

Dynamic combinatorics, complex systems and applications to physics

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Mathematics

Chaos Theory

Continuous & Discrete Modelisation

Image Processing

Computer Science

Complex Systems Complexity

Business Banking

Computation Techniques

Decision Making

Physics

Artificial Intelligence

Mechatronics

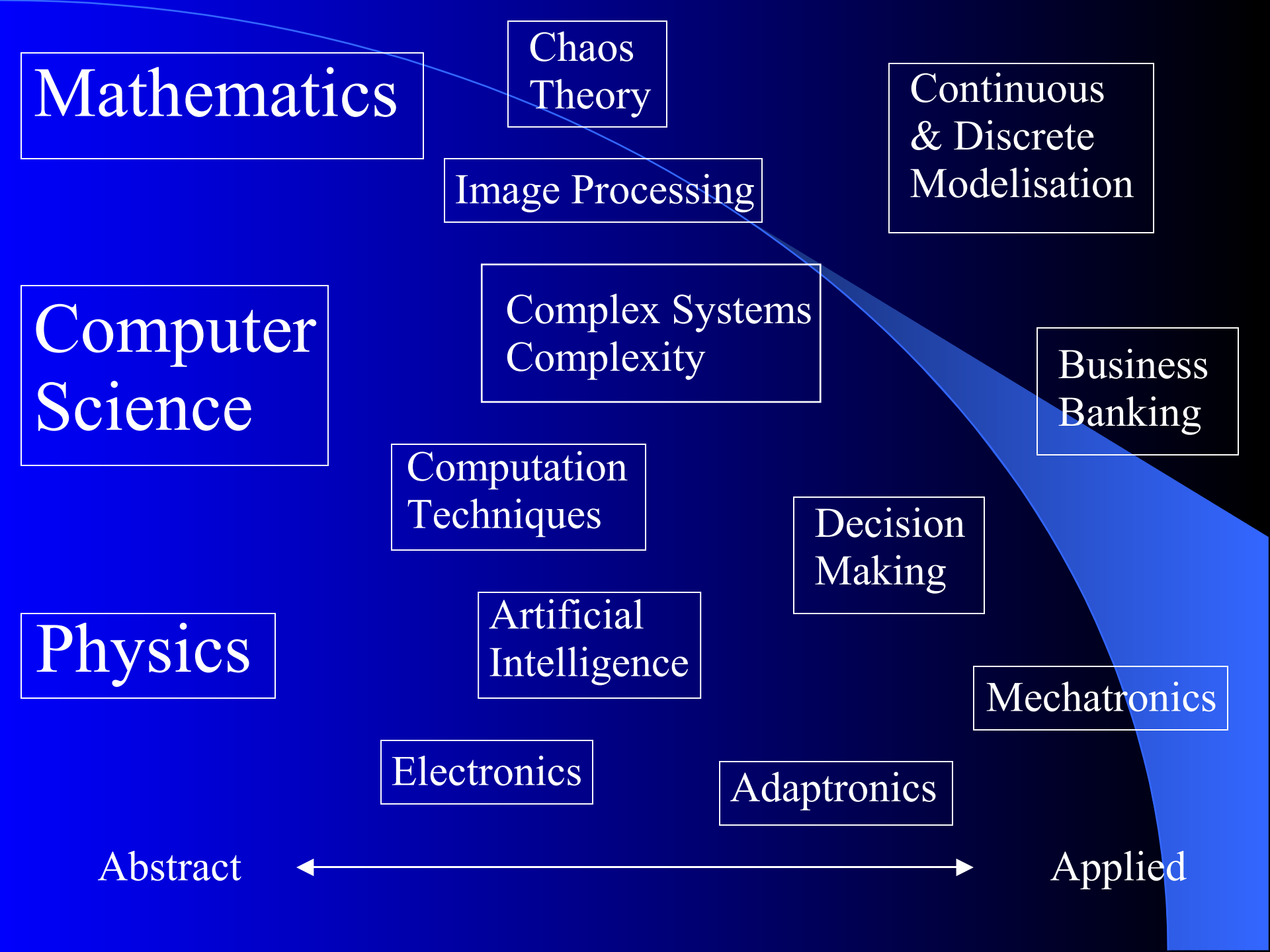
Electronics

Adaptronics

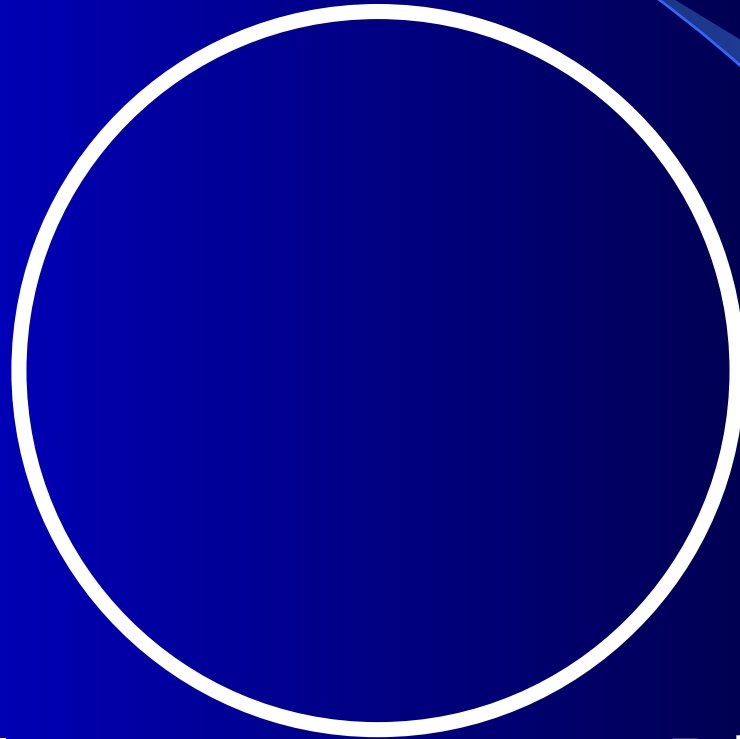
Abstract



Applied



Combinatorics (discrete mathematics)



Information
(computer sci.)

Physics
(classical/quant.)

Combinatorics
(discrete mathematics)



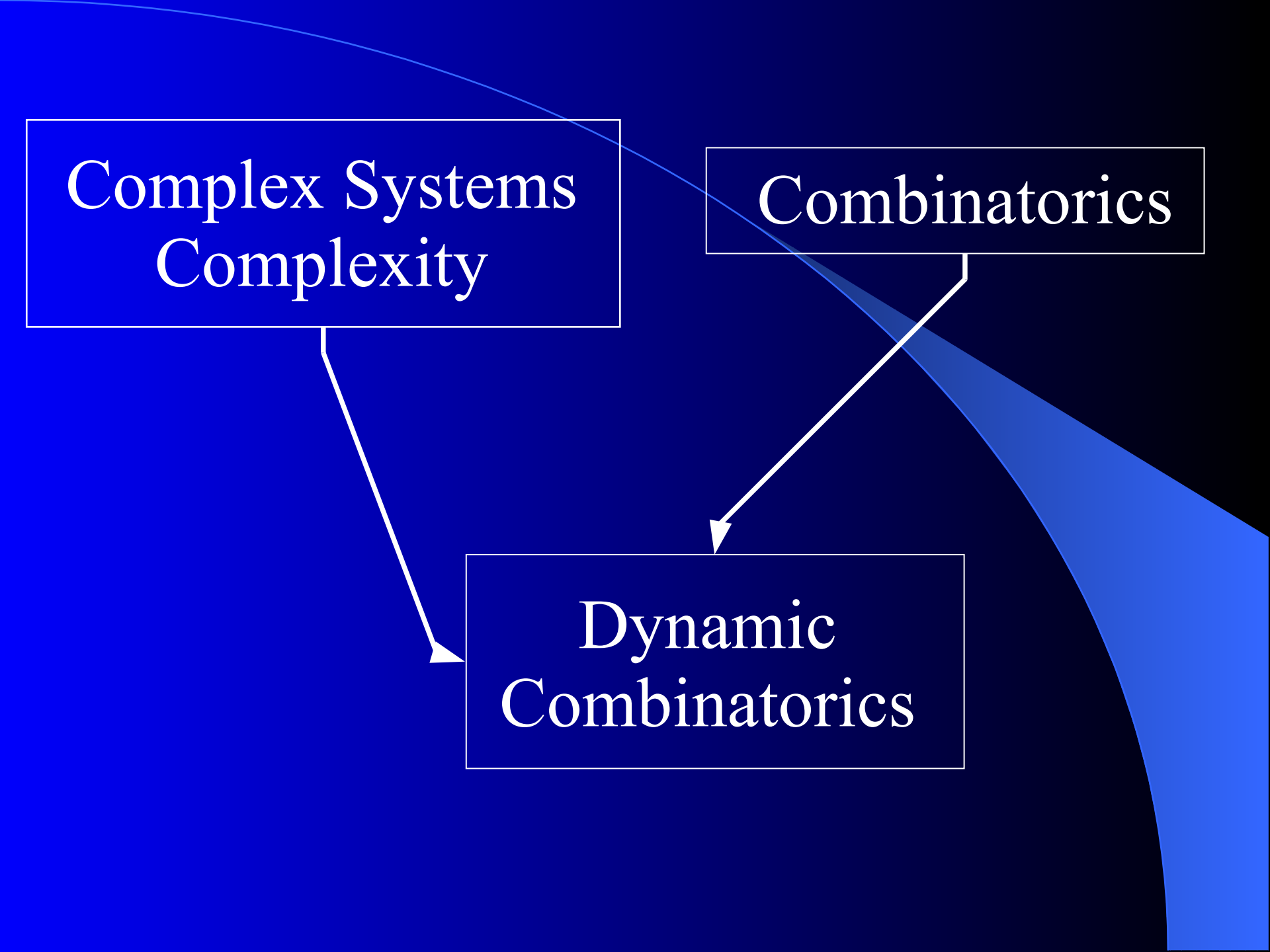
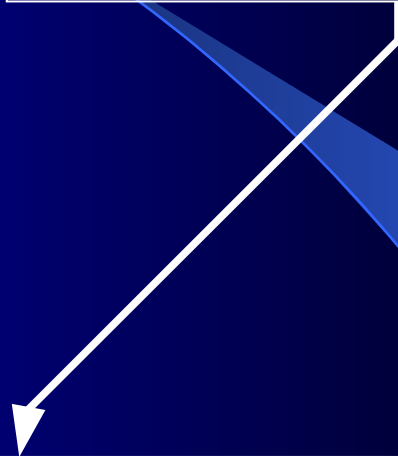
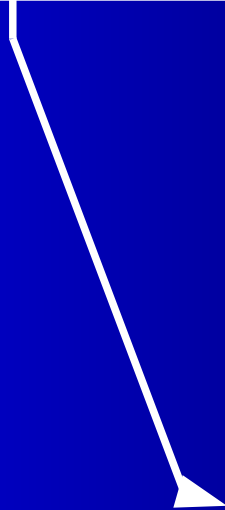
Information
(computer sci.)

