chaos, Catastrophes)
- Diagrams
- Quantum Groups

Combinatorics



# Combinatorics

## ... on words

- Languages
- Theory of codes
- Automata
- Transition structures
- Grammars
- Transducers
- Rational and algebraic expressions
- ...

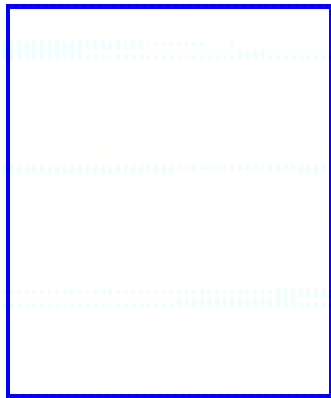
## enumerative, analytic

- Polyominoes
- Paths (Dyck, ...)
- Configurations
- q-grammars
- Generating Functions
- Continued Fractions (mono, multivariate, ..)
- Orthogonal Polynomials
- ...

## algebraic

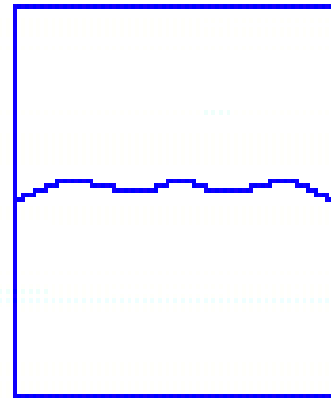
- Non commutative Continued fractions
- Representations of groups and deformations
- Quantum Groups
- Functors
- Characters
- Special Functions
- ...

A first example . . .



P

$n \rightarrow n+1$   
i = one unit



R

$n = n+1$

- $a_1 = 0$
- $a_{i+1} \leq a_i + 1$

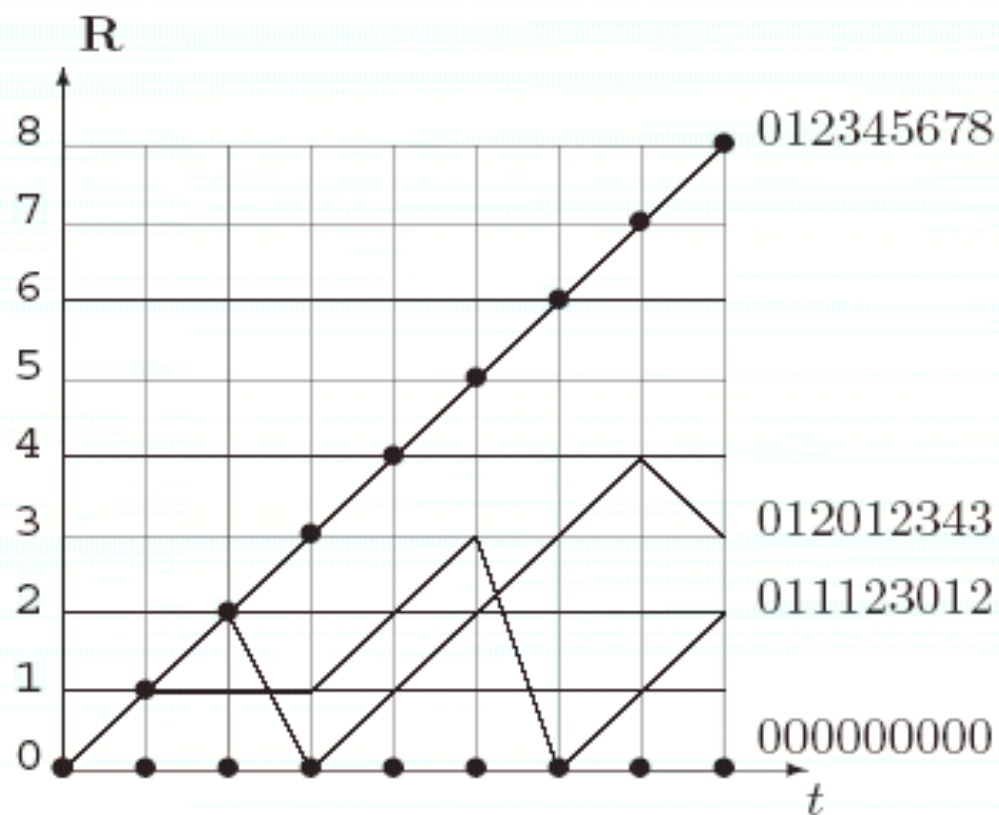
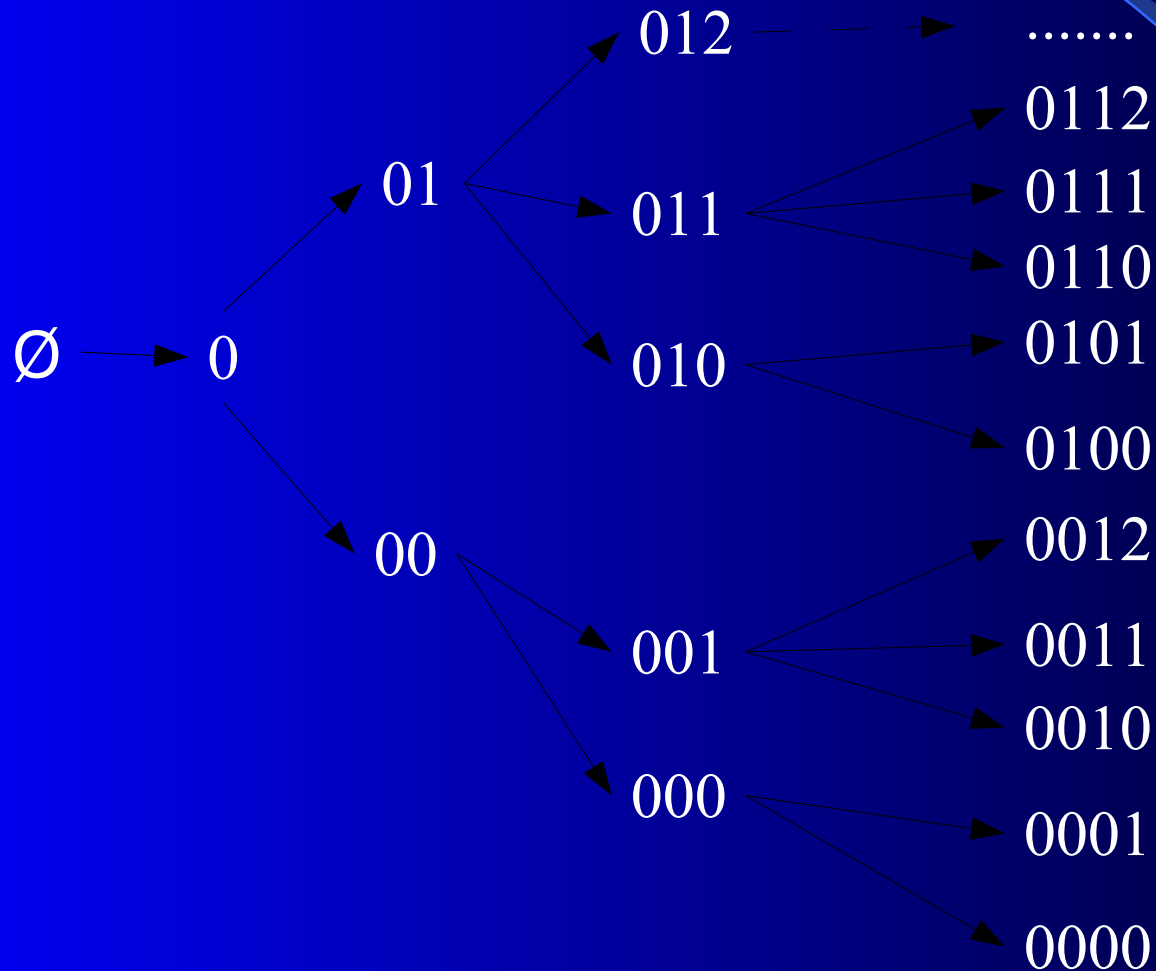


Figure 4.2: Maximal, minimal (dotted) and two intermediate trajectories. Their codes are on the right.

One can arrange all the trajectories on a dynamic graph





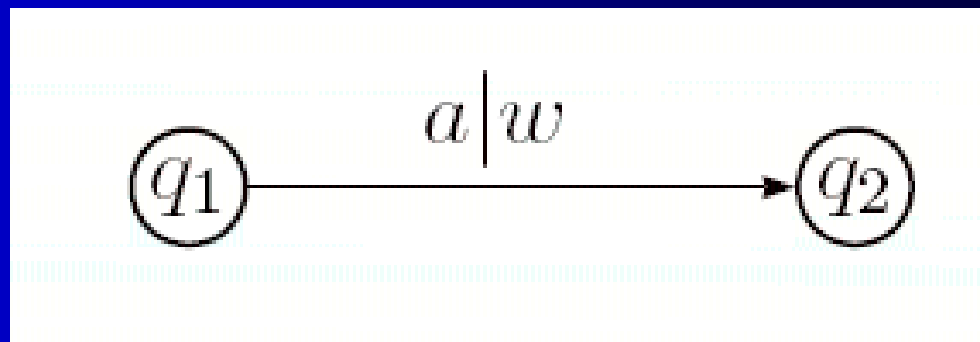


We will return to the Dyck paths later on.  
For the moment let us define what is a transition structure.

**Definition** (transition structure) : It is a graph (finite or infinite) with its arcs marked with pairs

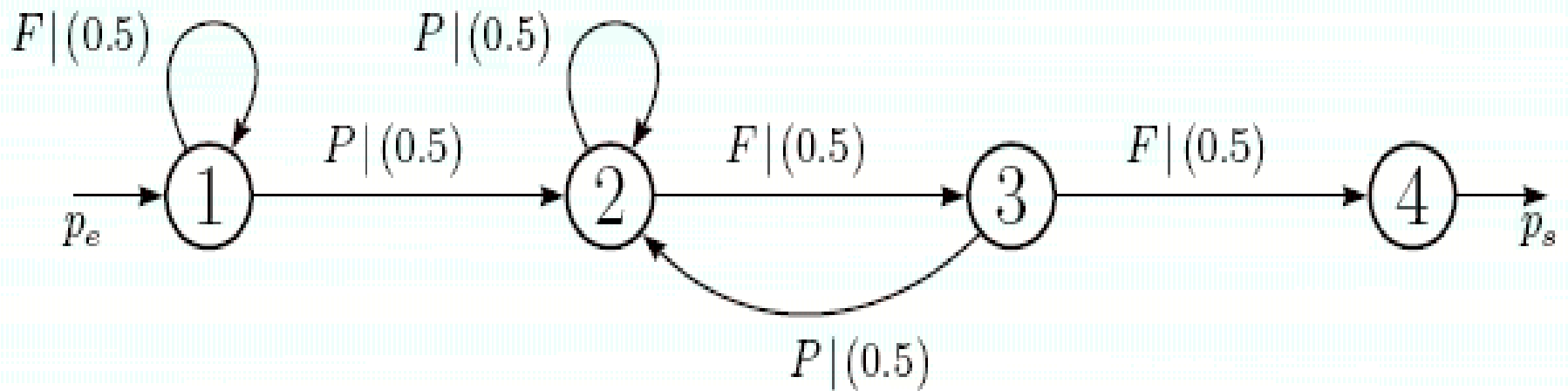
(command letter | coefficient)

Examples : Prisoner's dilemma, Markov chains, classical engineering.