Physics
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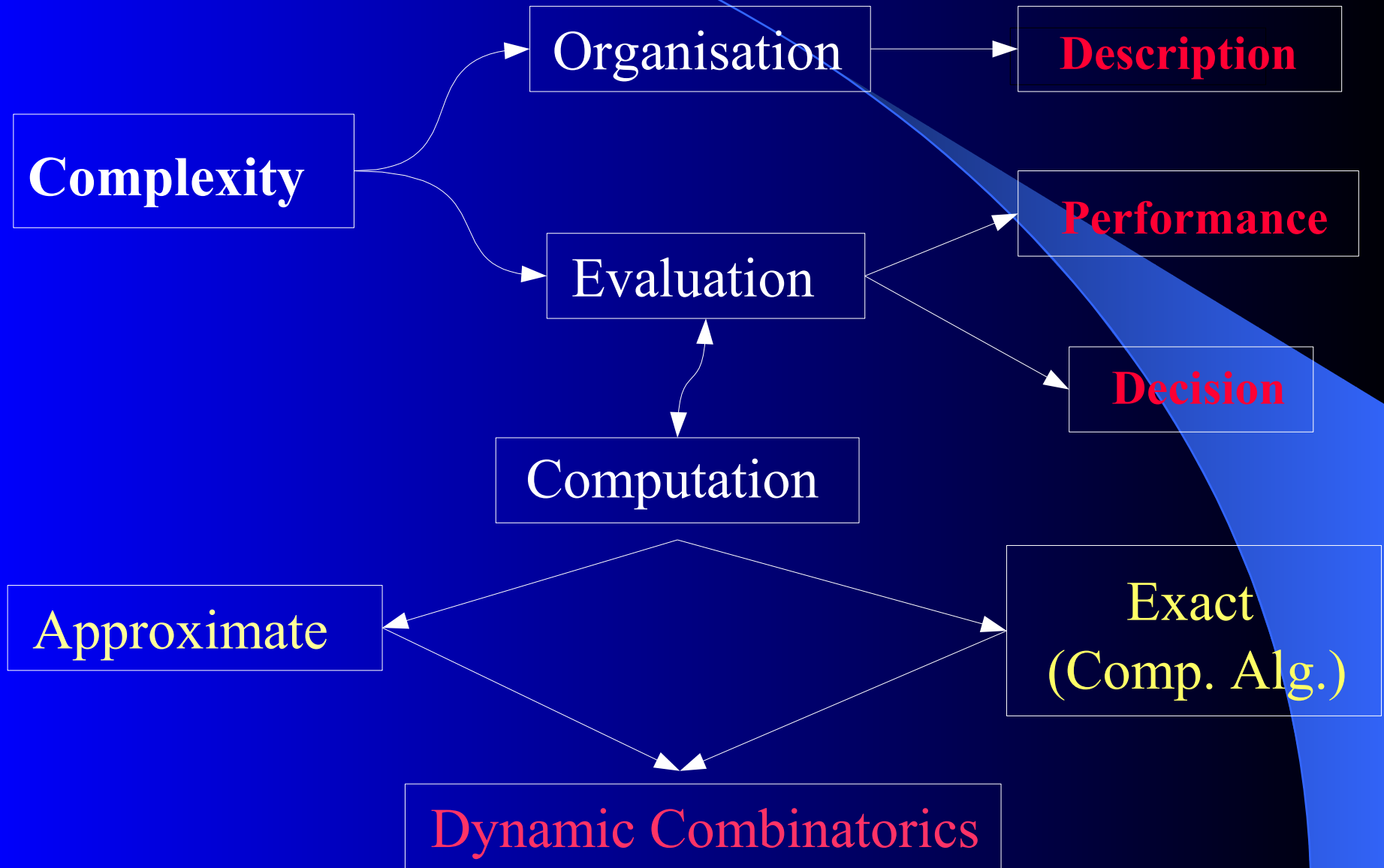
Complex Systems
Complexity

Combinatorics

Dynamic
Combinatorics



Problematics of Dynamic Combinatorics



What is the Legacy ?

Mathematics

- Noncommutative
- Representations
- Formulas, Universal Algebra
- Deformations

Comp. Sciences

- Words
- Automata Transition Structures
- Trees with Operators
- q-analogues

Physics

- Strings of operators
- Fields, Flows, Dynamic Systems (Chaos, Catastrophes)
- Diagrams
- Quantum Groups

Combinatorics

Combinatorics

... on words

- Languages
- Theory of codes
- Automata
- Transition structures
- Grammars
- Transducers
- Rational and algebraic expressions
- ...

enumerative, analytic

- Polyominoes
- Paths (Dyck,...)
- Configurations
- q-grammars
- Generating Functions
- Continued Fractions (mono, multivariate,..)
- Orthogonal Polynomials
- ...

algebraic

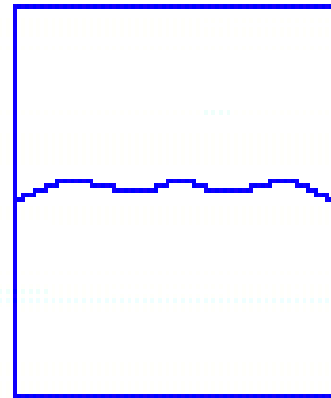
- Non commutative Continued fractions
- Representations of groups and deformations
- Quantum Groups
- Functors
- Characters
- Special Functions
- ...

A first example . . .



P

$n \rightarrow n+1$
i = one unit



R

$n = n+1$

- $a_1 = 0$
- $a_{i+1} \leq a_i + 1$

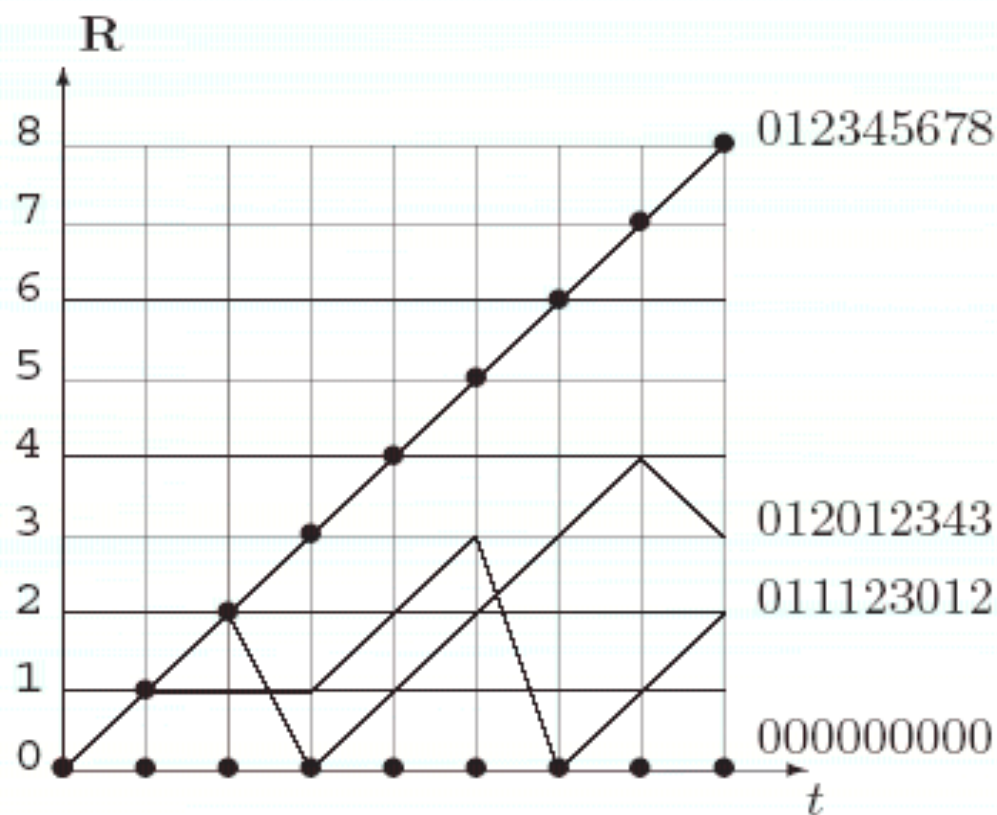


Figure 4.2: Maximal, minimal (dotted) and two intermediate trajectories. Their codes are on the right.

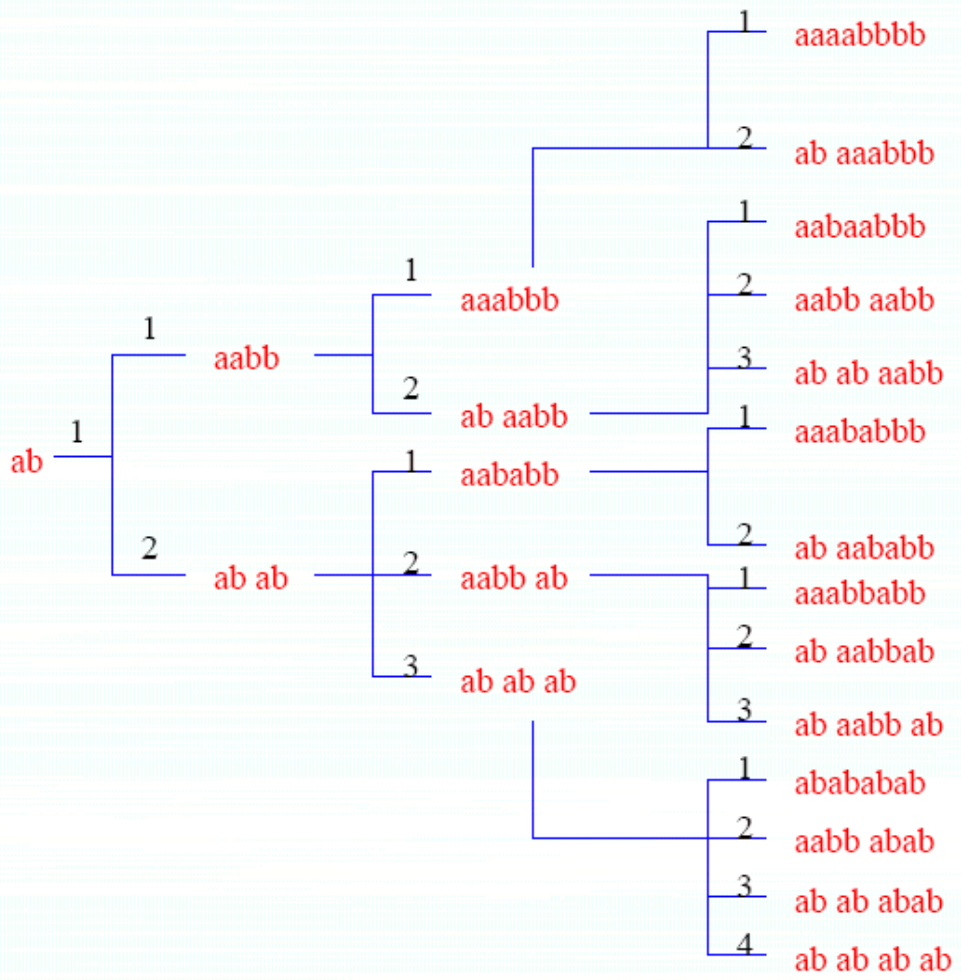


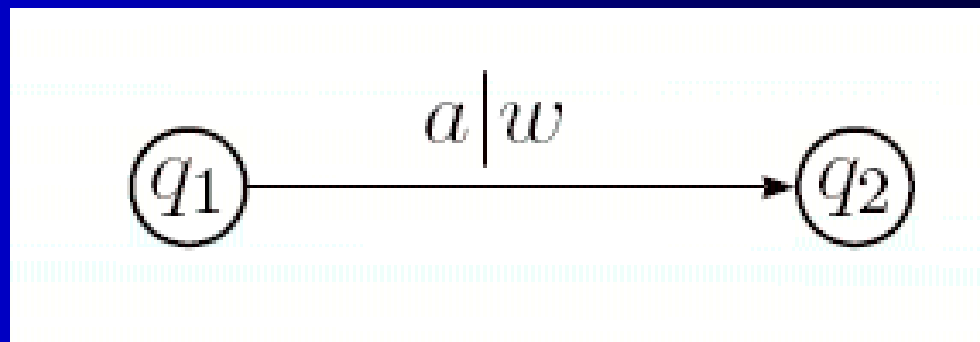
Figure 2.9: Dyck words of length $2n$ with k factors

We will return to the Dyck paths later on.
For the moment let us define what is a transition structure.