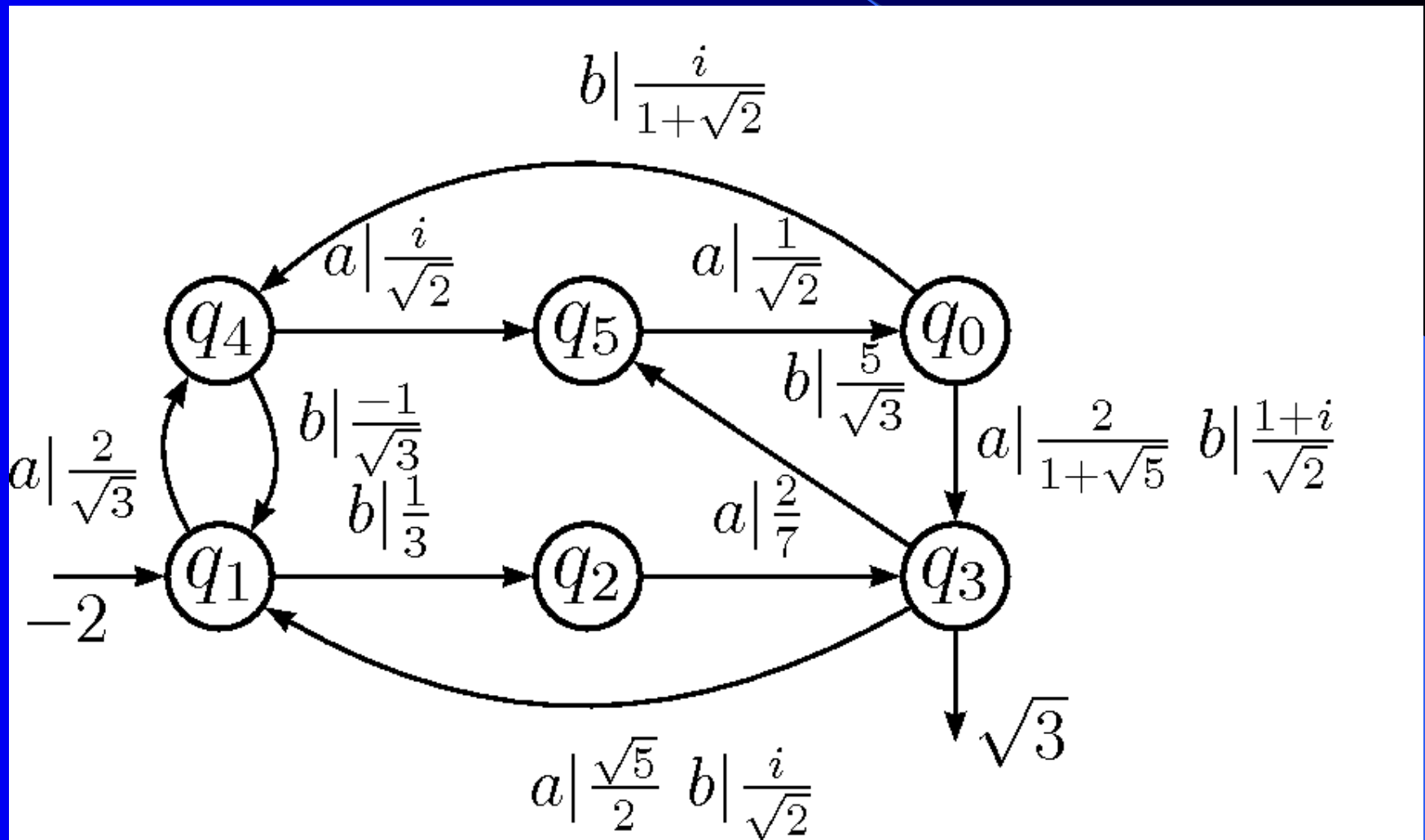




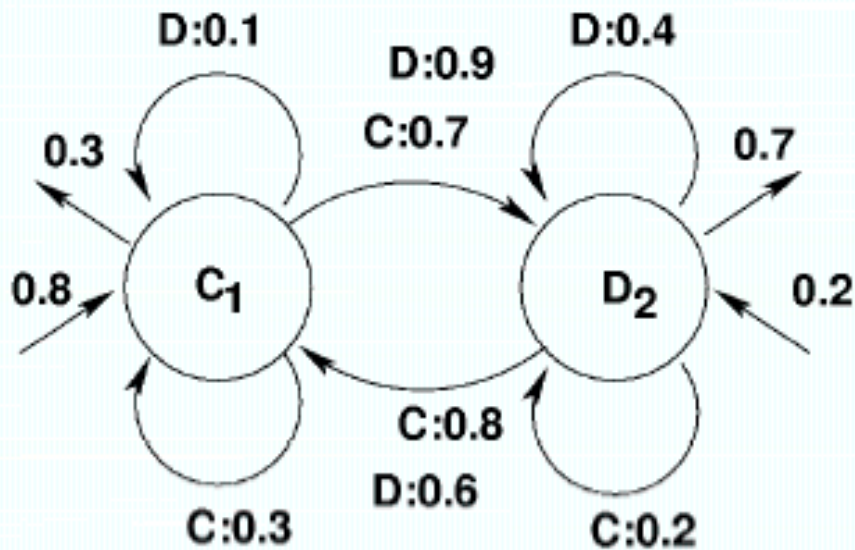
Example : A Markov chain generated by a game.



Example : An automaton generated by arbitrary transition coefficients.



# Example of Probabilistic Automaton



## LINEAR REPRESENTATION

$\begin{bmatrix} 0.8 & 0.2 \end{bmatrix}$  input vector

	1	2		1	2
1	0.3	0.7	0.1	0.9	1
2	0.8	0.2	0.6	0.4	2

$M(C)$

$M(D)$

$\begin{bmatrix} 0.3 \\ 0.7 \end{bmatrix}$  output vector

# Behaviour of an Automaton and how to compute it effectively

An automaton is a **machine** which takes a string (sequence of letters) and returns a **value**.

This value is computed as follows :

1) The **weight** of a path is the product of the weights (or coefficients) of its edges

2) The **label** of a path is the product (concatenation) of the labels of its edges

## Behaviour ... (cont'd)

3) The **behaviour** between two states « r,s » w.r.t. A word « w » is the product of

3a) the ingoing coefficient of the first state (here « r ») by

3b) the sum of the weights of the paths going from « r » to « s » with label « w » by

3c) the outgoing coefficient of the second state (here « s »)

## Behaviour ... (cont'd)

4) The **behaviour** of the automaton under consideration w.r.t. a word «  $w$  » is then the sum of all the behaviours of the automaton between two states «  $r,s$  » for all possible pairs of states.

## Behaviour ... (cont'd)

There is a simple formula using the linear representation. The **behaviour** of an automaton with linear representation  $(I, M, T)$  is the product

$$IM(w)T$$

where  $M(w)$  is the canonical extension of  $M$  to the strings.

$$M(a_1 a_2 \dots a_n) = M(a_1) M(a_2) \dots M(a_n)$$

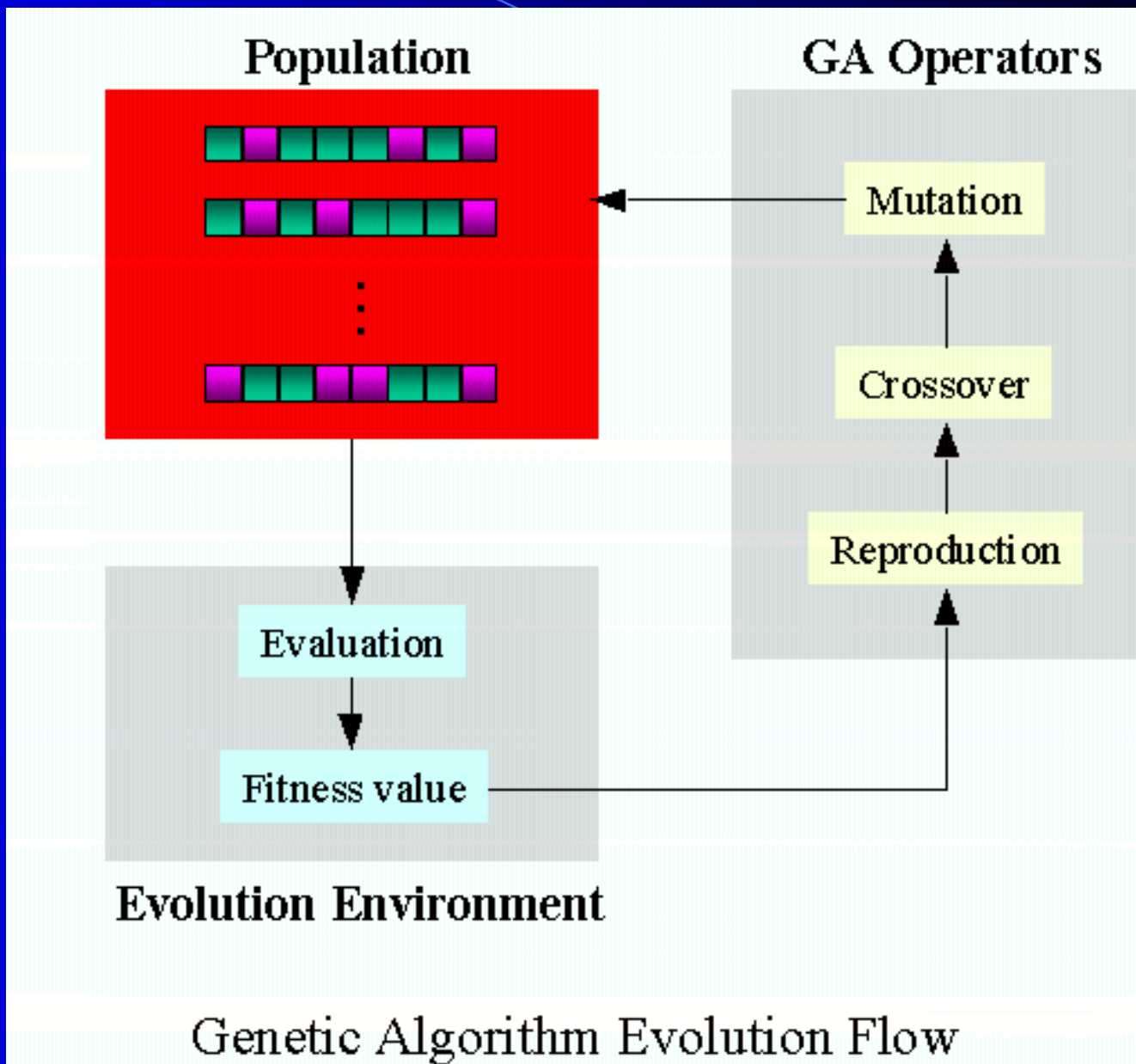
## Behaviour ... (end)