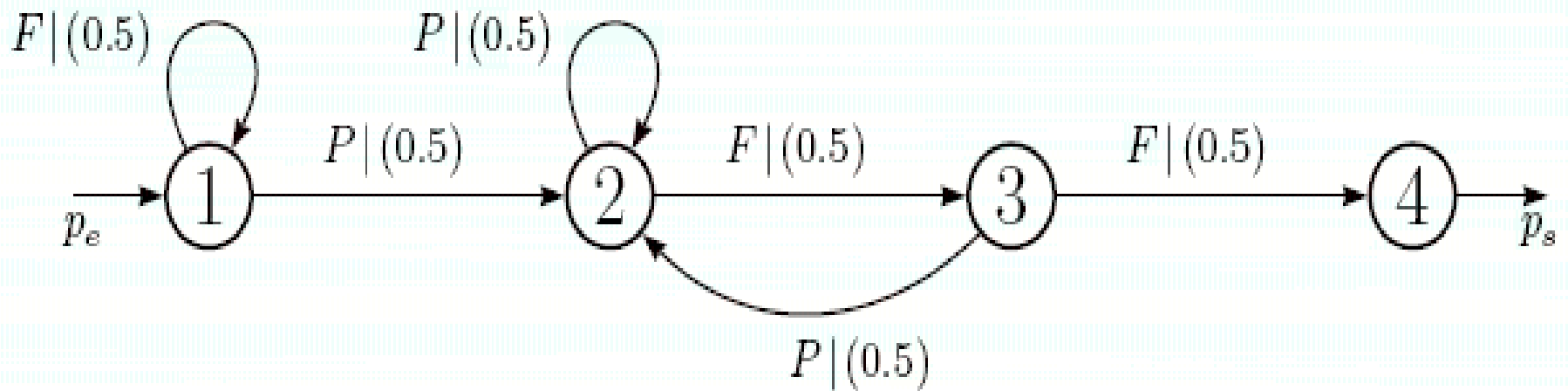
Definition (transition structure) : It is a graph (finite or infinite) with its arcs marked with pairs

(command letter | coefficient)

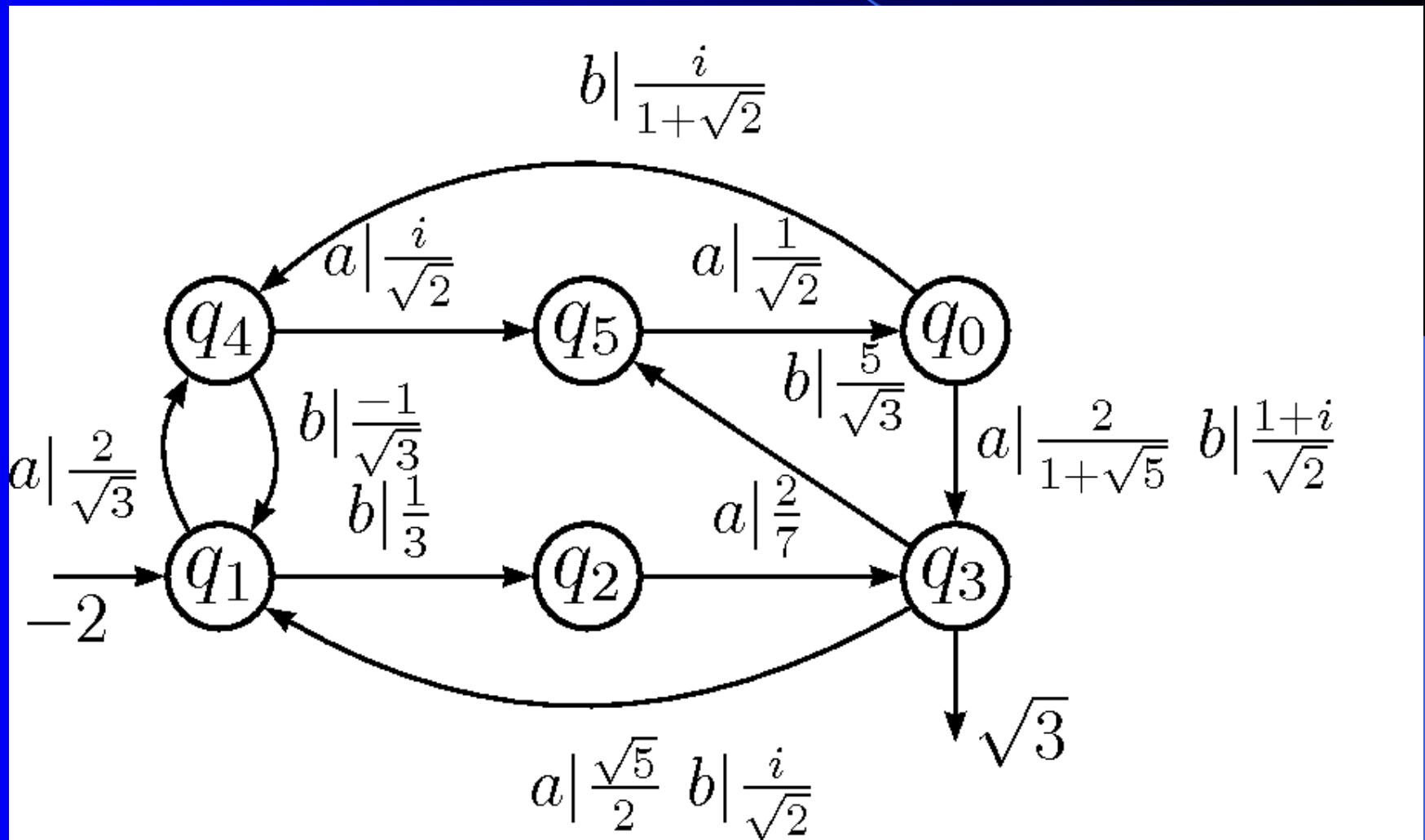
Examples : Prisoner's dilemma, Markov chains, classical engineering.



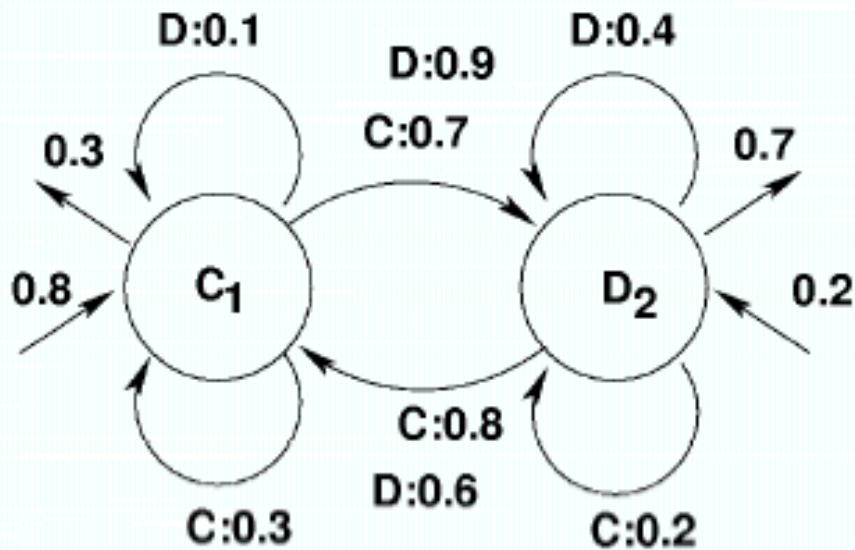
Example : A Markov chain generated by a game.



Example : An automaton generated by arbitrary transition coefficients.



Example of Probabilistic Automaton



LINEAR REPRESENTATION

$\begin{bmatrix} 0.8 & 0.2 \end{bmatrix}$ input vector

	1	2		1	2
1	0.3	0.7	0.1	0.9	1
2	0.8	0.2	0.6	0.4	2

$M(C)$

$M(D)$

$\begin{bmatrix} 0.3 \\ 0.7 \end{bmatrix}$ output vector

Behaviour of an Automaton and how to compute it effectively

An automaton is a **machine** which takes a string (sequence of letters) and returns a **value**.

This value is computed as follows :

- 1) The **weight** of a path is the product of the weights (or coefficients) of its edges

- 2) The **label** of a path is the product (concatenation) of the labels of its edges

Behaviour ... (cont'd)

3) The **behaviour** between two states « r,s » w.r.t. A word « w » is the product of

3a) the ingoing coefficient of the first state (here « r ») by

3b) the sum of the weights of the paths going from « r » to « s » with label « w » by

3c) the outgoing coefficient of the second state (here « s »)

Behaviour ... (cont'd)

4) The **behaviour** of the automaton under consideration w.r.t. a word « w » is then the sum of all the behaviours of the automaton between two states « r,s » for all possible pairs of states.

Behaviour ... (cont'd)

There is a simple formula using the linear representation. The **behaviour** of an automaton with linear representation (I, M, T) is the product

$$IM(w)T$$

where $M(w)$ is the canonical extension of M to the strings.

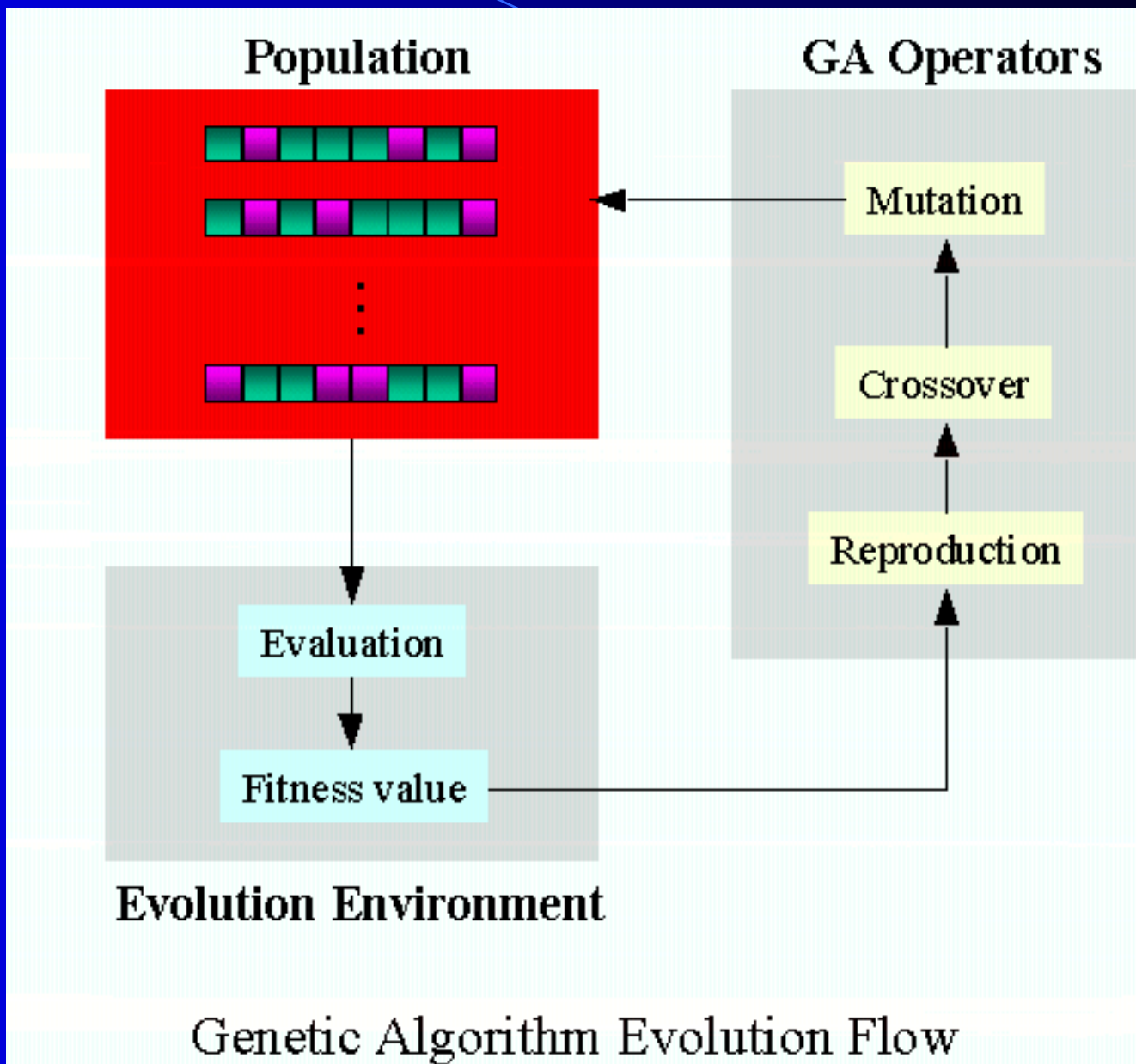
$$M(a_1 a_2 \dots a_n) = M(a_1) M(a_2) \dots M(a_n)$$

Behaviour ... (end)

The behaviour, as a function on words belongs to the rational class. If time permits, we will return to its complete calculation as a **rational expression** and the problem of its algorithmic evaluation by means of special cancellation operators. Linear representations can also be used to compute **distances between automata**.

Example -> use of genetic algorithms to control indirect (set of) parameters : the spectrum of a matrix.

Genetic algorithms : general pattern



Genetic algorithms : implementation

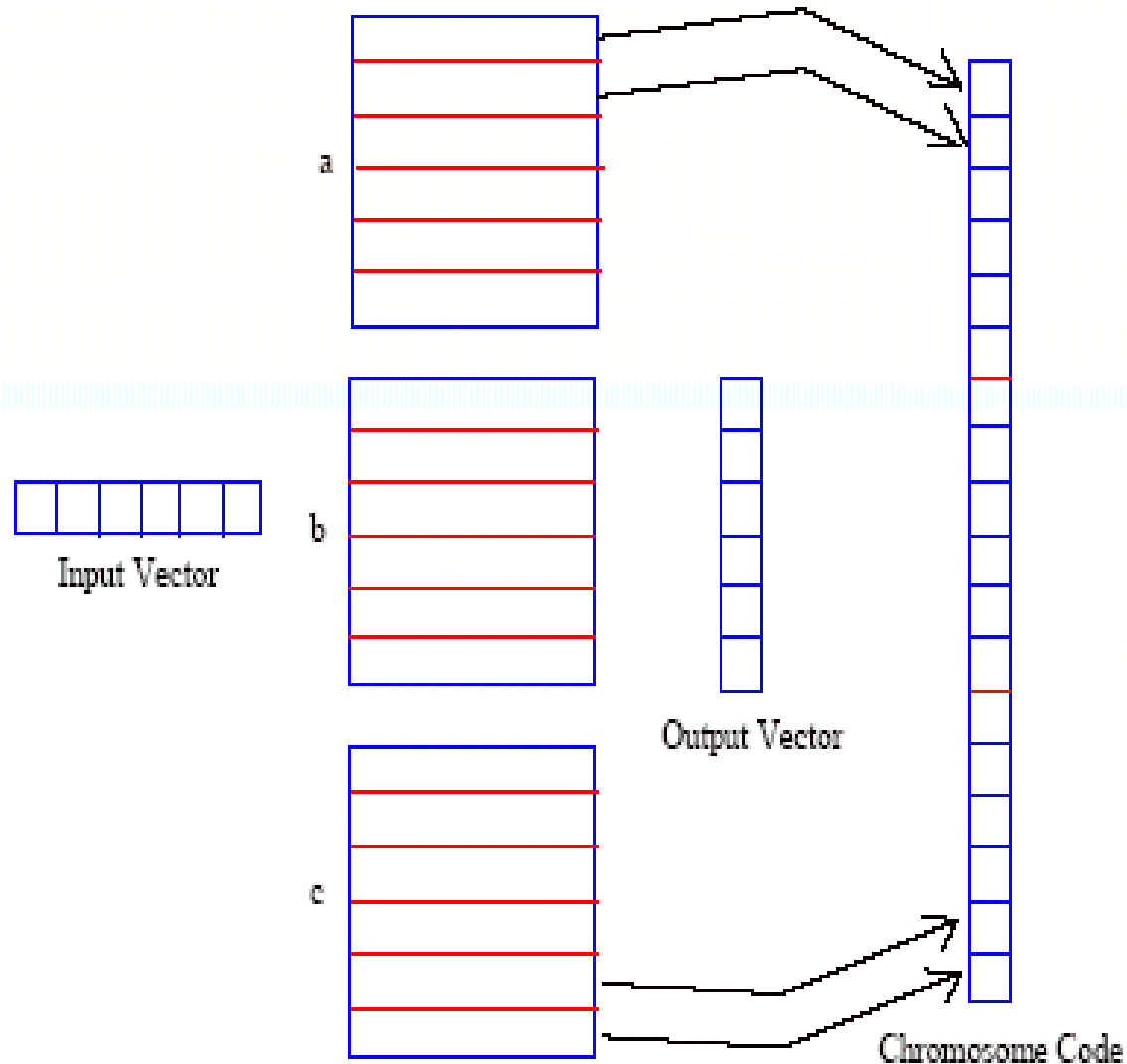


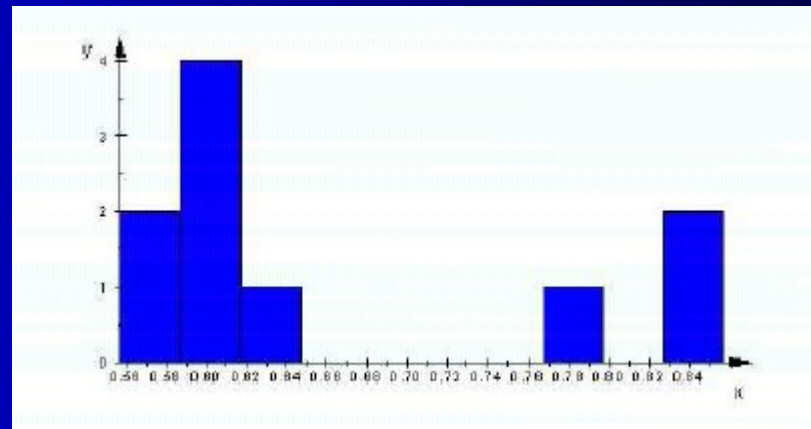
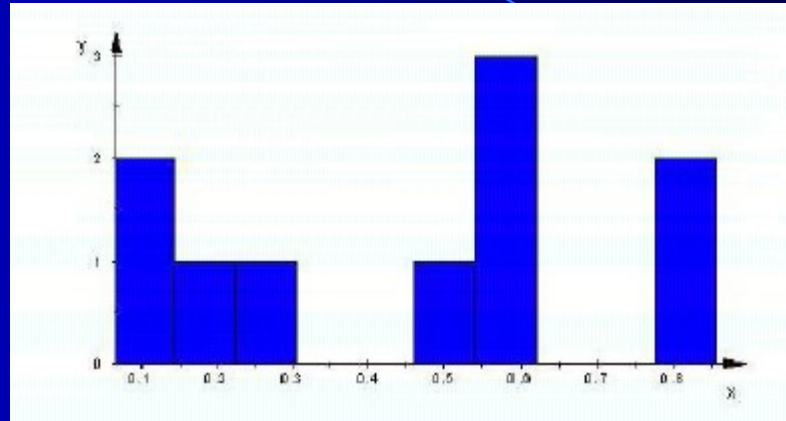
Figure 4.13: Chromosome code

Genetic algorithms : implementation

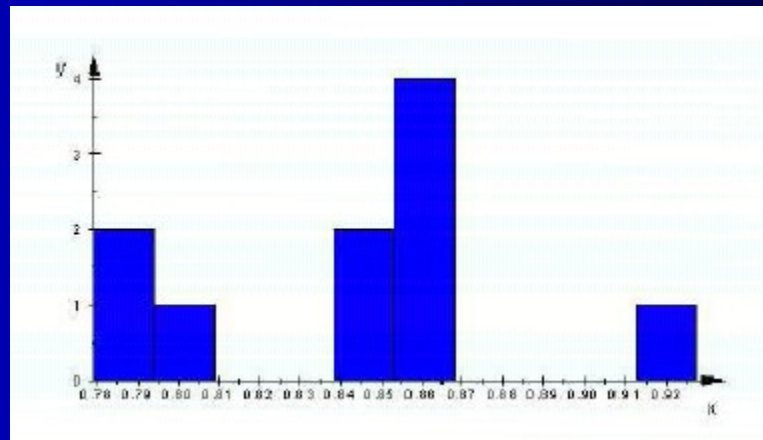
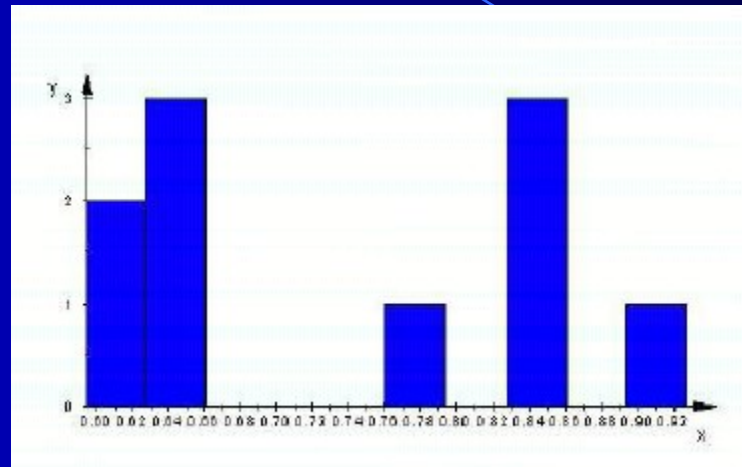
Below, the results of an experiment aiming to control the second greatest eigenvalue of the transfer matrix of a population of probabilistic automata.

- The fitness function of each automaton corresponds to the second greatest eigenvalue (in module). The first being, of course, of value 1.

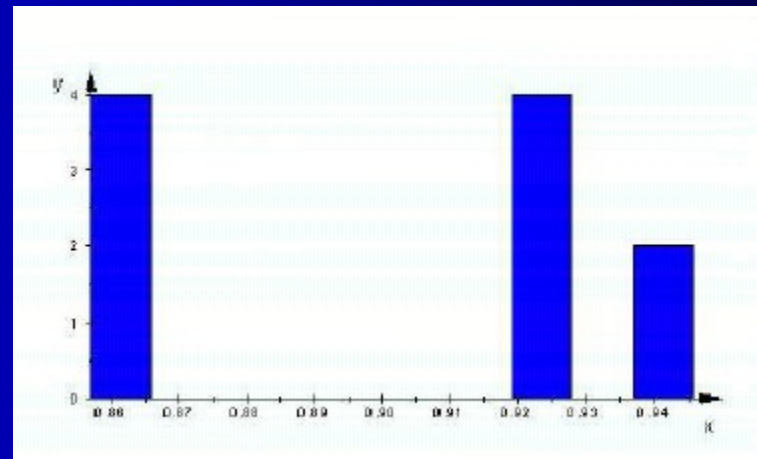
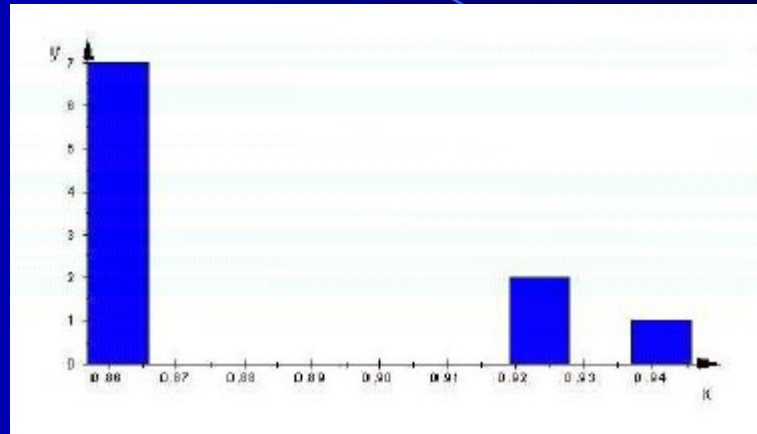
Genetic algorithms ; results



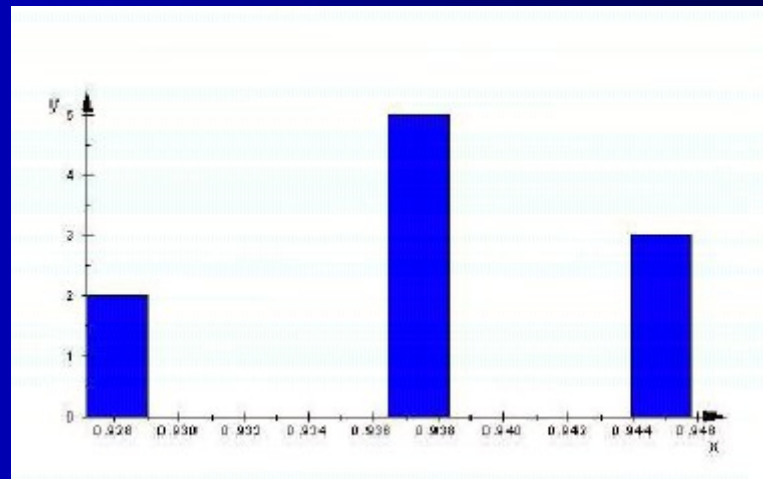
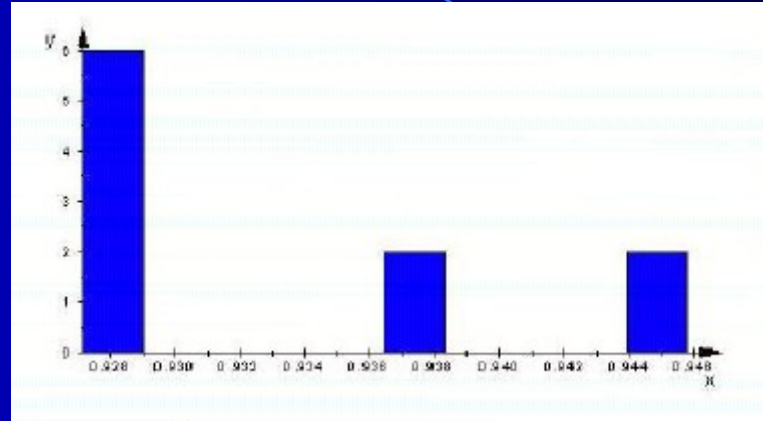
Genetic algorithms ; results