The behaviour, as a function on words belongs to the rational class. If time permits, we will return to its complete calculation as a **rational expression** and the problem of its algorithmic evaluation by means of special cancellation operators. Linear representations can also be used to compute **distances between automata**.



Example -> use of genetic algorithms to control indirect (set of) parameters : the spectrum of a matrix.

# Genetic algorithms : general pattern



# Genetic algorithms : implementation

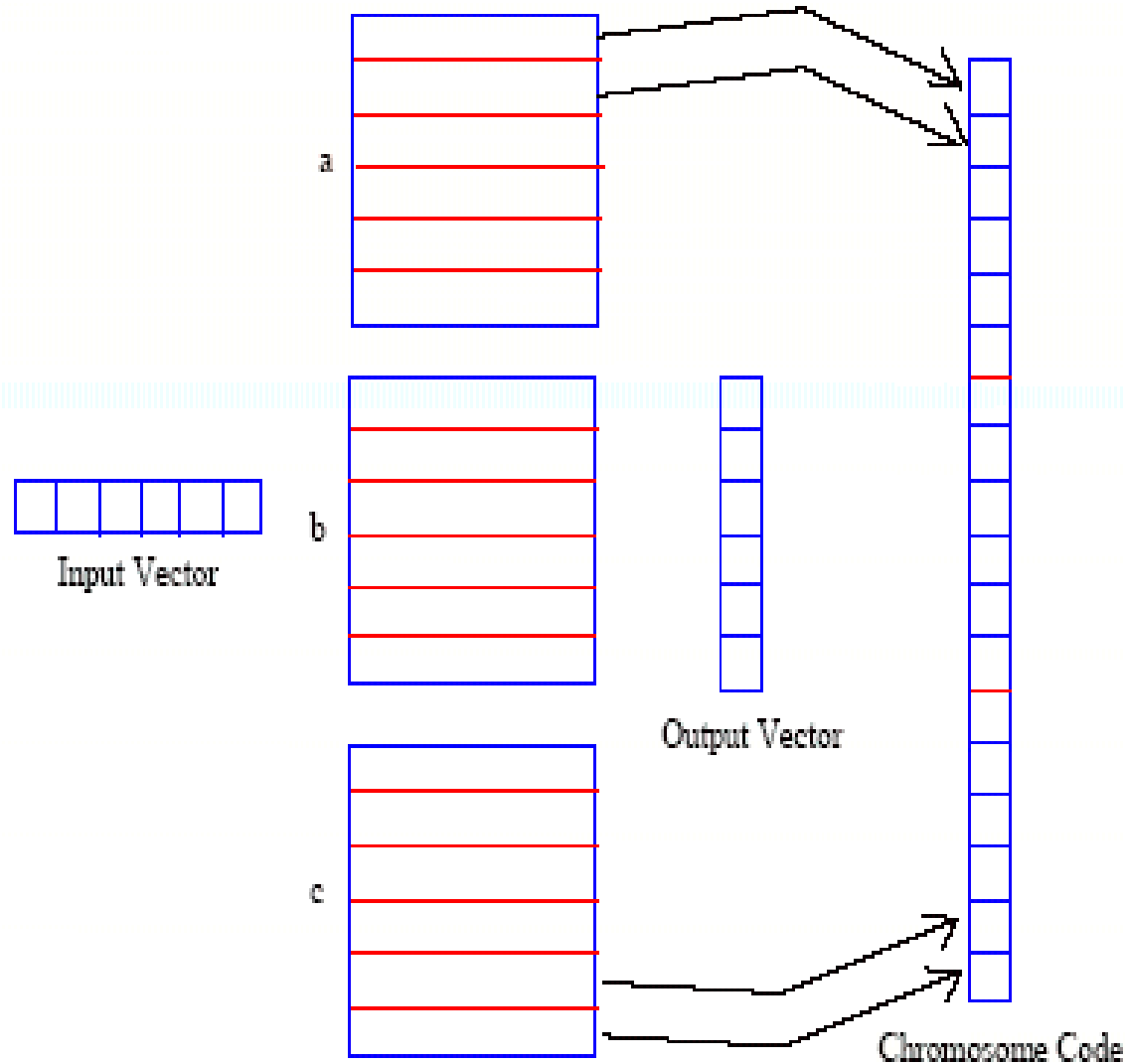


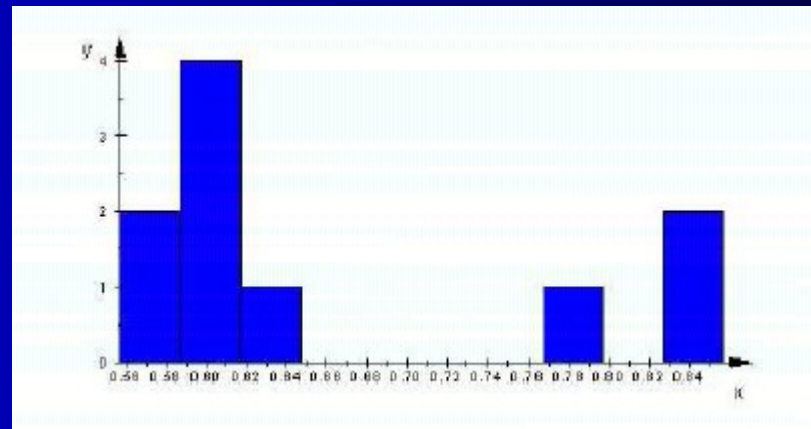
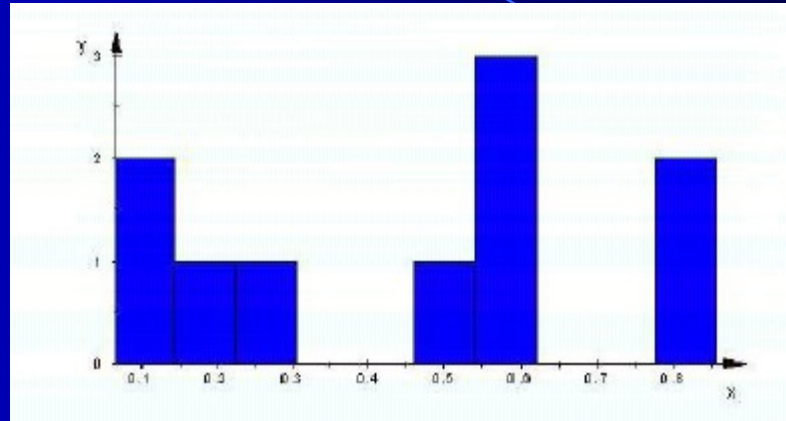
Figure 4.13: Chromosome code

# Genetic algorithms : implementation

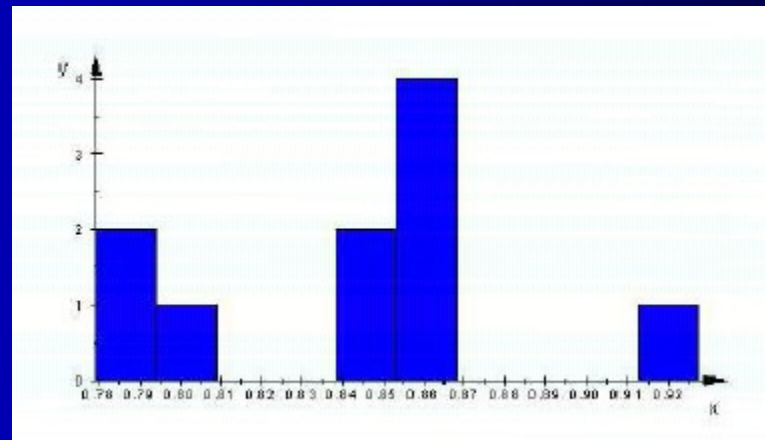
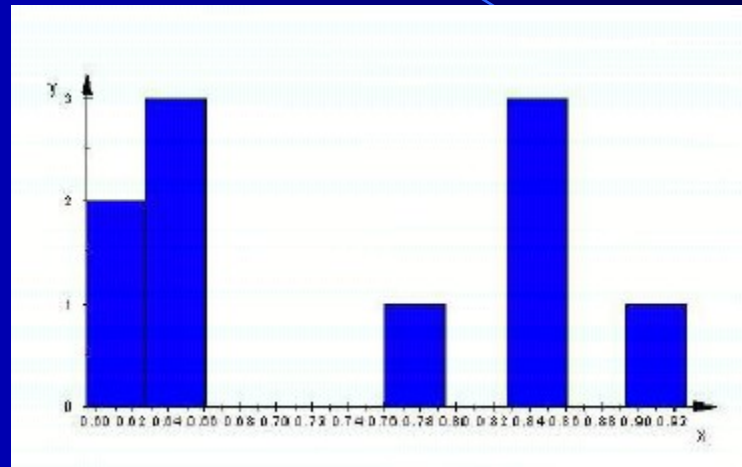
Below, the results of an experiment aiming to control the second greatest eigenvalue of the transfer matrix of a population of probabilistic automata.

- The fitness function of each automaton corresponds to the second greatest eigenvalue (in module). The first being, of course, of value 1.