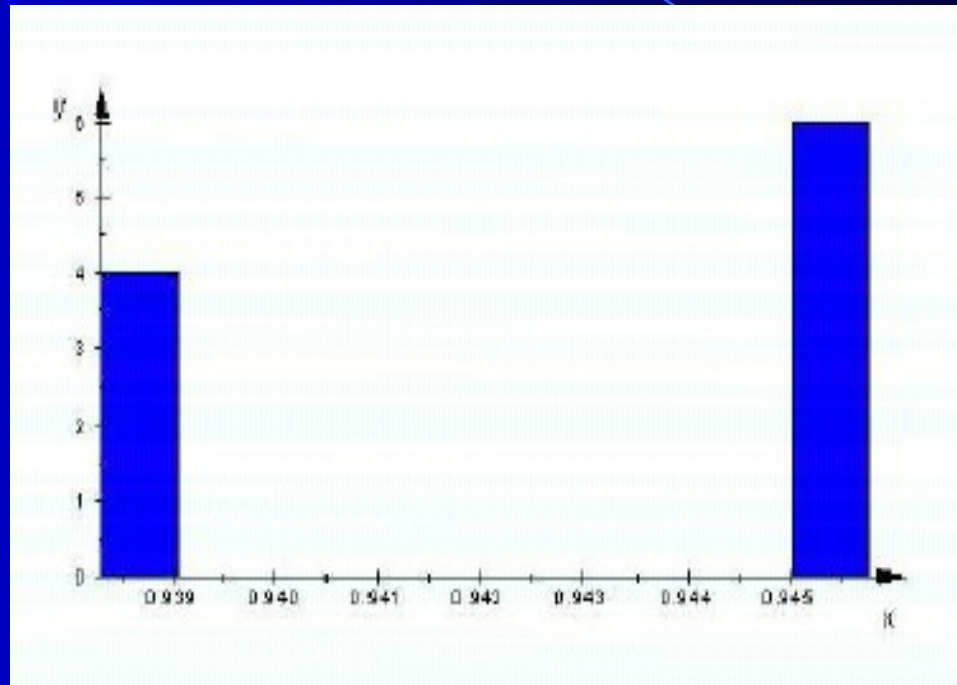
Genetic algorithms ; results



Genetic algorithms ; results



Genetic algorithms ; results



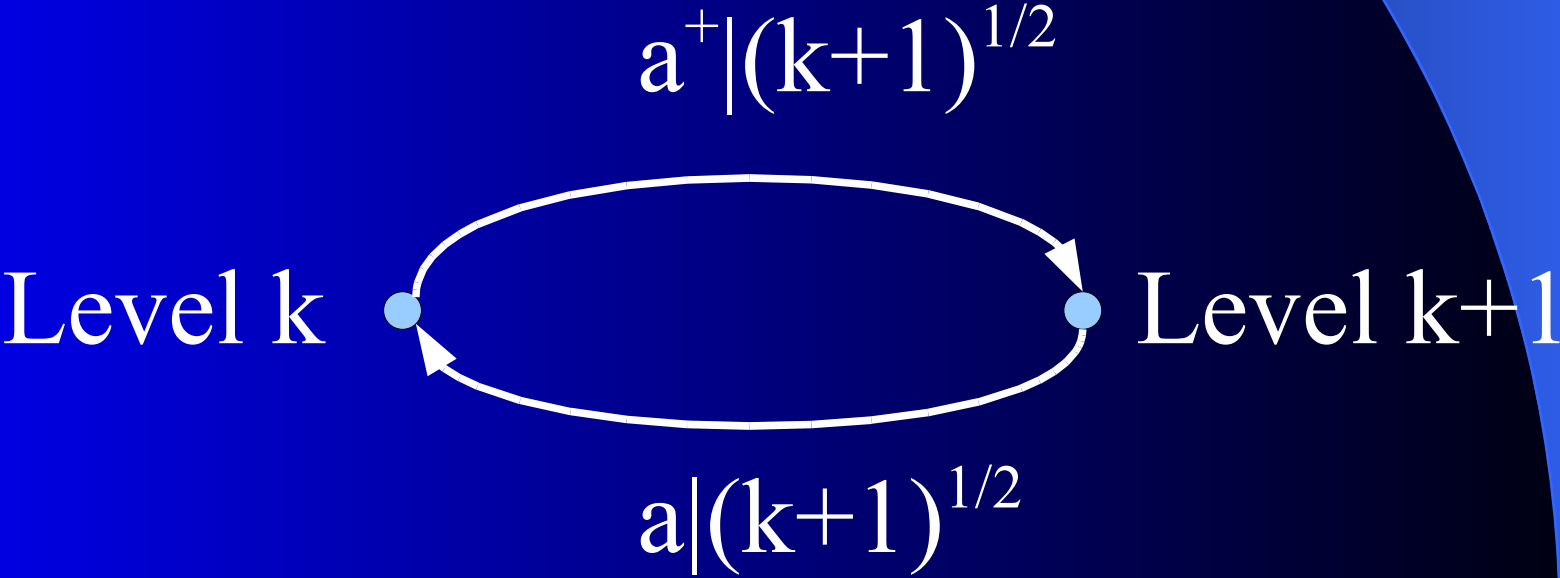
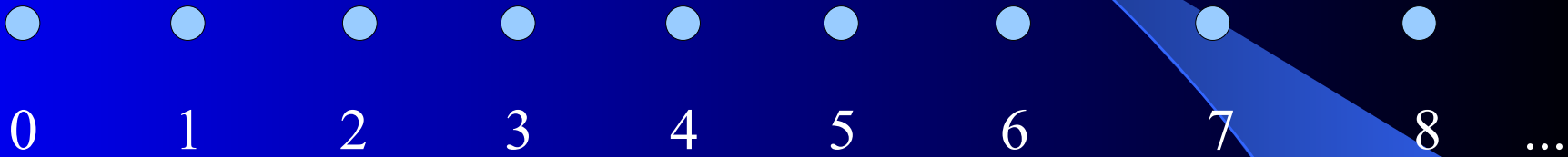
Final result : the population is rendered homogeneous

General transition systems



- Automata (finite number of edges)
- Sweedler's duals (physics, finite number of states)
- Representations
- Level systems (Quantum Physics)
- Markov chains (prob. automata when finite)

Example in Physics : annihilation/creation operators



The (classical, for bosons) normal ordering problem goes as follows.

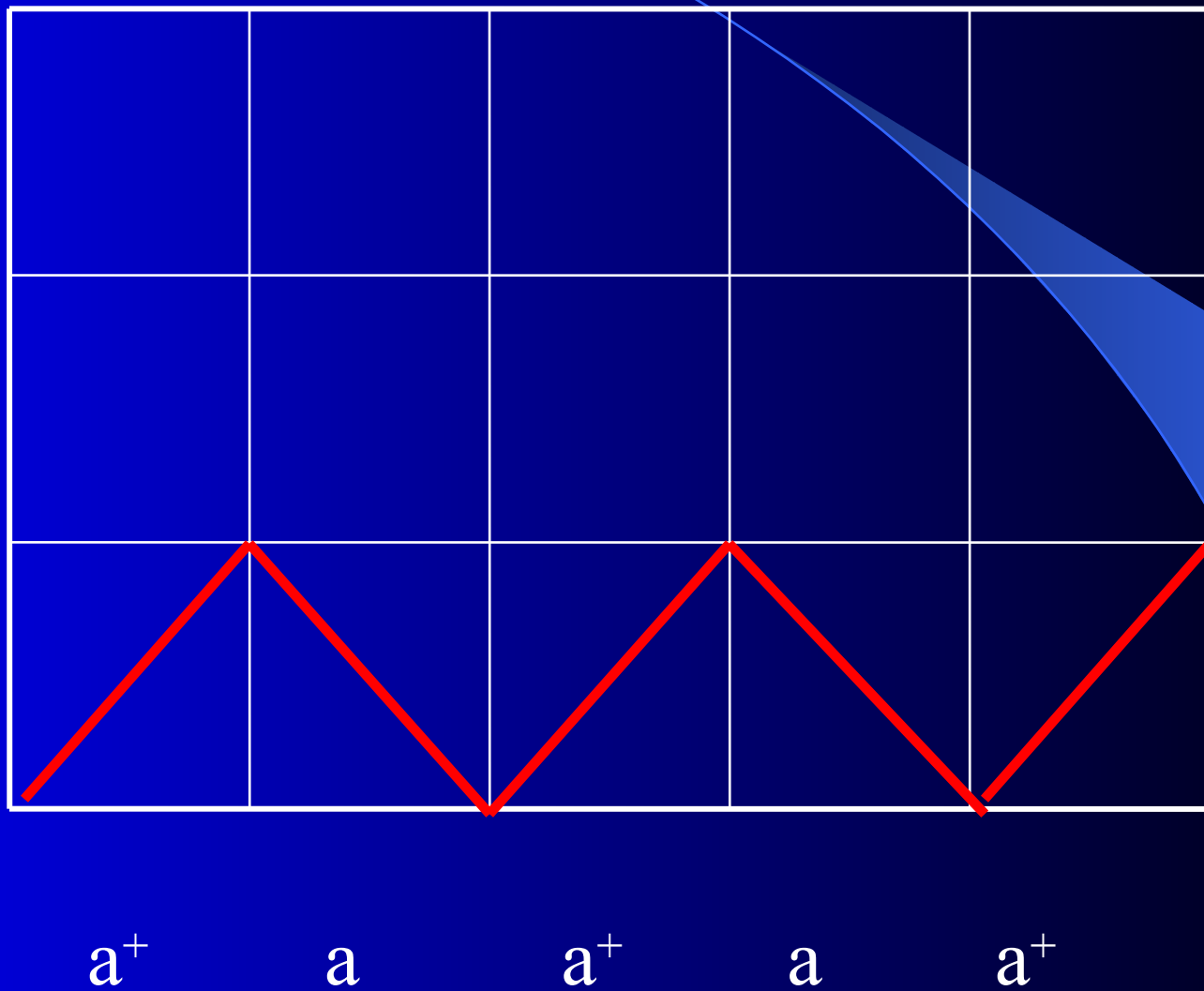
- Weyl (two-dimensional) algebra defined as
 $\langle a^+, a ; [a, a^+] = 1 \rangle$
- Known to have no (faithful) representation by bounded operators in a Banach space.