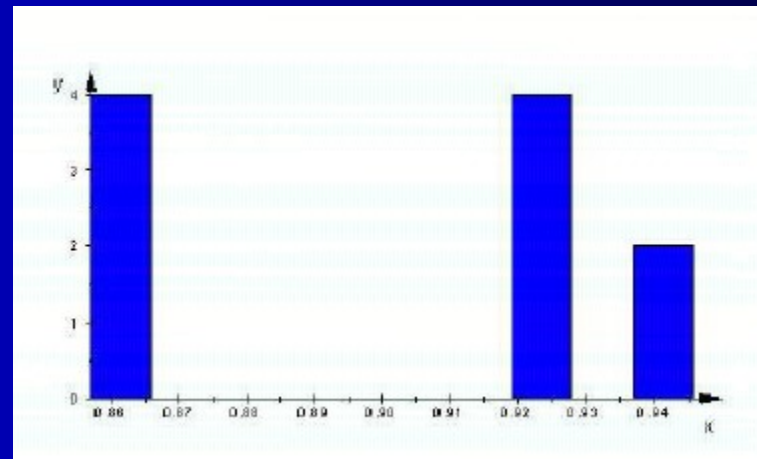
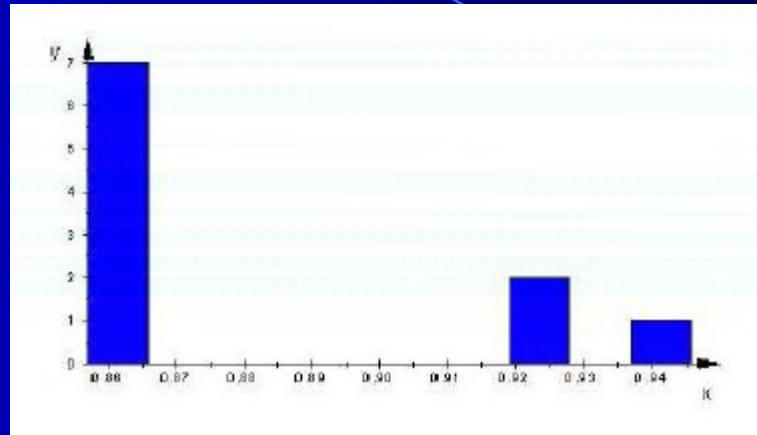
# Genetic algorithms ; results



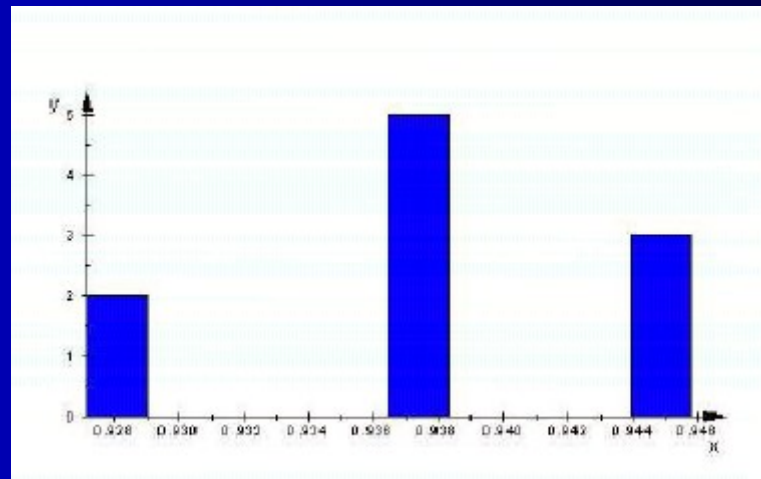
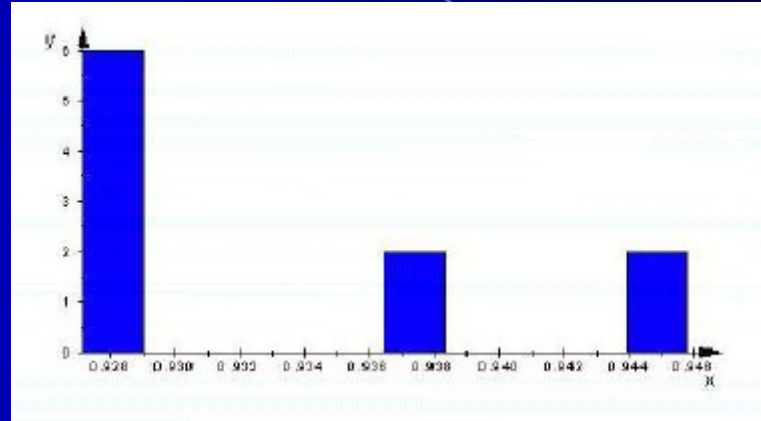
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# Genetic algorithms ; results

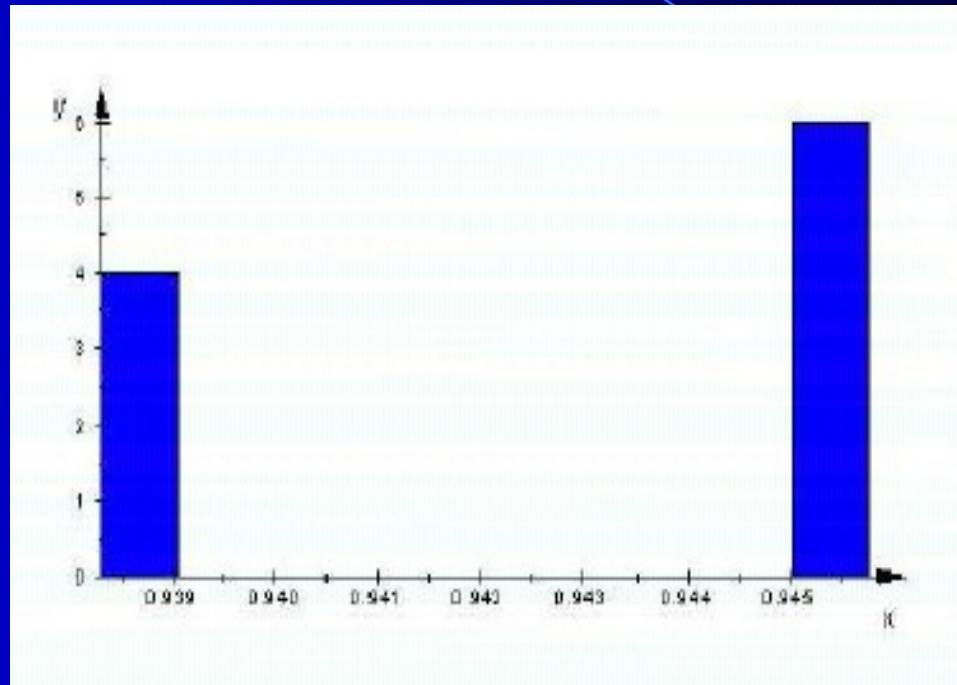


# Genetic algorithms ; results





# Genetic algorithms ; results



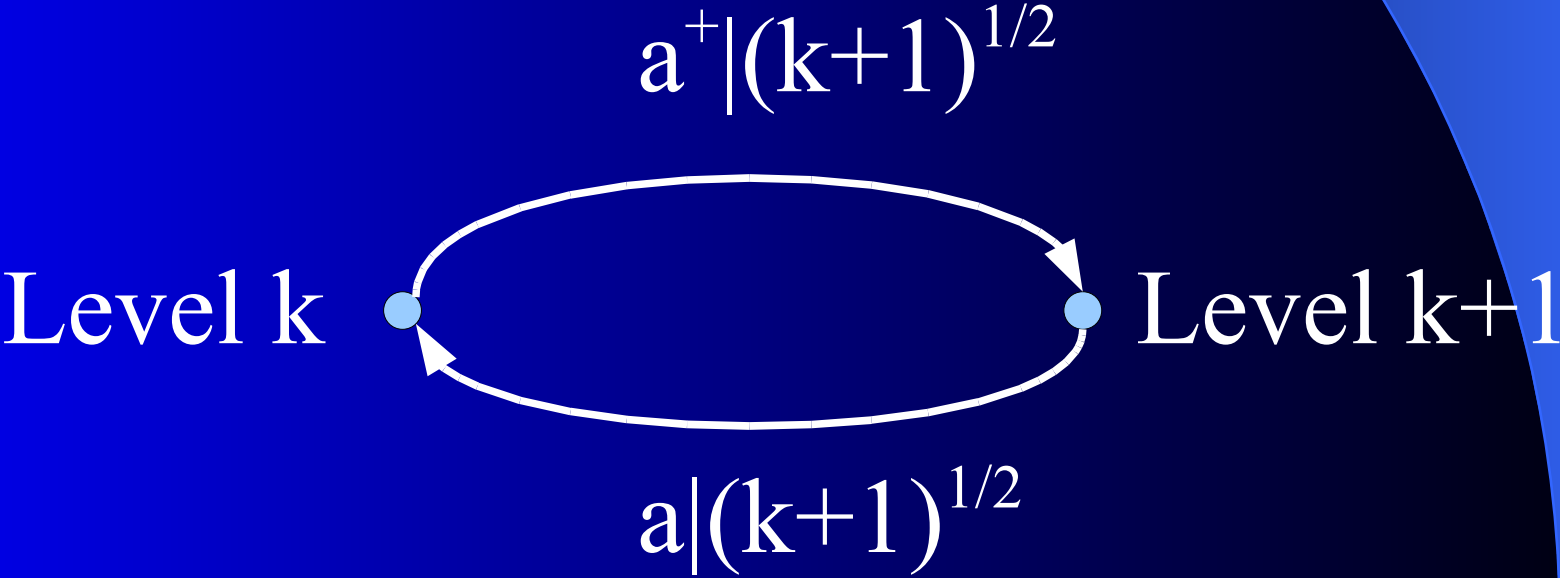
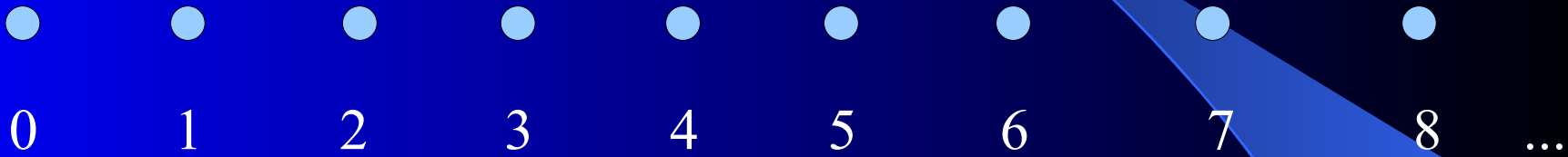
Final result : the population is rendered homogeneous

# General transition systems



- Automata (finite number of edges)
- Sweedler's duals (physics, finite number of states)
- Representations
- Level systems (Quantum Physics)
- Markov chains (prob. automata when finite)

# Example in Physics : annihilation/creation operators



The (classical, for bosons) normal ordering problem goes as follows.

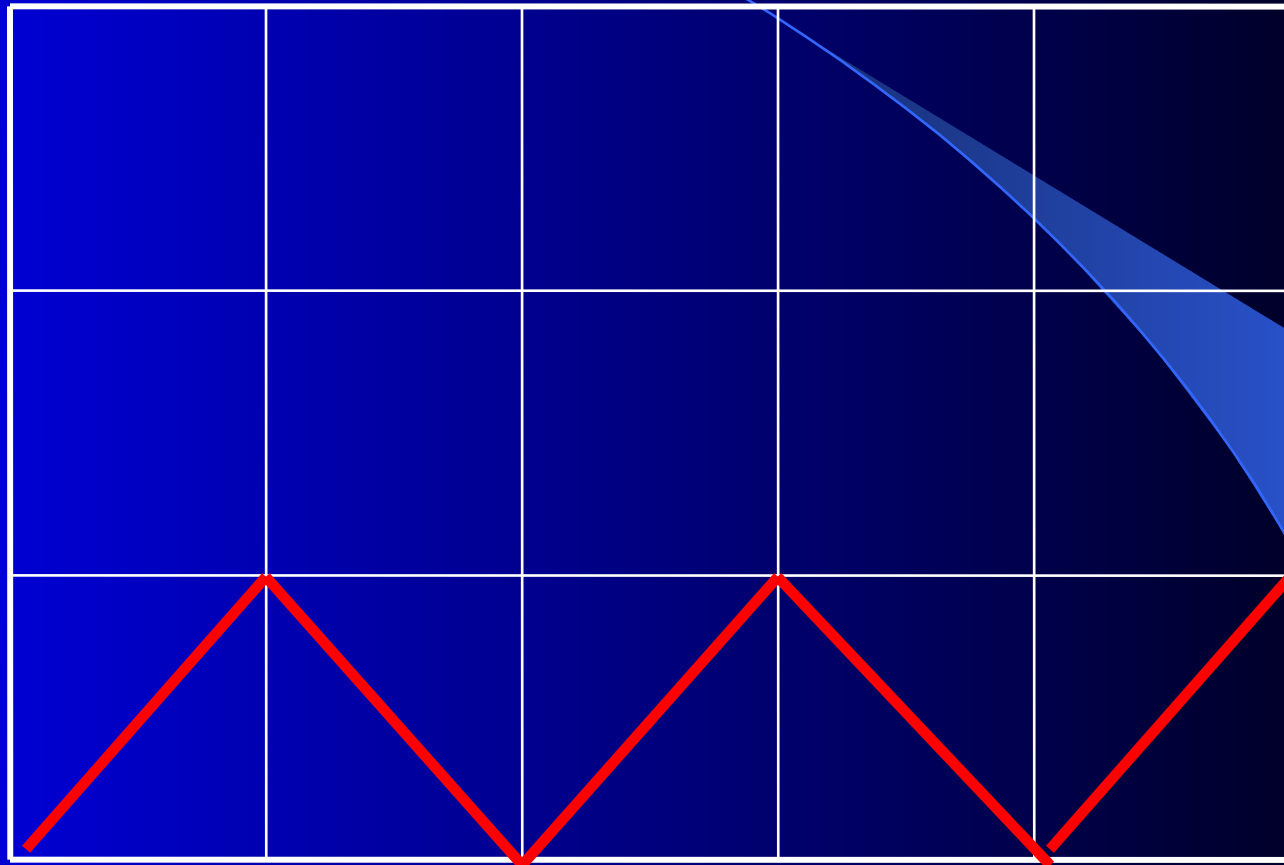
- Weyl (two-dimensional) algebra defined as  
 $\langle a^+, a ; [a, a^+] = 1 \rangle$
- Known to have no (faithful) representation by bounded operators in a Banach space.

There are many « combinatorial » (faithful) representations by operators. The most famous one is the Bargmann-Fock representation

$$a \rightarrow d/dx ; a^+ \rightarrow x$$

where  $a$  has degree  $-1$  and  $a^+$  has degree  $1$ .