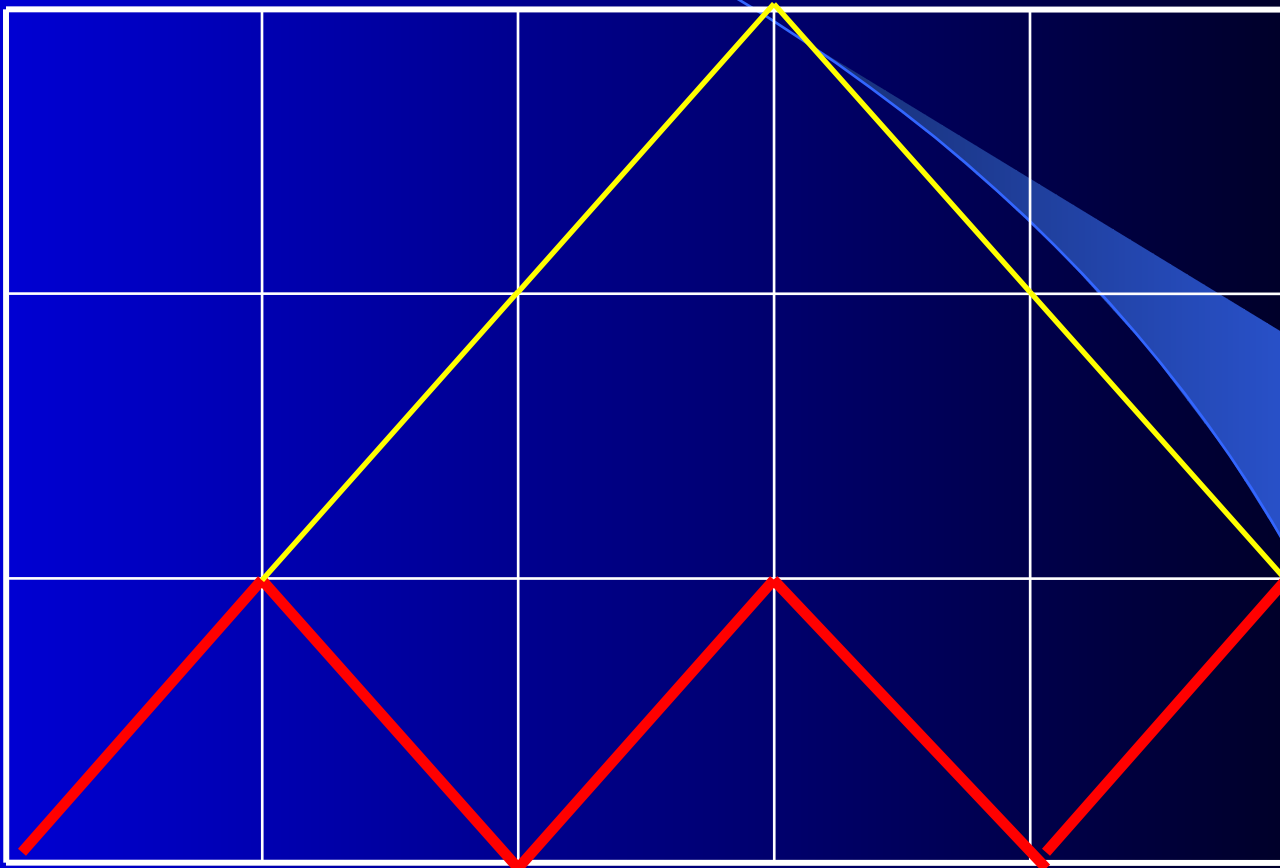
There are many « combinatorial » (faithful) representations by operators. The most famous one is the Bargmann-Fock representation

$$a \rightarrow d/dx ; a^+ \rightarrow x$$

where a has degree -1 and a^+ has degree 1 .

Example with $\Omega = a^+ a a^+ a a^+$





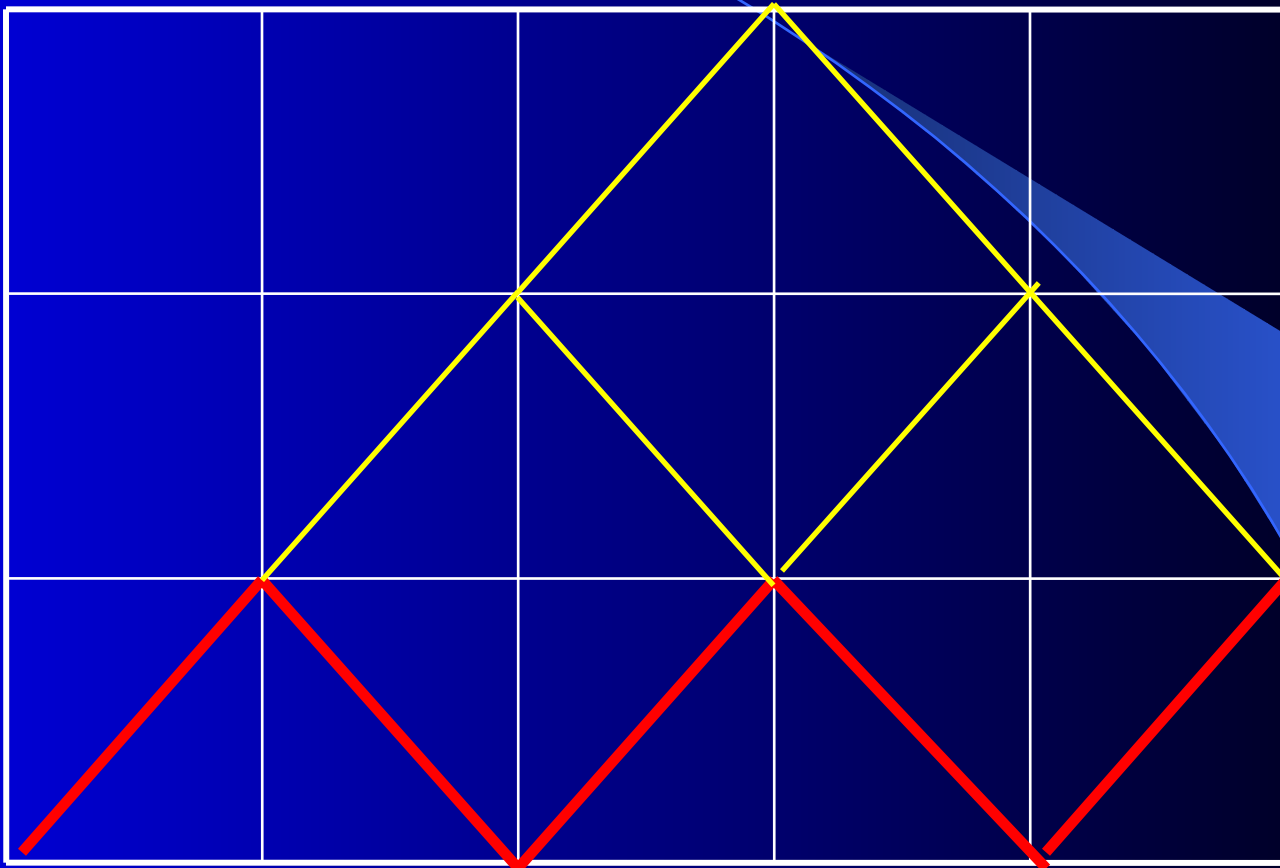
a^+

a

a^+

a

a^+



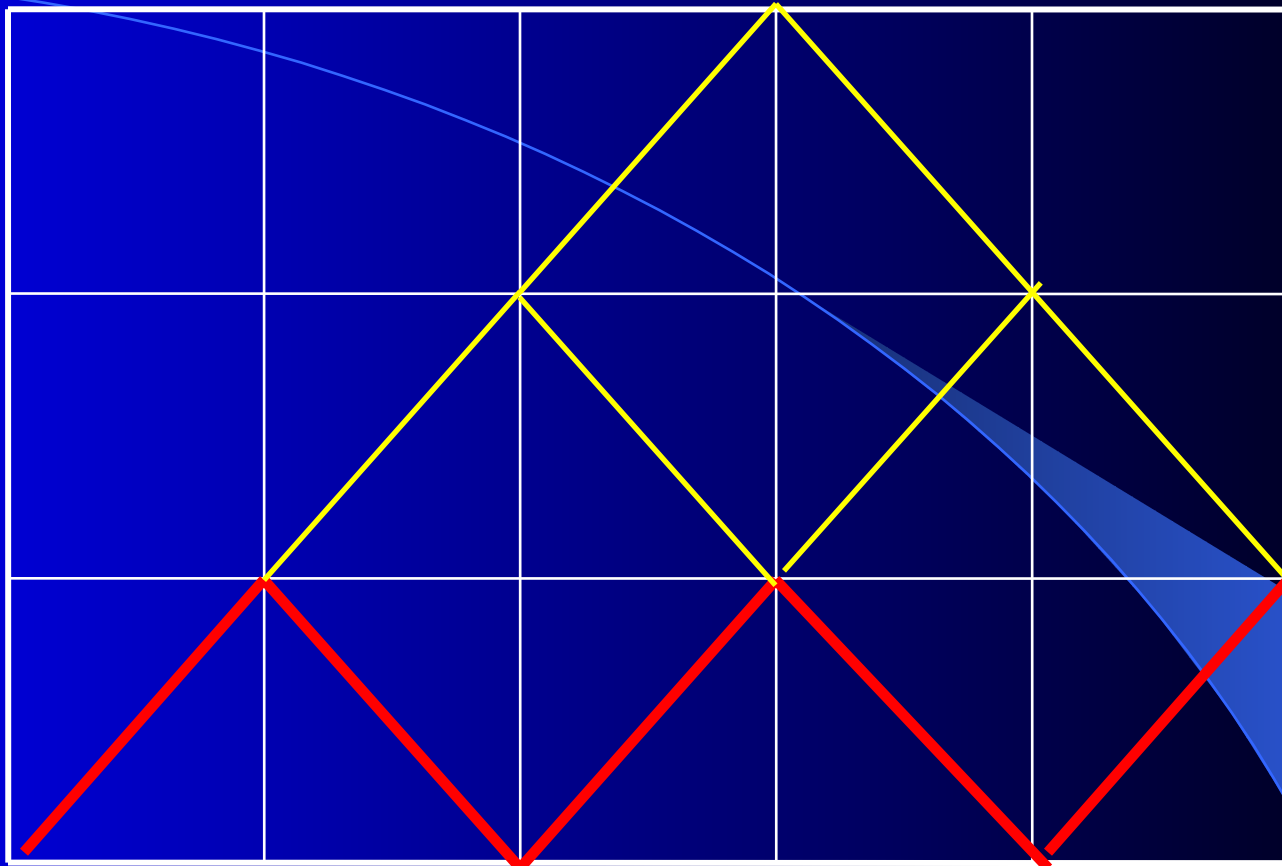
a^+

a

a^+

a

a^+



$a^+ \quad a \quad a^+ \quad a \quad a^+$

$$a^+aa^+aa^+ = 1 a^+a^+a^+aa + 3 a^+a^+a + 1 a^+$$

Through Bargmann-Fock representation

$$a \rightarrow d/dx ; a^+ \rightarrow x$$

Operators who have only one annihilation have exponentials who act as one-parameter groups of substitutions.

One can thus use computer algebra to determine their generating function.

For example, with

$$\Omega = a^{+2}a a^+ + a^+a a^{+2}$$

the computation reads



One parameter group by $f(v(u(x)+\lambda))$; v is reciprocal of u

```
> T1(lambda, x) := (-4 * (-1 / (4 * x^2) + lambda)) ^ (-1/2);
```

$$T1(\lambda, x) := \frac{1}{\sqrt{\frac{1}{x^2} - 4\lambda}}$$

We suppose $x > 0$

```
> T1 := (lambda, x) -> x / ((1 - 4 * lambda * x^2) ^ (1/2));
```

$$T1 := (\lambda, x) \rightarrow \frac{x}{\sqrt{1 - 4\lambda x^2}}$$

Checking the tangent vector

```
> subs(lambda=0, diff(T1(lambda, x), lambda));
```

$$2x^3$$

... and the one-parameter group property

```
> simplify(T1(lambda1, T1(lambda2, x)) - T1(lambda1+lambda2, x));
```

$$0$$

And the action of $\exp(\lambda \omega)$ on $[f(x)]$ is

$$\begin{aligned}
 U_\lambda (f) &= x^{-\frac{3}{2}} f(s_\lambda (x)).(s_\lambda (x))^{\frac{3}{2}} \\
 &= \sqrt[4]{\frac{1}{(1-4\lambda x^2)^3}} f\left(\sqrt{\frac{x^2}{1-4\lambda x^2}}\right)
 \end{aligned}$$