Example with  $\Omega = a^+ a a^+ a a^+$



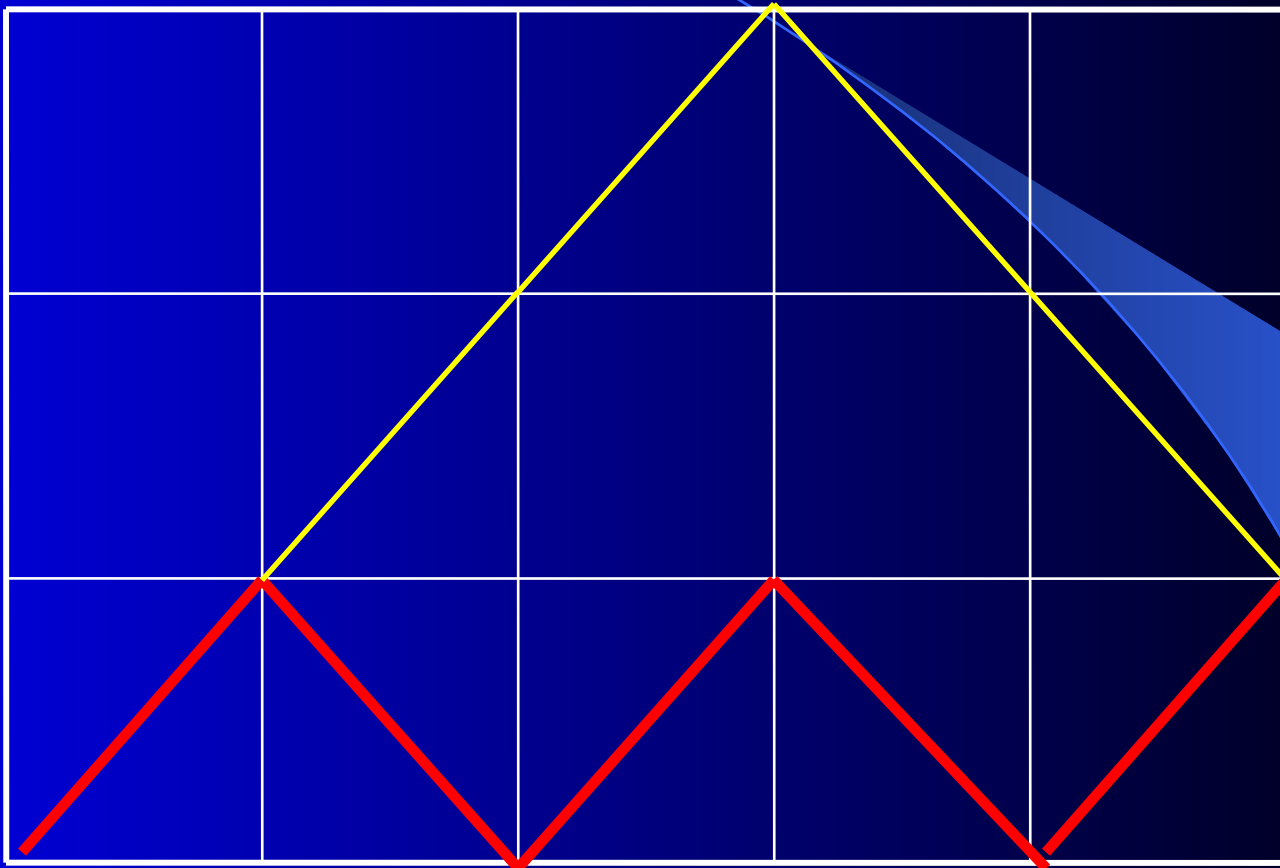
$a^+$

$a$

$a^+$

$a$

$a^+$



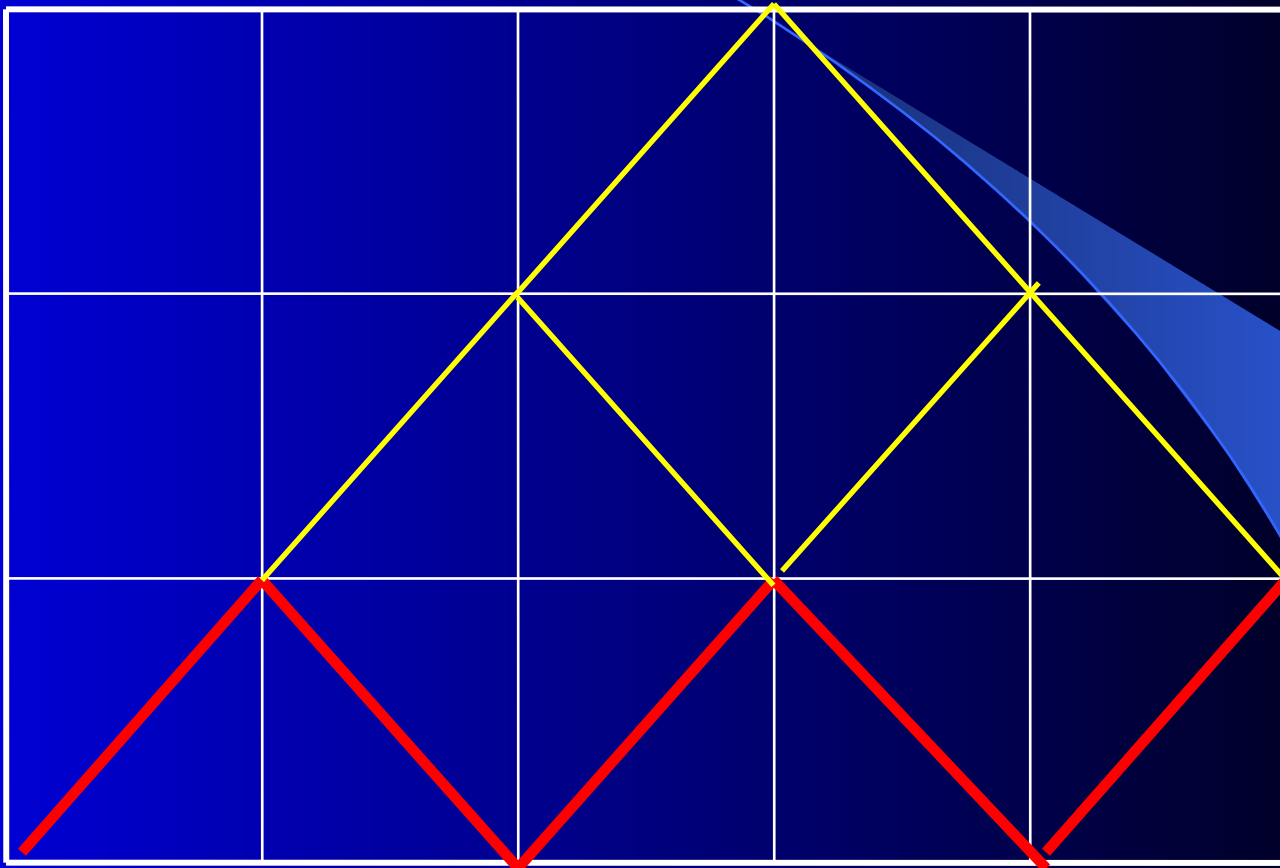
$a^+$

$a$

$a^+$

$a$

$a^+$



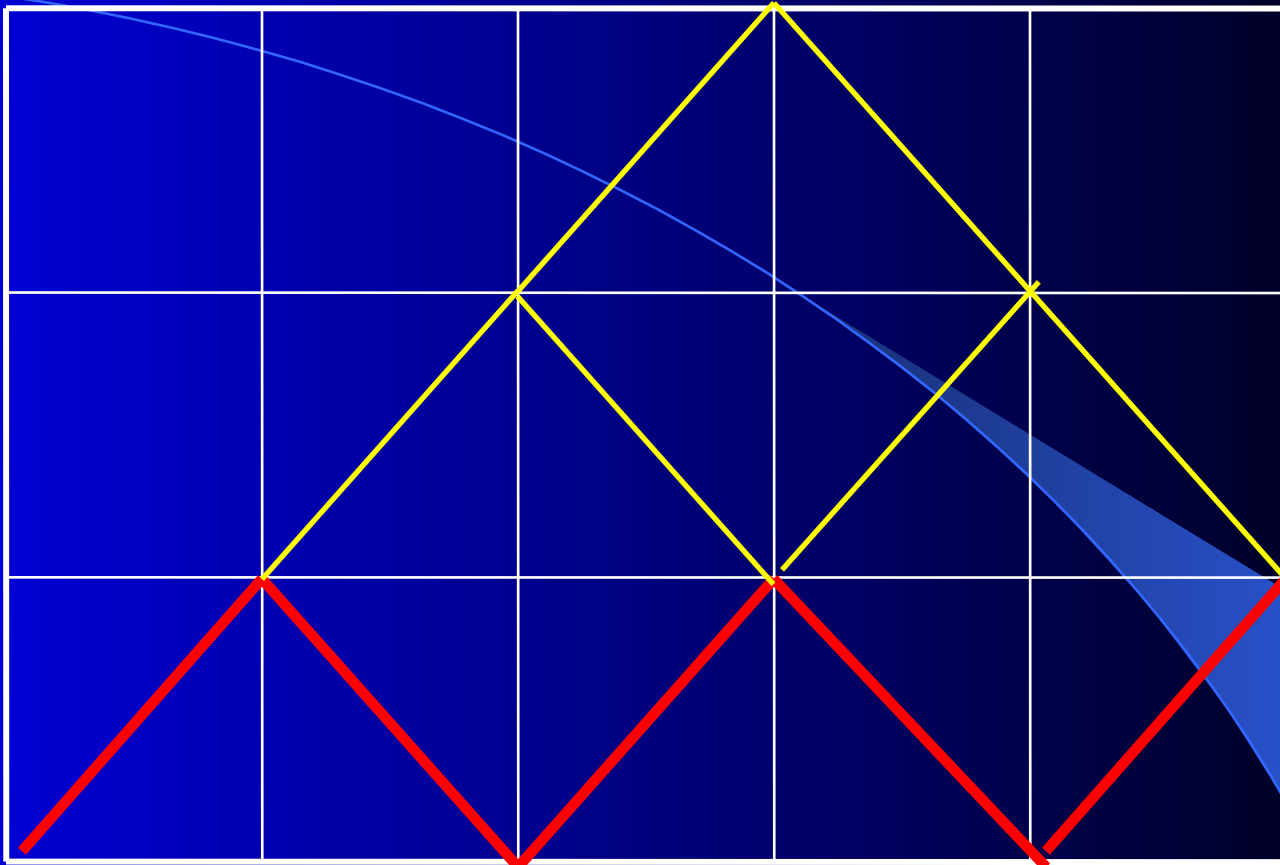
$a^+$

$a$

$a^+$

$a$

$a^+$



$a^+$     $a$     $a^+$     $a$     $a^+$

$$a^+aa^+aa^+ = 1 a^+a^+a^+aa + 3 a^+a^+a + 1 a^+$$



Through Bargmann-Fock representation

$$a \rightarrow d/dx ; a^+ \rightarrow x$$

Operators who have only one annihilation have exponentials who act as one-parameter groups of substitutions.

One can thus use computer algebra to determine their generating function.

For example, with

$$\Omega = a^{+2}a a^+ + a^+a a^{+2}$$

the computation reads



One parameter group by  $f(v(u(x)+\lambda))$ ;  $v$  is reciprocal of  $u$

```
> T1(lambda, x) := (-4 * (-1 / (4 * x^2) + lambda)) ^ (-1/2);
```

$$T1(\lambda, x) := \frac{1}{\sqrt{\frac{1}{x^2} - 4\lambda}}$$

We suppose  $x > 0$

```
> T1 := (lambda, x) -> x / ((1 - 4 * lambda * x^2) ^ (1/2));
```

$$T1 := (\lambda, x) \rightarrow \frac{x}{\sqrt{1 - 4\lambda x^2}}$$

Checking the tangent vector

```
> subs(lambda=0, diff(T1(lambda, x), lambda));
```

$$2x^3$$

... and the one-parameter group property

```
> simplify(T1(lambda1, T1(lambda2, x)) - T1(lambda1+lambda2, x));
```

$$0$$

And the action of  $\exp(\lambda \omega)$  on  $[f(x)]$  is

$$\begin{aligned}
 U_\lambda (f) &= x^{-\frac{3}{2}} f(s_\lambda (x)).(s_\lambda (x))^{\frac{3}{2}} \\
 &= \sqrt[4]{\frac{1}{(1-4\lambda x^2)^3}} f\left(\sqrt{\frac{x^2}{1-4\lambda x^2}}\right)
 \end{aligned}$$

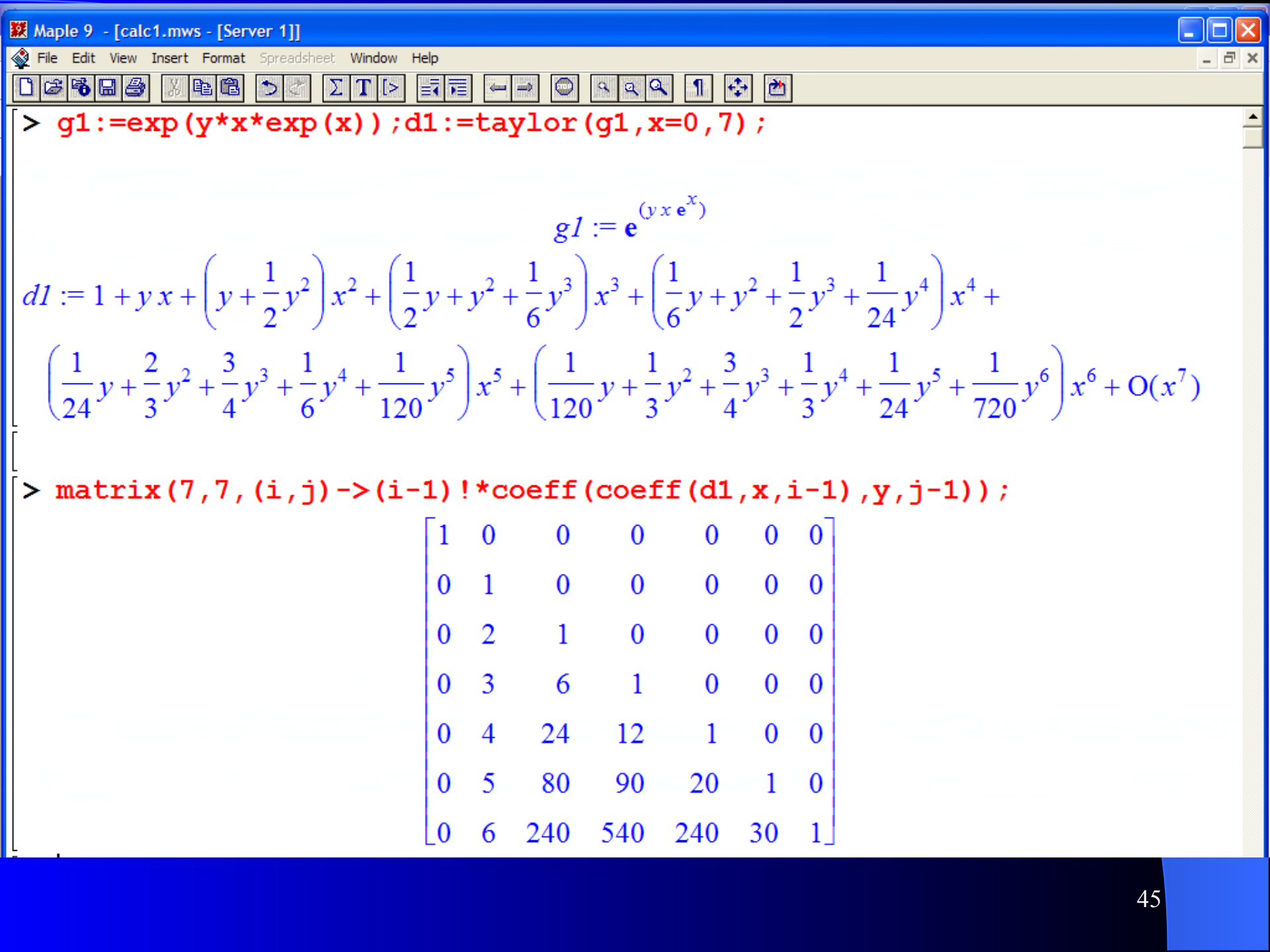
which explains the prefactor. Again we can check by computation that the composition of  $(U_\lambda)$ s amounts to simple addition of parameters !!