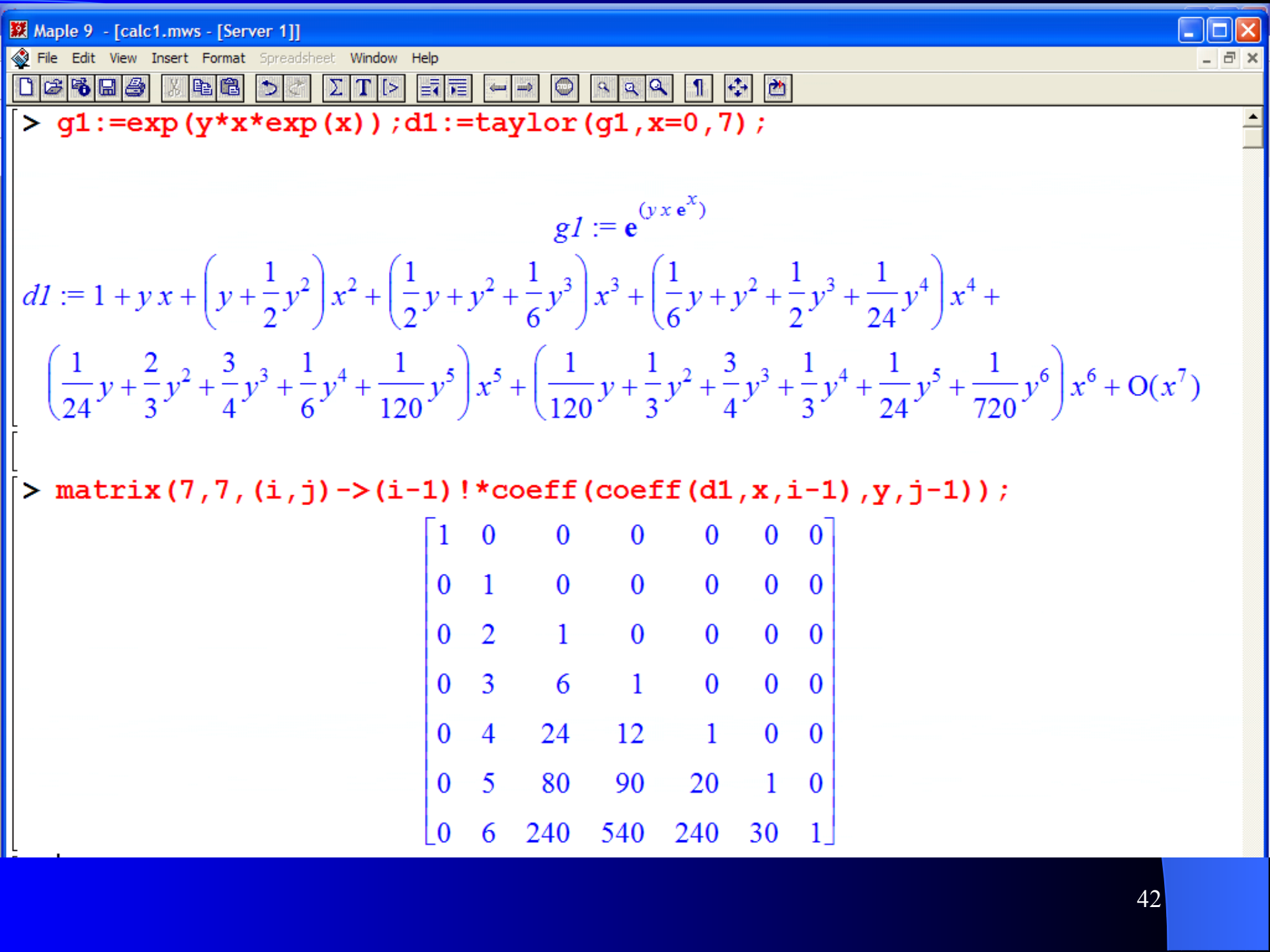
which explains the prefactor. Again we can check by computation that the composition of (U_λ) s amounts to simple addition of parameters !!

Now suppose that $\exp(\lambda \omega)$ is in normal form.

In view of Eq1 (slide 15) we must have

$$\exp(\lambda \omega) = \sum_{n \geq 0} \frac{\lambda^n \omega^n}{n!} = \sum_{n \geq 0} \frac{\lambda^n}{n!} x^{ne} \sum_{k=0}^{ne} S_\omega(n, k) x^k \left(\frac{d}{dx}\right)^k$$

So, using this new technique, one can compute easily the coefficients of the matrix giving the normal forms.



For these one-parameter groups and conjugates of vector fields

G. H. E. Duchamp, K.A. Penson, A.I. Solomon, A. Horzela and P. Blasiak,

One-parameter groups and combinatorial physics,

Third International Workshop on Contemporary Problems in Mathematical Physics (COPROMAPH3), Porto-Novo (Benin), November 2003. arXiv : quant-ph/0401126.

For the Sheffer-type sequences and coherent states

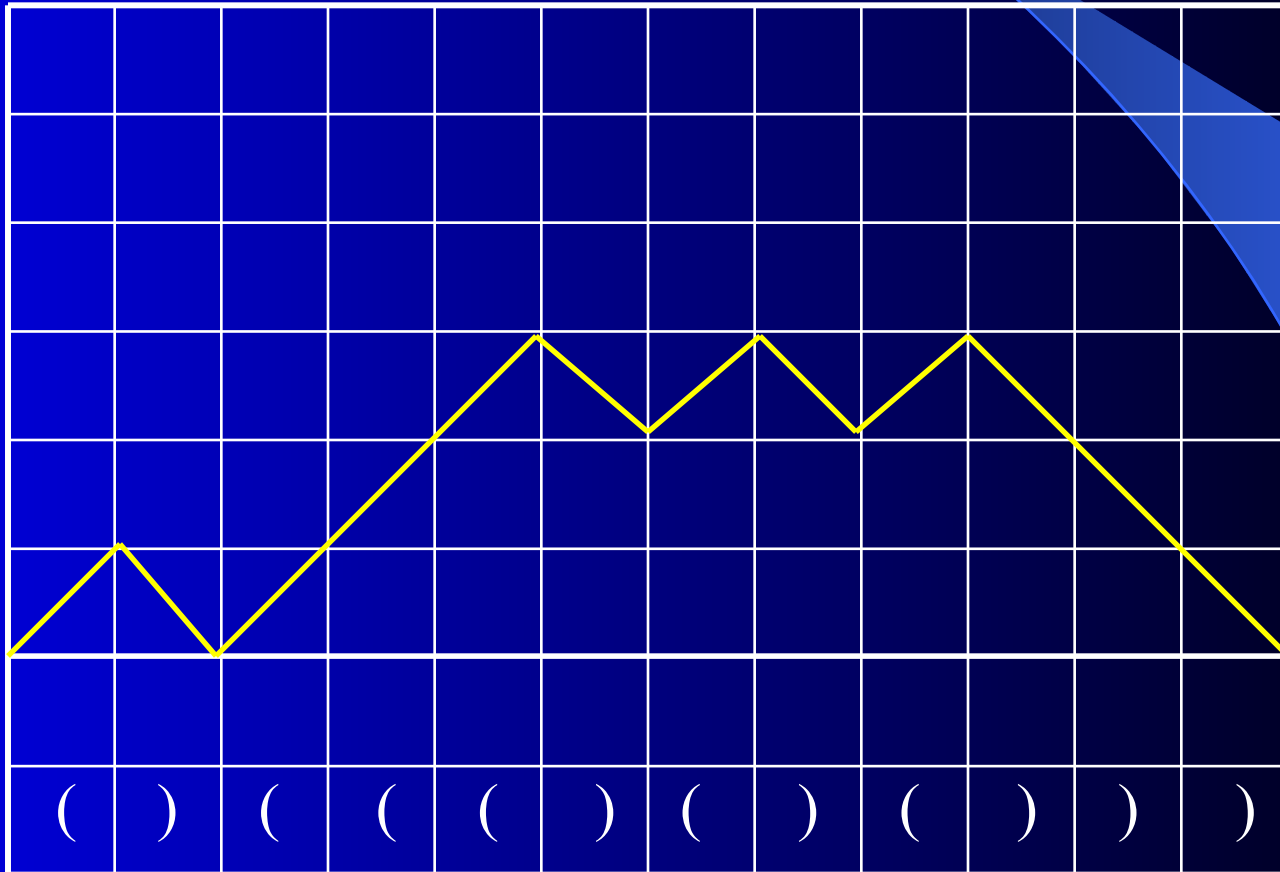
K A Penson, P Blasiak, G H E Duchamp, A Horzela and A I Solomon,

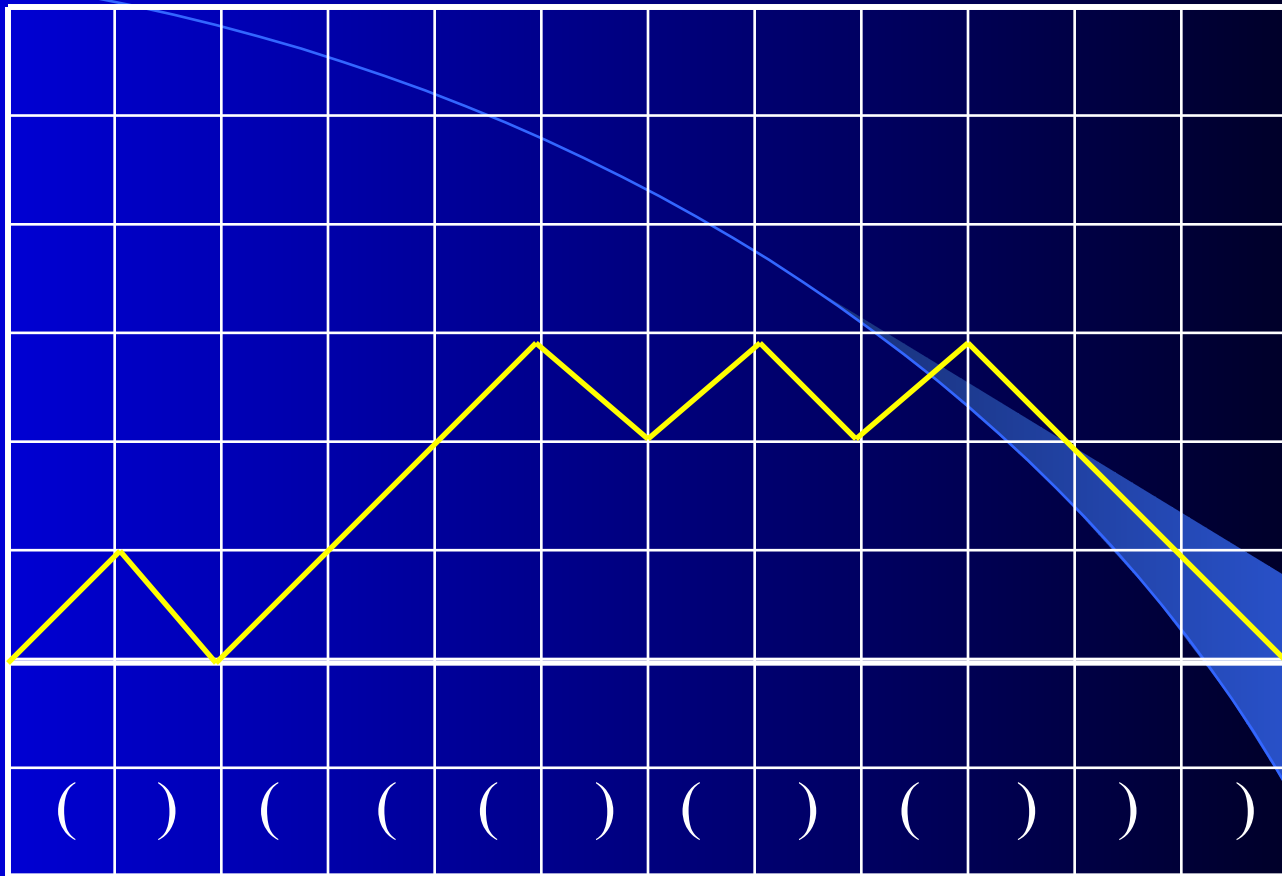
Hierarchical Dobinski-type relations via substitution and the moment problem,

J. Phys. A: Math. Gen. 37 3457 (2004) arXiv : quant-ph/0312202

A second application : Dyck paths

(systems of brackets, trees, physics, ...)

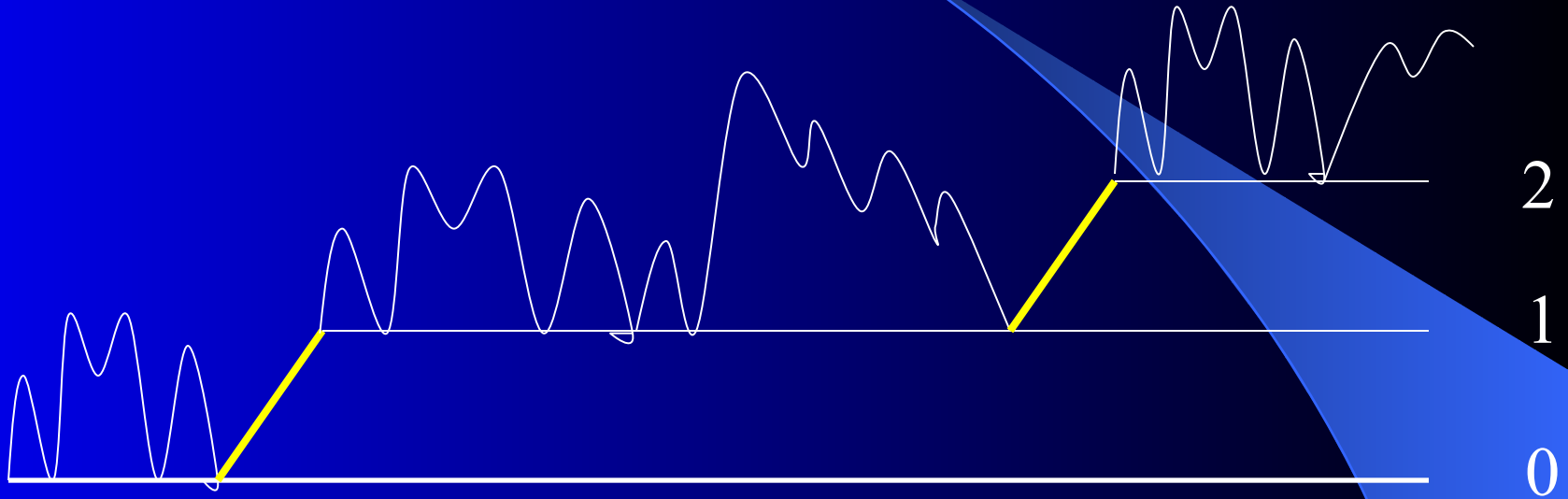




Equation : $D = \text{void} + (D) D \dots$ one counts strings using an « x » by bracket and one finds $T(x) = x^0 + x^2 T^2(x)$ which can be solved by elementary methods ...

$$x^2 T^2 - T + 1 = 0 \quad \text{Variable : } T \quad \text{Parameter : } x$$

Change of level (physics)



$$\text{Positifs} = D(aD)^*$$

$$Pos := \frac{Dyck}{1 - x Dyck}$$

```
> solve(x^2*T^2-T+1=0, T);
```

$$\frac{1 + \sqrt{1 - 4x^2}}{2x^2}, \frac{1 - \sqrt{1 - 4x^2}}{2x^2}$$

```
> f:=1/(2*x^2)*(1-(1-4*x^2)^(1/2));
```

$$f := \frac{1 - \sqrt{1 - 4x^2}}{2x^2}$$

```
> taylor(f, x=0, 20);
```

$$1 + x^2 + 2x^4 + 5x^6 + 14x^8 + 42x^{10} + 132x^{12} + 429x^{14} + 1430x^{16} + O(x^{18})$$

```
> seq(binomial(2*k, k)/(k+1), k=1..8);
```

1, 2, 5, 14, 42, 132, 429, 1430

```
>
```

```
> Pos:=simplify(Dyck/(1-x*Dyck));
```

$$Pos := -\frac{2}{-1 - \sqrt{1 - 4xy + 2x}}$$

```
> coeftayl(Pos, [x,y]=[0,0], [6,4]);
```

90

```
> S:=0:for l from 0 to 6 do for k from 0 to 6 do  
S:=S+coeftayl(Pos, [x,y]=[0,0], [k,l])*x^k*y^l od  
od:S;
```

$1 + x + xy + 20x^6y^2 + 14x^5y^2 + 5x^3y^3 + 2x^2y^2 + x^3 + 28x^5y^3 + x^4 + x^5$
 $+ x^6 + x^2 + 132x^6y^5 + 2x^2y + 5x^3y^2 + 90x^6y^4 + 42x^5y^5 + 3x^3y$
 $+ 132x^6y^6 + 4x^4y + 14x^4y^4 + 14x^4y^3 + 5x^5y + 9x^4y^2 + 48x^6y^3$
 $+ 42x^5y^4 + 6x^6y$

```
> |
```


Automata and rationality

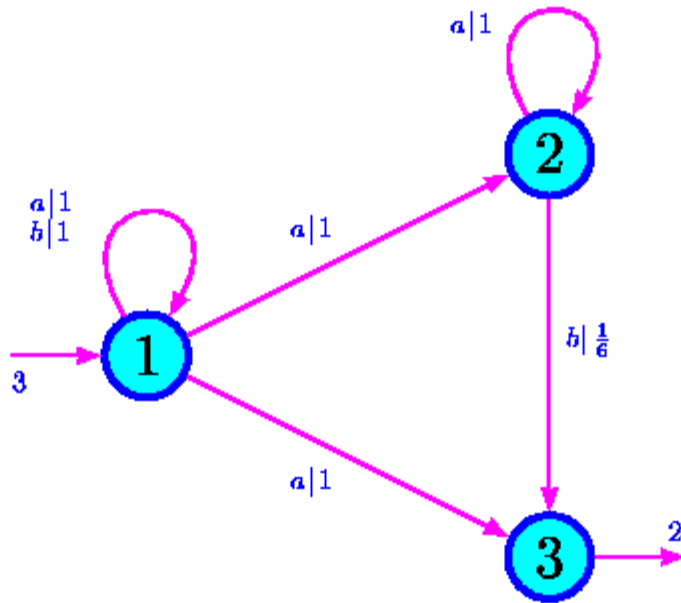


FIG. 1 – Un \mathbb{Q} -automate \mathcal{A} .

Le comportement de \mathcal{A} est :

$$\text{comportement}(\mathcal{A}) = \sum_{a, b \in A} (a + b)^*(6 + a^*b).$$

Un type particulier d'automate à multiplicités est constitué des automates à multiplicités avec des ε -transitions.

Un k - ε -automate \mathcal{A}_ε est un k -automate sur l'alphabet $A_\varepsilon = A \cup \{\varepsilon\}$.

Exemple :

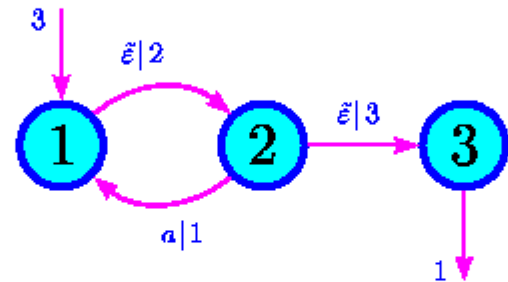
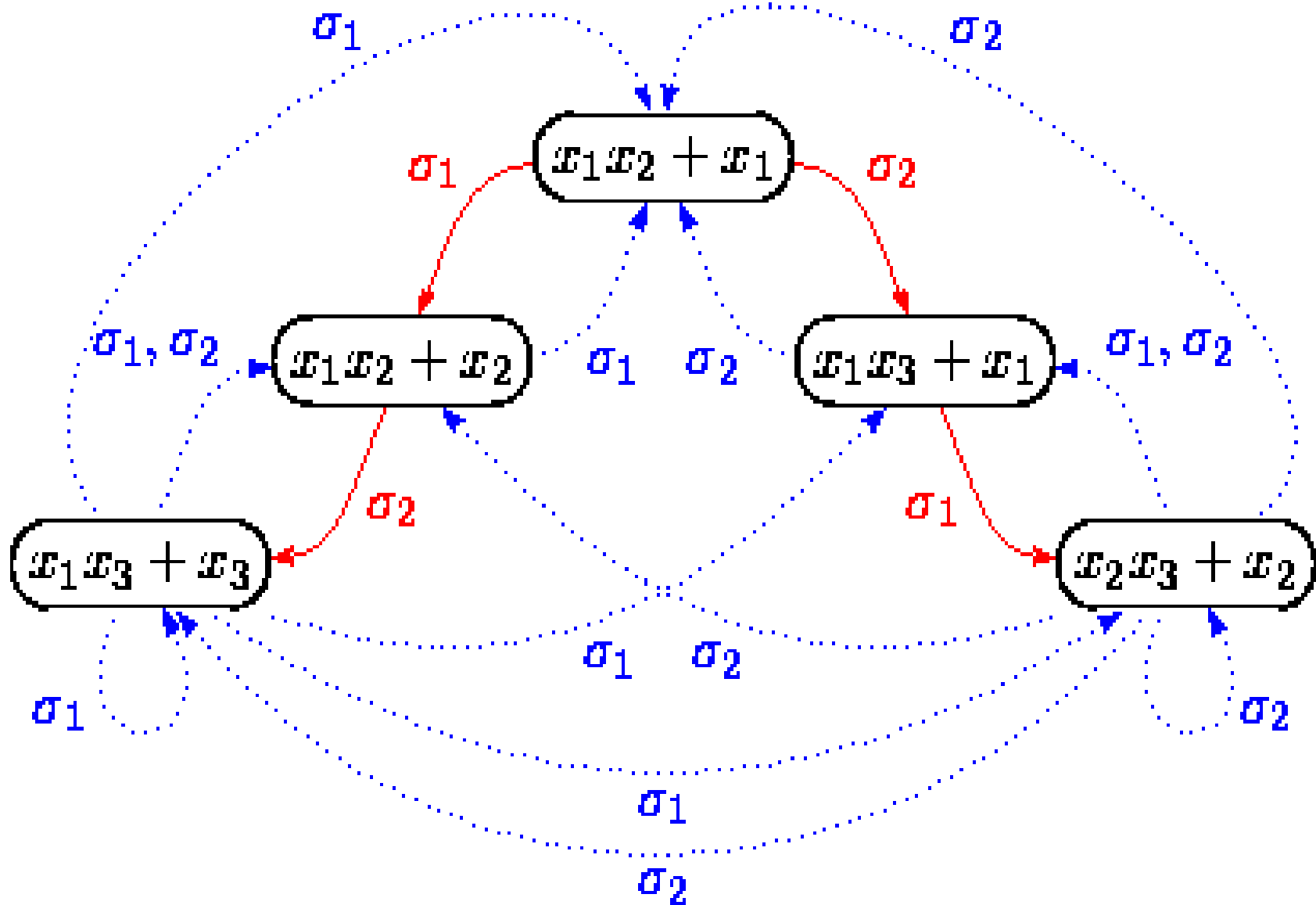


FIG. 2 – Un \mathbb{N} - ε -automate \mathcal{A}_ε

$$\text{comportement}(\mathcal{A}_\varepsilon) = 18\varepsilon \left(\sum_{i \in \mathbb{N}} 2^i (a\varepsilon)^i \right) \varepsilon.$$



A correct implementation of Schelling's model

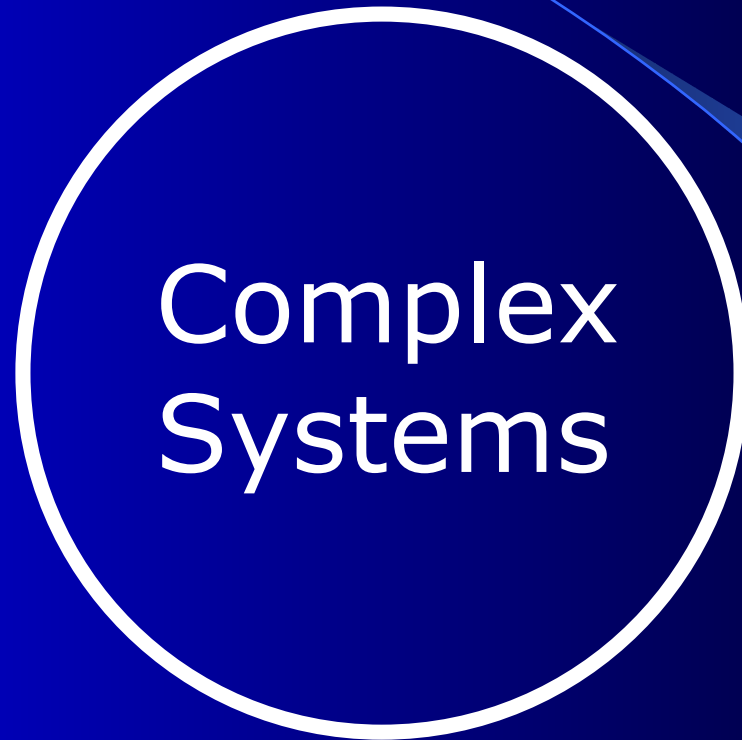
Problem : If one scans the board, addressing the inhabitants one after one, result is sensitive to the order of scanning.

Solution : Invent a (combinatorial) data structure which adapted to the parallel structure of the moving intentions of the inhabitants.

Problem : If one scans the board, addressing the inhabitants one after one, result is sensitive to the order of scanning.

Solution : Invent a (combinatorial) data structure which adapted to the parallel structure of the moving intentions of the inhabitants --> this must be a global model.

Combinatorics (mathematics)



Information
(comp. sci.)

Physics
(class. quant.)

Thank You