Now suppose that  $\exp(\lambda \omega)$  is in normal form.

In view of Eq1 (slide 15) we must have

$$\exp(\lambda \omega) = \sum_{n \geq 0} \frac{\lambda^n \omega^n}{n!} = \sum_{n \geq 0} \frac{\lambda^n}{n!} x^{ne} \sum_{k=0}^{ne} S_\omega(n, k) x^k \left(\frac{d}{dx}\right)^k$$

So, using this new technique, one can compute easily the coefficients of the matrix giving the normal forms.



**For these one-parameter groups and conjugates of vector fields**

**G. H. E. Duchamp, K.A. Penson, A.I. Solomon, A. Horzela and P. Blasiak,**

**One-parameter groups and combinatorial physics,**

**Third International Workshop on Contemporary Problems in Mathematical Physics (COPROMAPH3), Porto-Novo (Benin), November 2003. arXiv : quant-ph/0401126.**

**For the Sheffer-type sequences and coherent states**

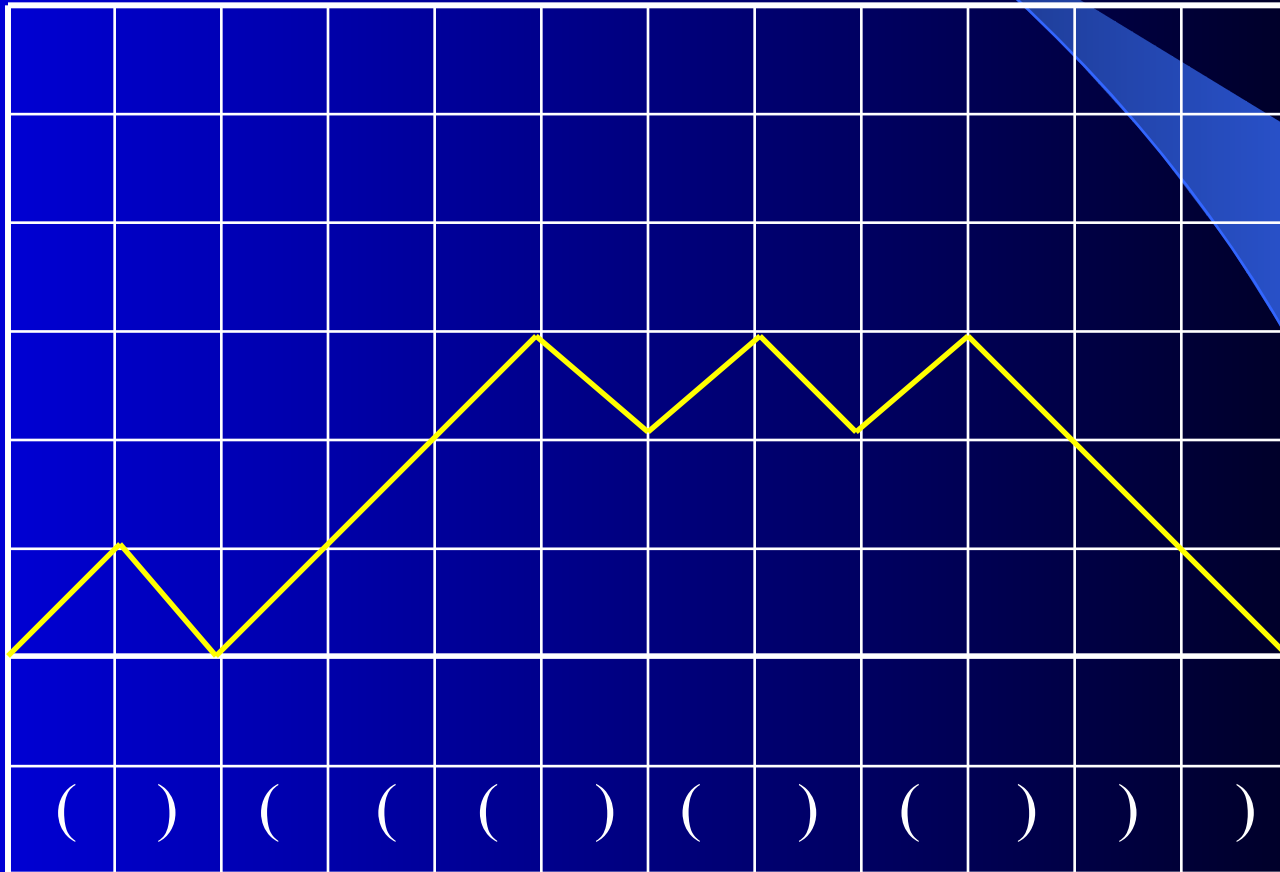
**K A Penson, P Blasiak, G H E Duchamp, A Horzela and A I Solomon,**

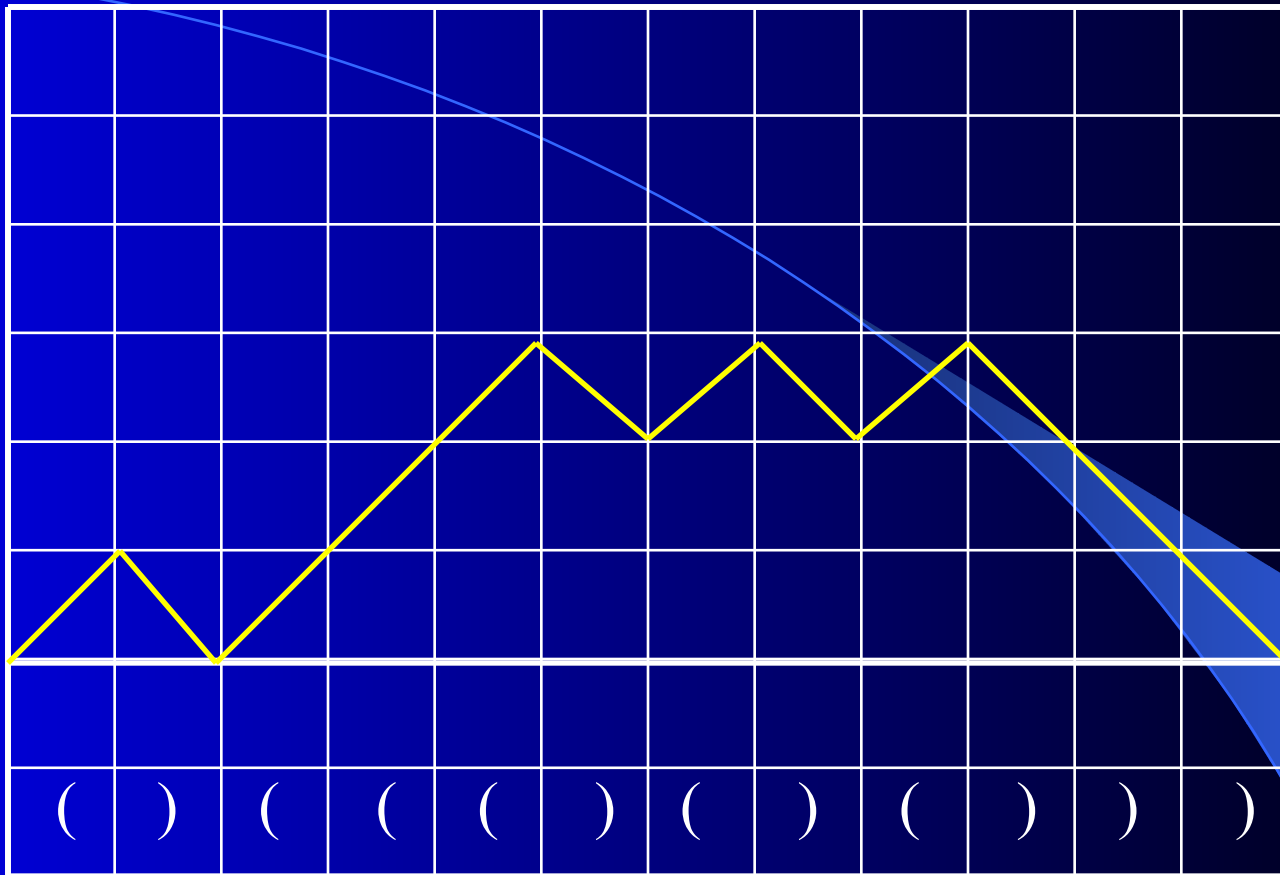
**Hierarchical Dobinski-type relations via substitution and the moment problem,**

**J. Phys. A: Math. Gen. 37 3457 (2004) arXiv : quant-ph/0312202**

# A second application : Dyck paths

(systems of brackets, trees, physics, ...)



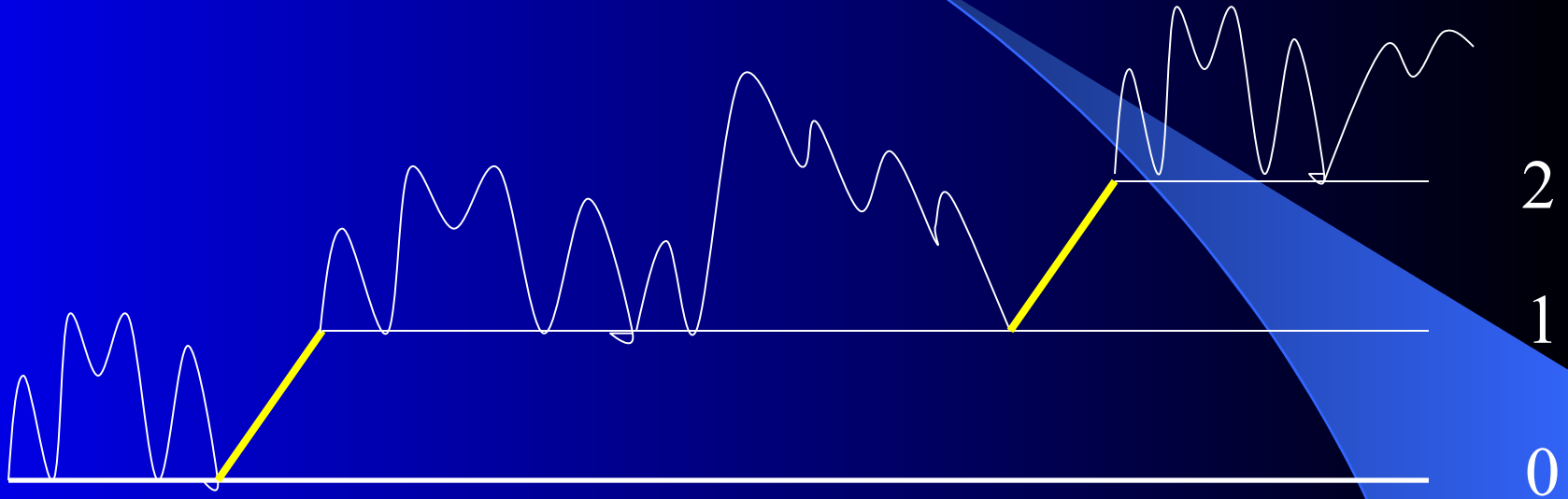


Equation :  $D = \text{void} + (D) D \dots$  one counts strings using an « x » by bracket and one finds  $T(x) = x^0 + x^2 T^2(x)$  which can be solved by elementary methods ...

$$x^2 T^2 - T + 1 = 0 \quad \text{Variable : } T \quad \text{Parameter : } x$$



# Change of level (physics)



Positifs =  $D(aD)^*$

$$Pos := \frac{Dyck}{1 - x Dyck}$$

```
> solve(x^2*T^2-T+1=0, T);
```

$$\frac{1 + \sqrt{1 - 4x^2}}{2x^2}, \frac{1 - \sqrt{1 - 4x^2}}{2x^2}$$

```
> f:=1/(2*x^2)*(1-(1-4*x^2)^(1/2));
```

$$f := \frac{1 - \sqrt{1 - 4x^2}}{2x^2}$$

```
> taylor(f, x=0, 20);
```

$$1 + x^2 + 2x^4 + 5x^6 + 14x^8 + 42x^{10} + 132x^{12} + 429x^{14} + 1430x^{16} + O(x^{18})$$

```
> seq(binomial(2*k, k)/(k+1), k=1..8);
```

1, 2, 5, 14, 42, 132, 429, 1430

```
>
```

> Pos:=simplify(Dyck/(1-x\*Dyck));

$$Pos := -\frac{2}{-1 - \sqrt{1 - 4xy + 2x}}$$

> coeftayl(Pos, [x,y]=[0,0], [6,4]);

90

> S:=0:for l from 0 to 6 do for k from 0 to 6 do  
S:=S+coeftayl(Pos, [x,y]=[0,0], [k,l])\*x^k\*y^l od  
od:S;

$1 + x + xy + 20x^6y^2 + 14x^5y^2 + 5x^3y^3 + 2x^2y^2 + x^3 + 28x^5y^3 + x^4 + x^5$   
 $+ x^6 + x^2 + 132x^6y^5 + 2x^2y + 5x^3y^2 + 90x^6y^4 + 42x^5y^5 + 3x^3y$   
 $+ 132x^6y^6 + 4x^4y + 14x^4y^4 + 14x^4y^3 + 5x^5y + 9x^4y^2 + 48x^6y^3$   
 $+ 42x^5y^4 + 6x^6y$

> |

# Automata and rationality

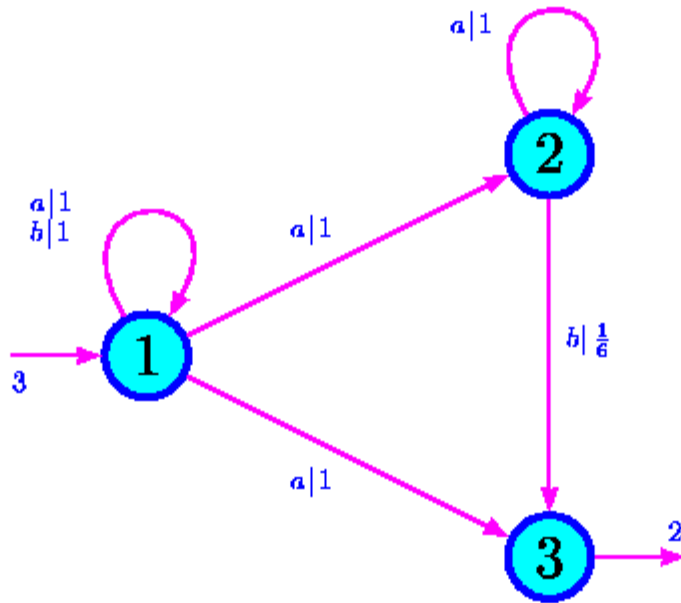


FIG. 1 – Un  $\mathbb{Q}$ -automate  $\mathcal{A}$ .

Le comportement de  $\mathcal{A}$  est :

$$\text{comportement}(\mathcal{A}) = \sum_{a, b \in A} (a + b)^*(6 + a^*b).$$

Un type particulier d'automate à multiplicités est constitué des automates à multiplicités avec des  $\varepsilon$ -transitions.

Un  $k$ - $\varepsilon$ -automate  $\mathcal{A}_\varepsilon$  est un  $k$ -automate sur l'alphabet  $A_\varepsilon = A \cup \{\varepsilon\}$ .

Exemple :

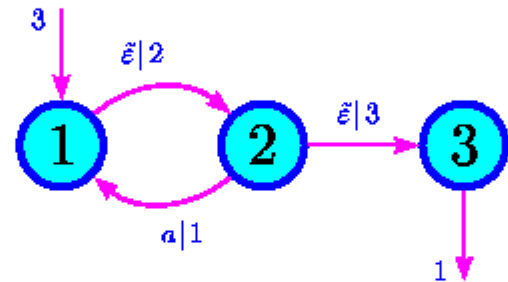
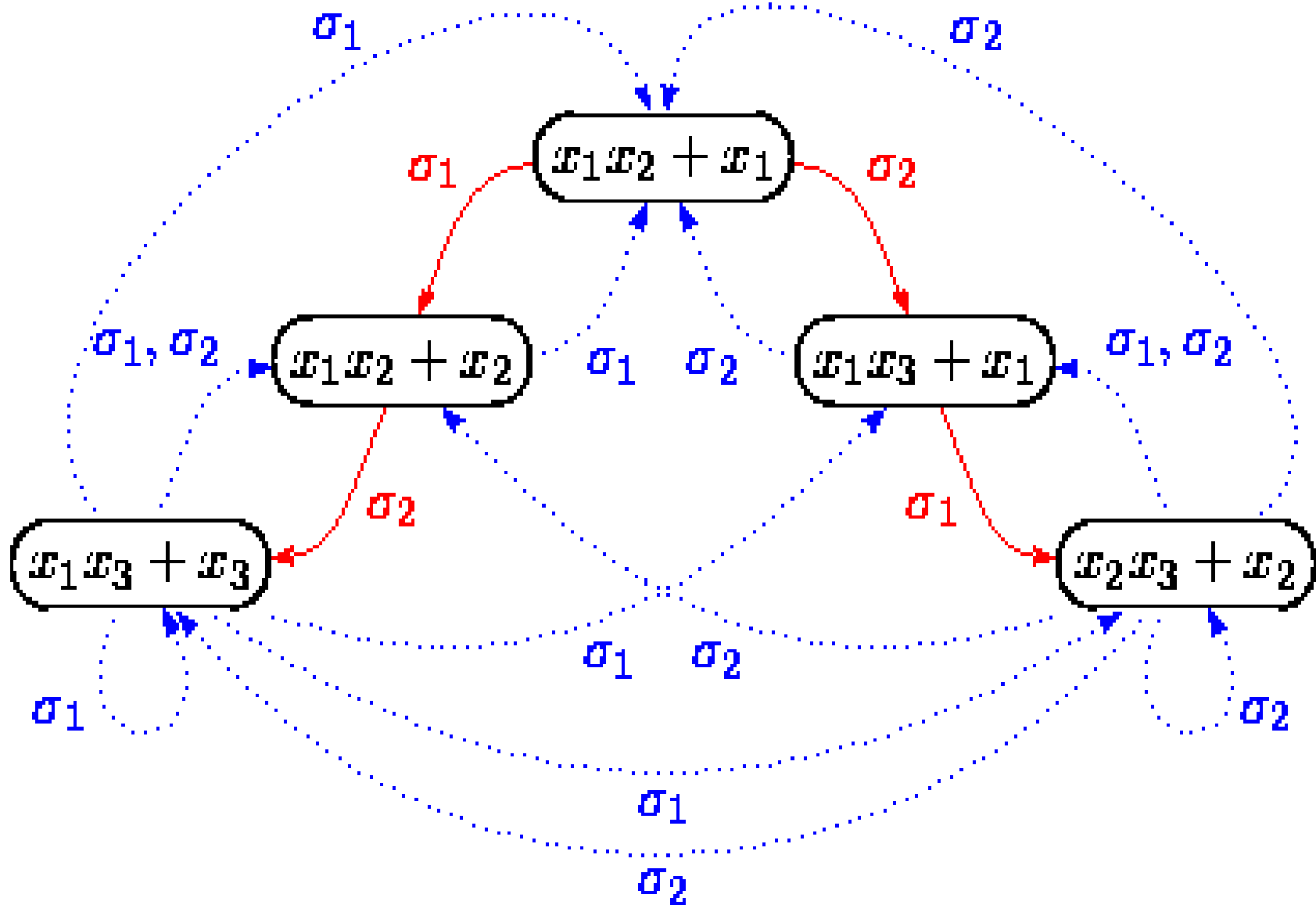


FIG. 2 – Un  $\mathbb{N}$ - $\varepsilon$ -automate  $\mathcal{A}_\varepsilon$ .

$$\text{comportement}(\mathcal{A}_\varepsilon) = 18\varepsilon \left( \sum_{i \in \mathbb{N}} 2^i (a\varepsilon)^i \right) \varepsilon.$$



# A correct implementation of Schelling's model

Problem : If one scans the board, addressing the inhabitants one after one, result is sensitive to the order of scanning.

Solution : Invent a (combinatorial) data structure which adapted to the parallel structure of the moving intentions of the inhabitants.

Problem : If one scans the board, addressing the inhabitants one after one, result is sensitive to the order of scanning.

Solution : Invent a (combinatorial) data structure which adapted to the parallel structure of the moving intentions of the inhabitants --> this must be a global model.

# Combinatorics (mathematics)



Complex  
Systems

Information  
(comp. sci.)

Physics  
(class. quant.)



Thank You